# TARANET: Traffic-Analysis Resistant Anonymity at the Network Layer

Chen Chen (CMU), Daniele E. Asoni, Adrian Perrig (ETH Zürich),

David Barrera (Polytechnique Montreal),

George Danezis (UCL), Carmela Troncoso (EPFL)

EuroS&P 2018



Our vision:

Our vision:

**Anonymous communication** 

**Anonymous communication** Non-discrimination

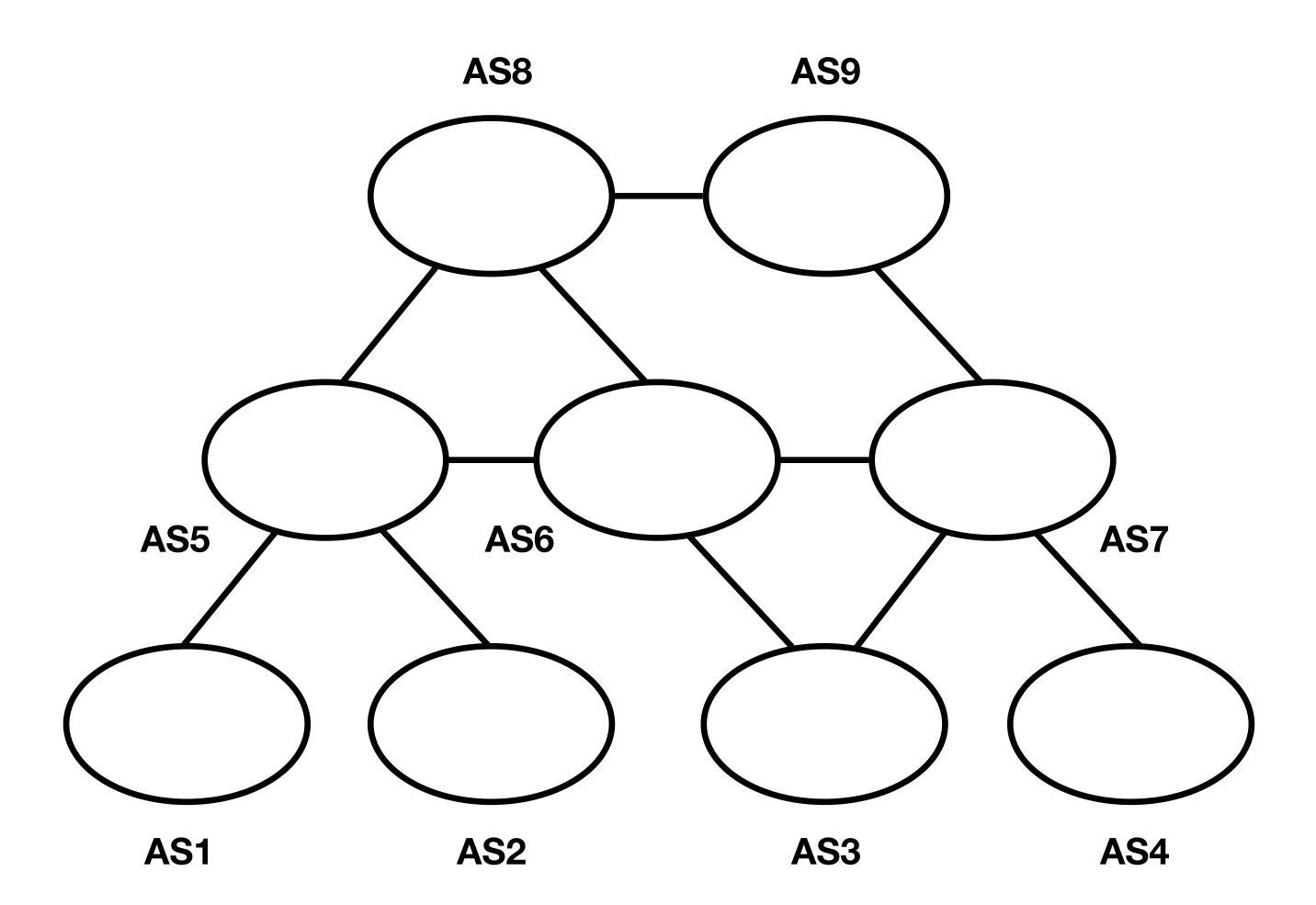
Our vision:

Our vision:

**Anonymous communication** 

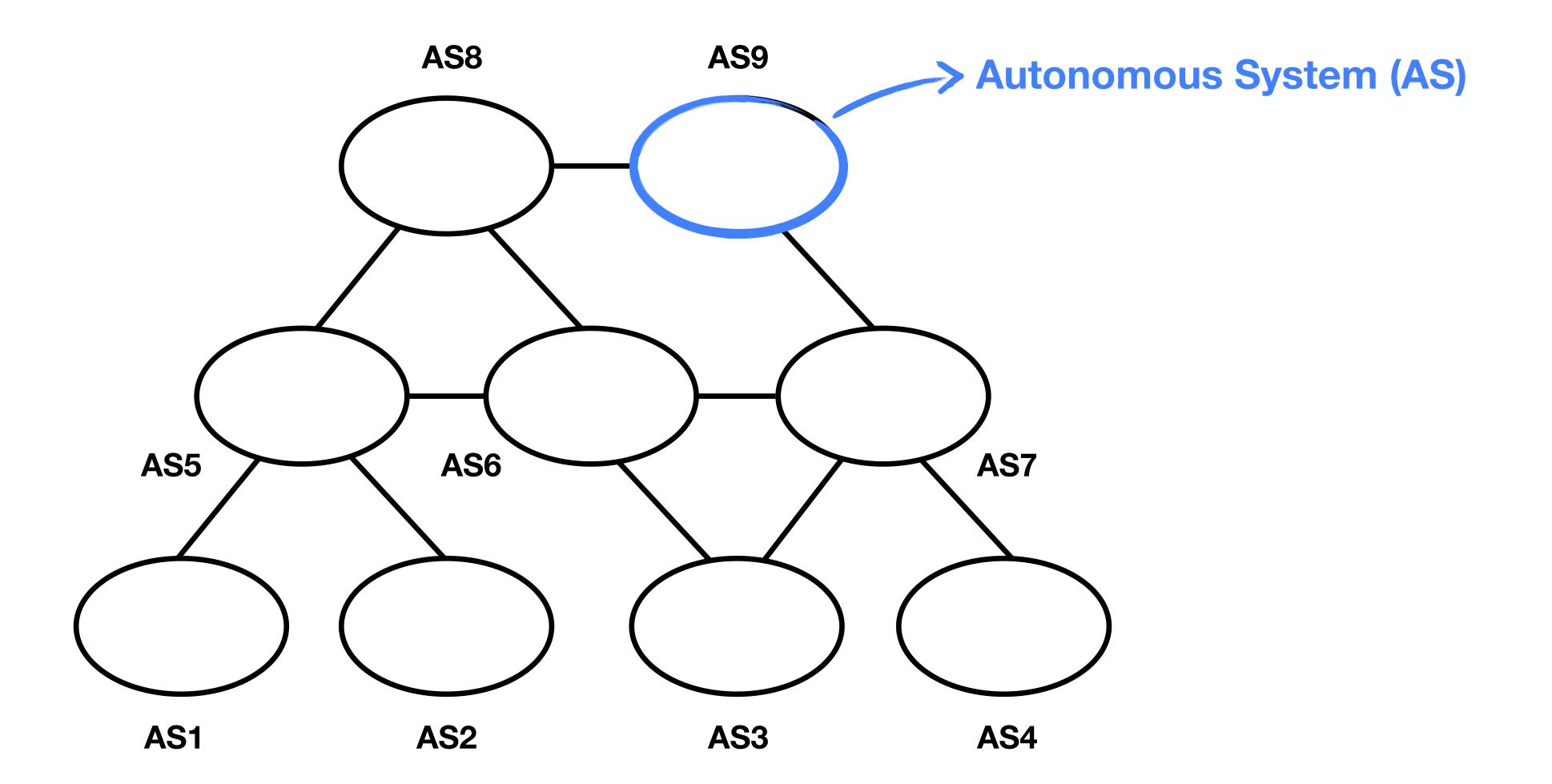
Non-discrimination

Prevent industrial espionage



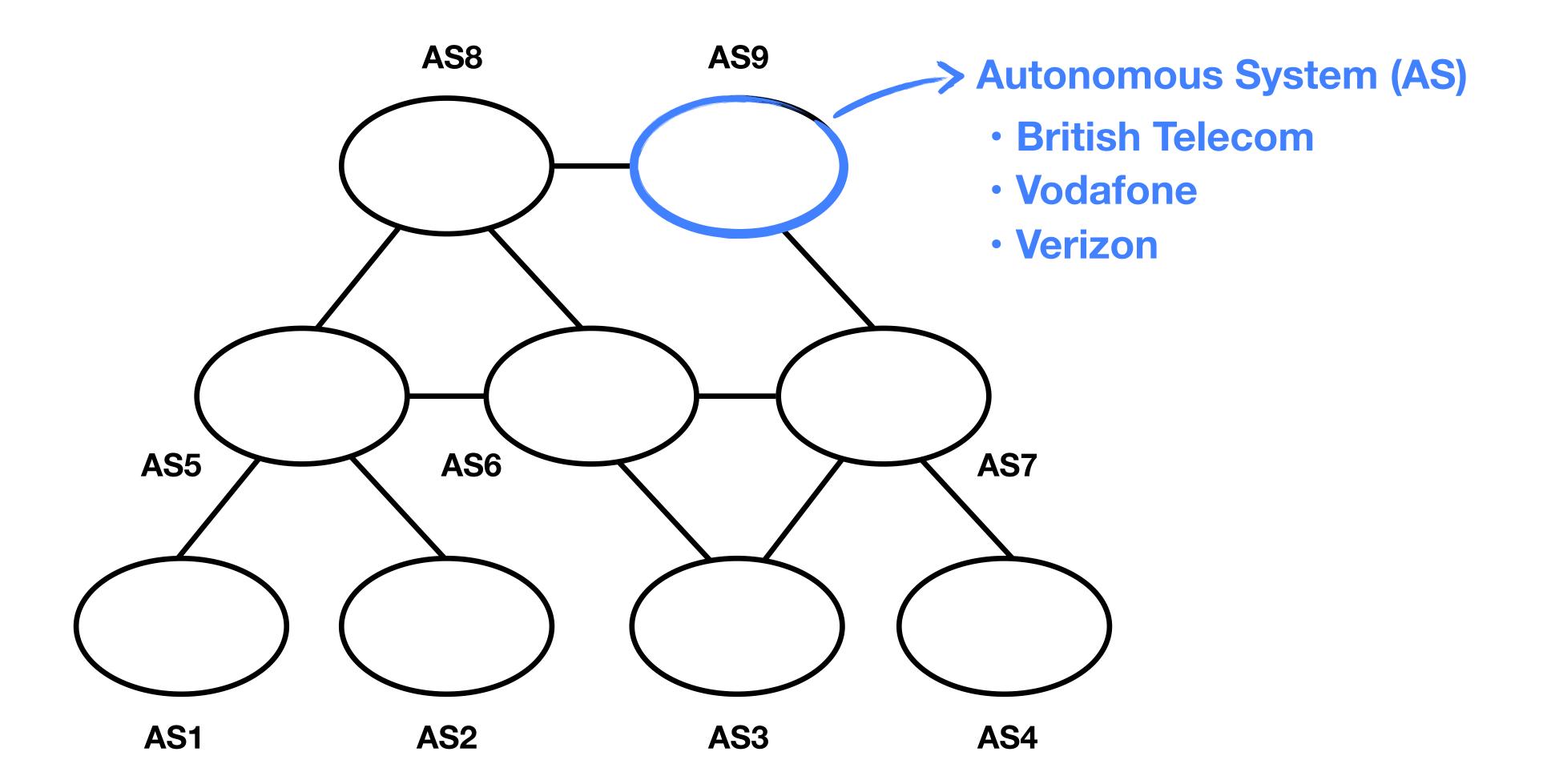






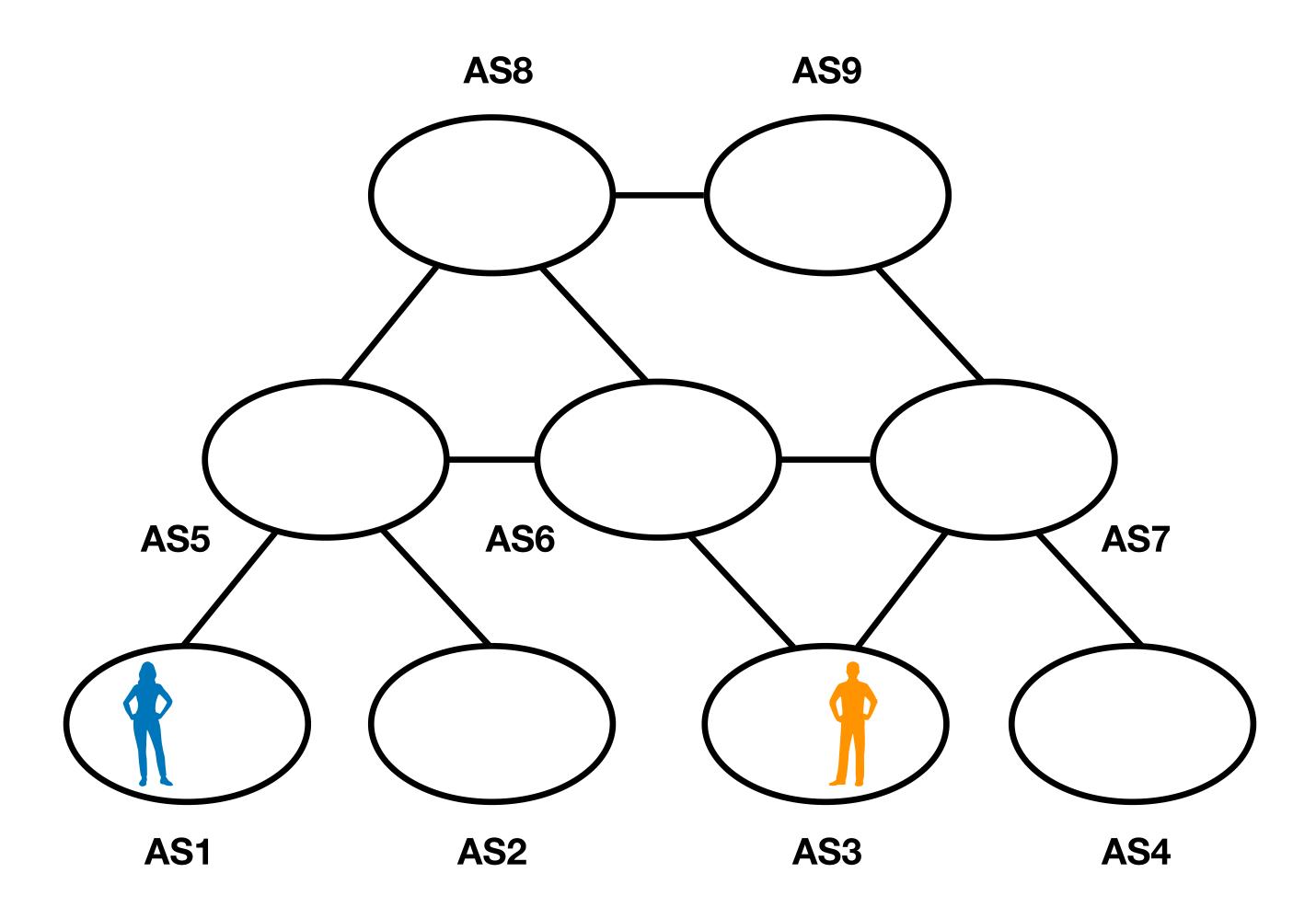






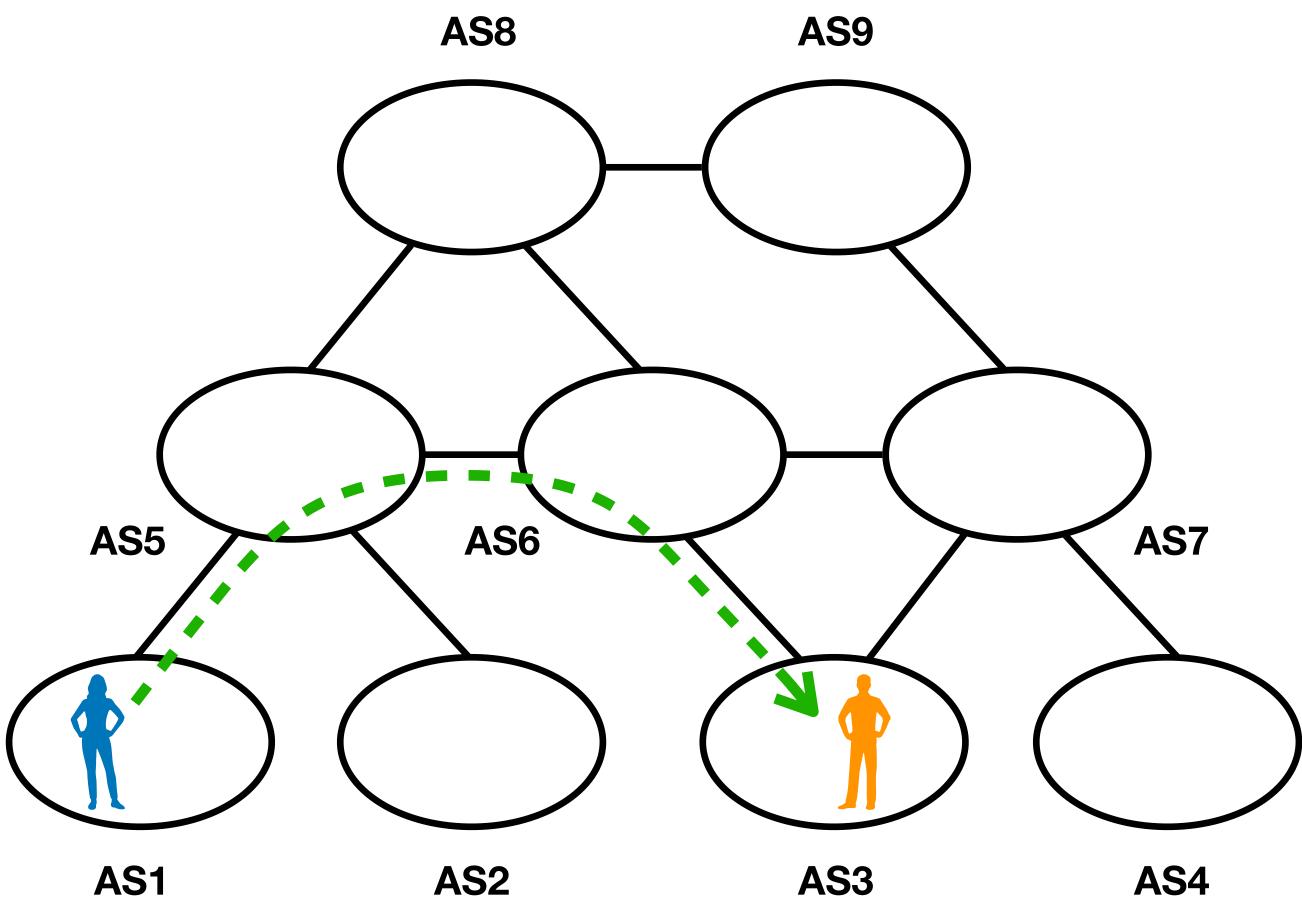






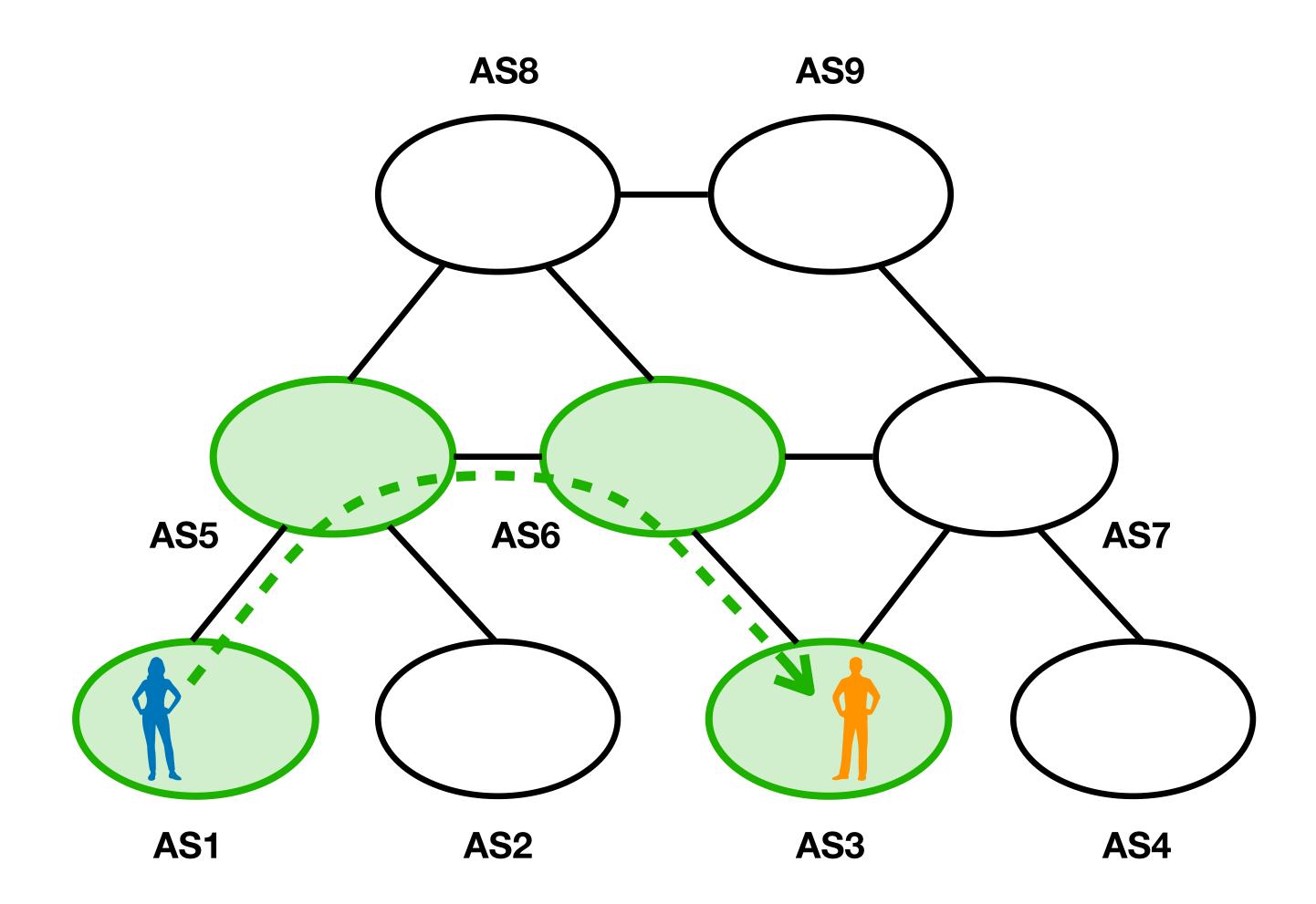






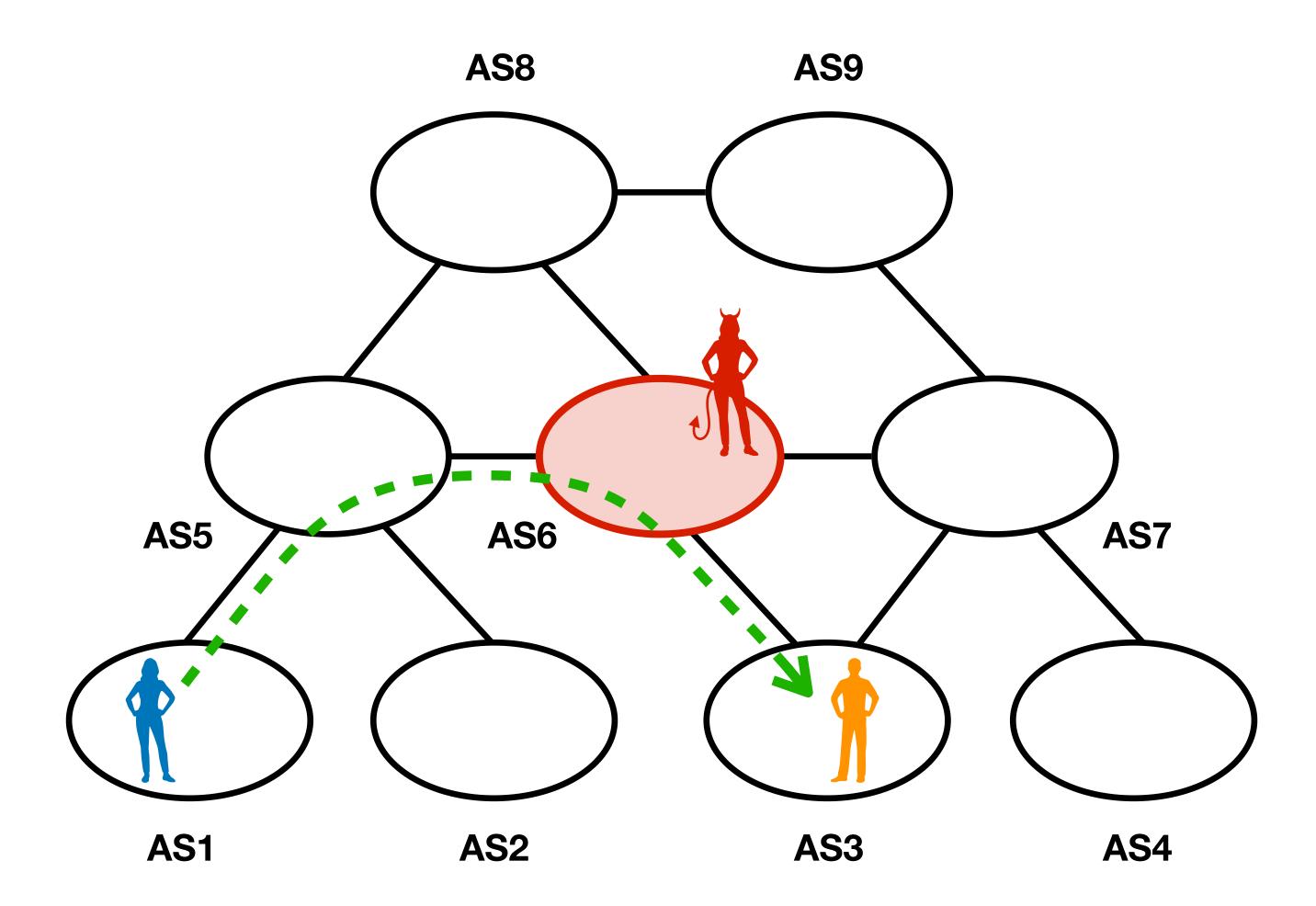






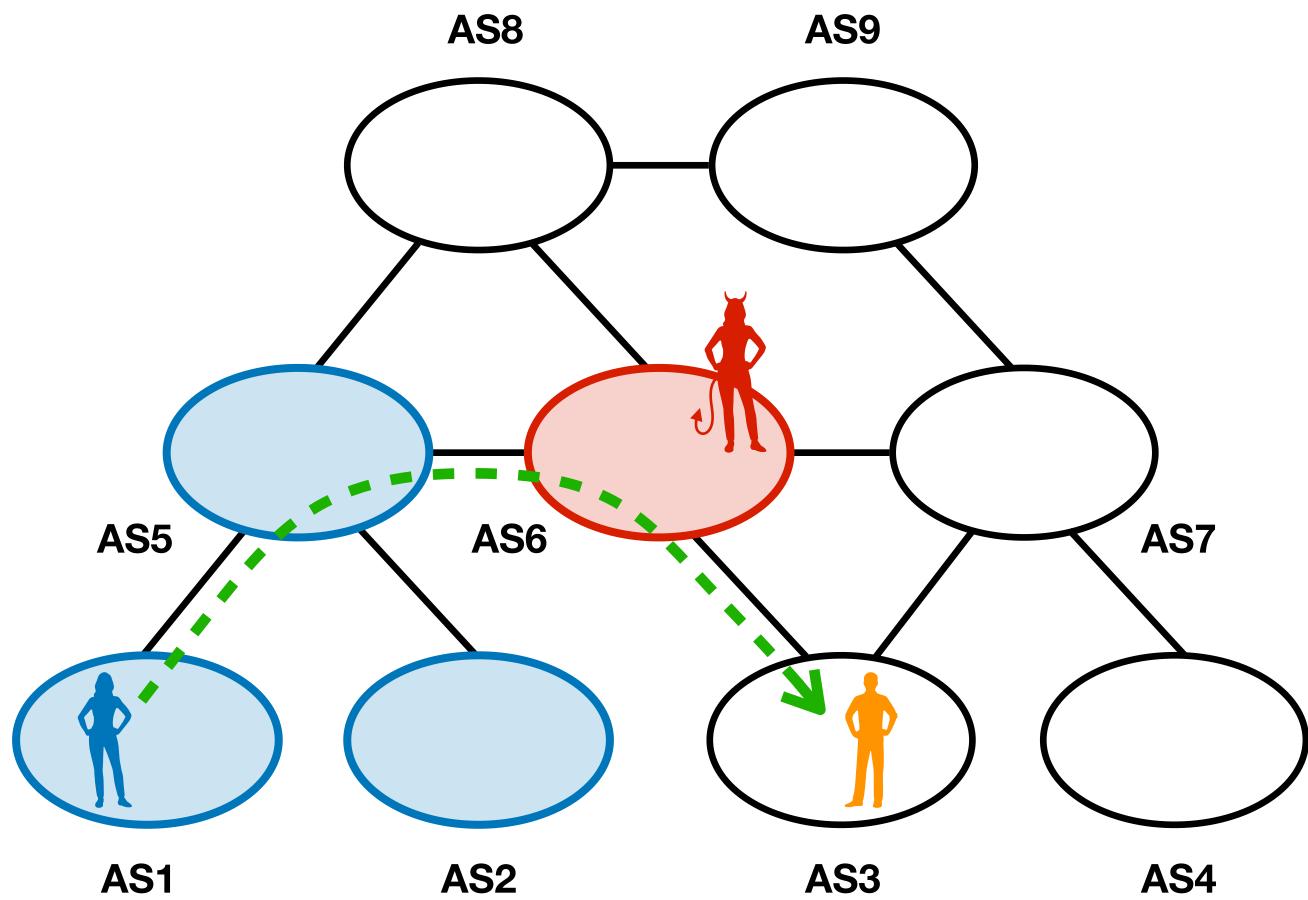






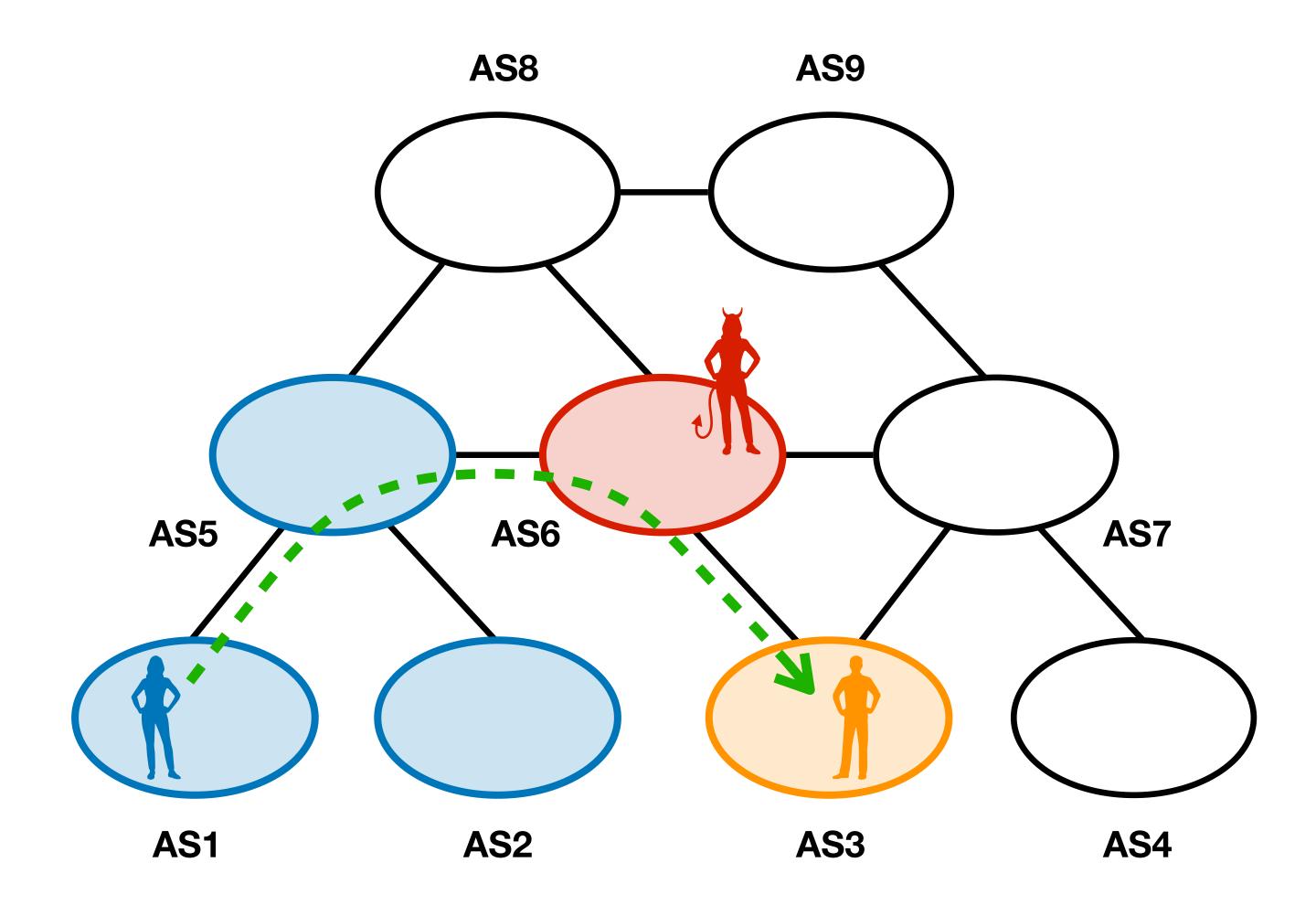








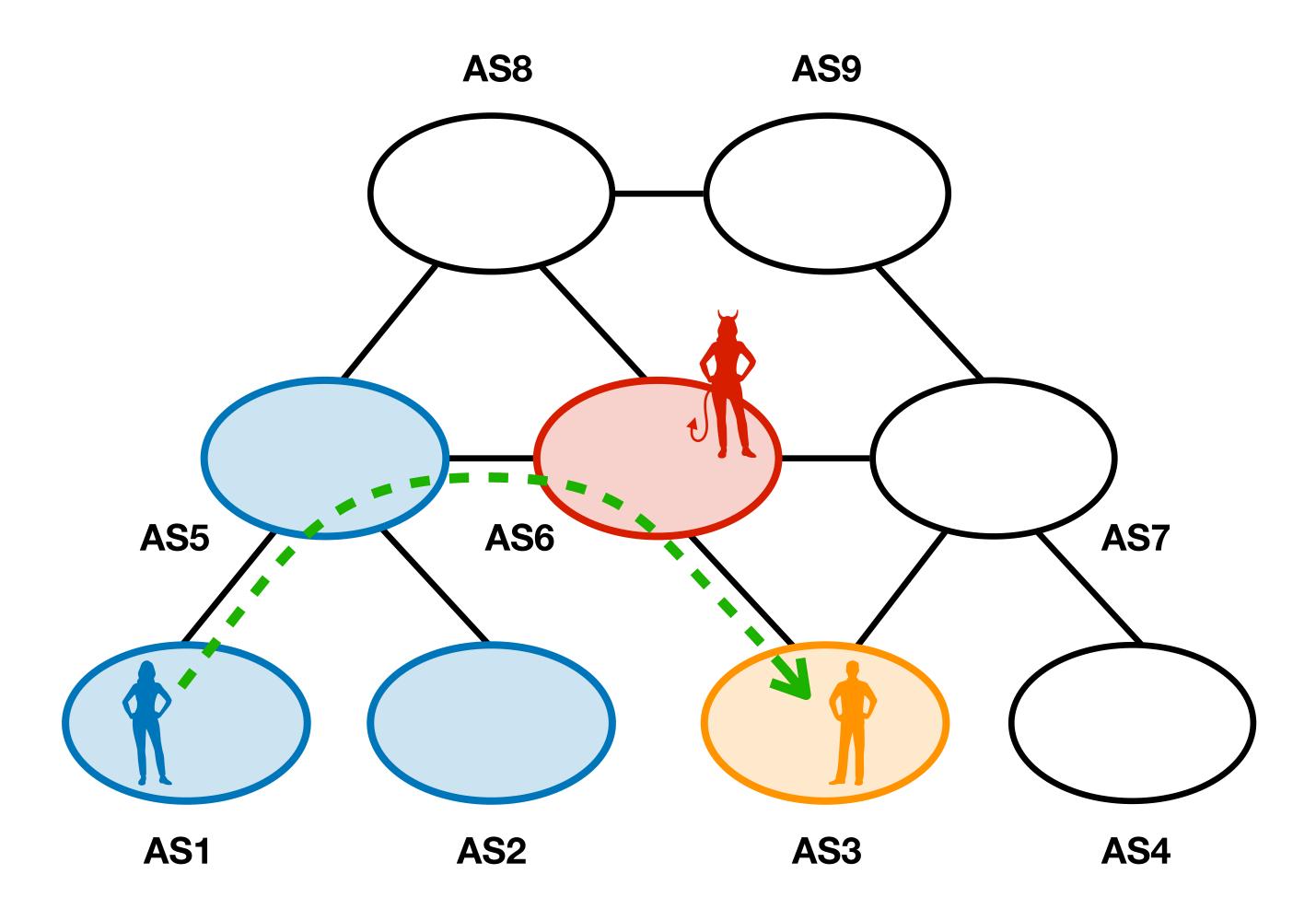






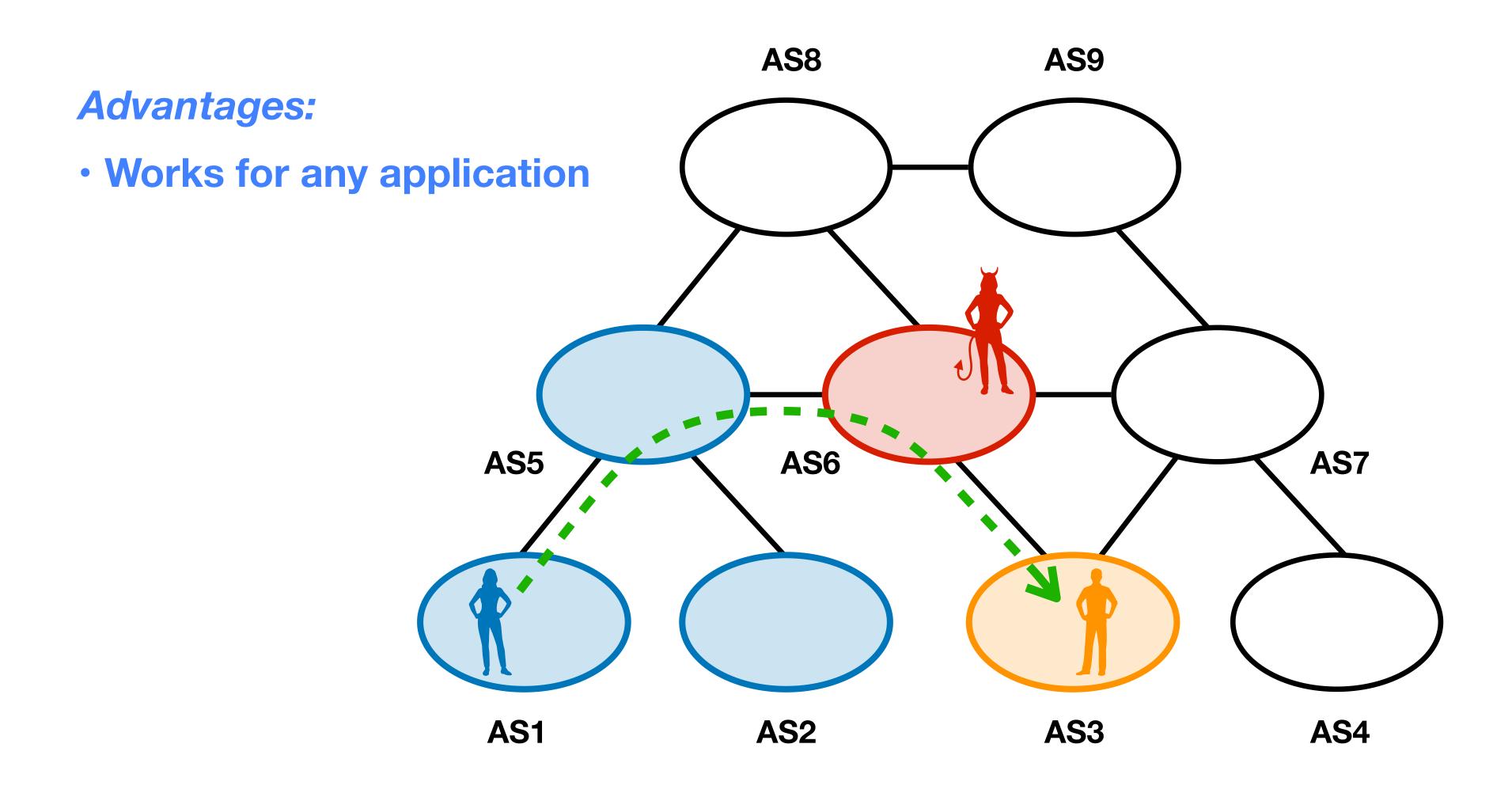






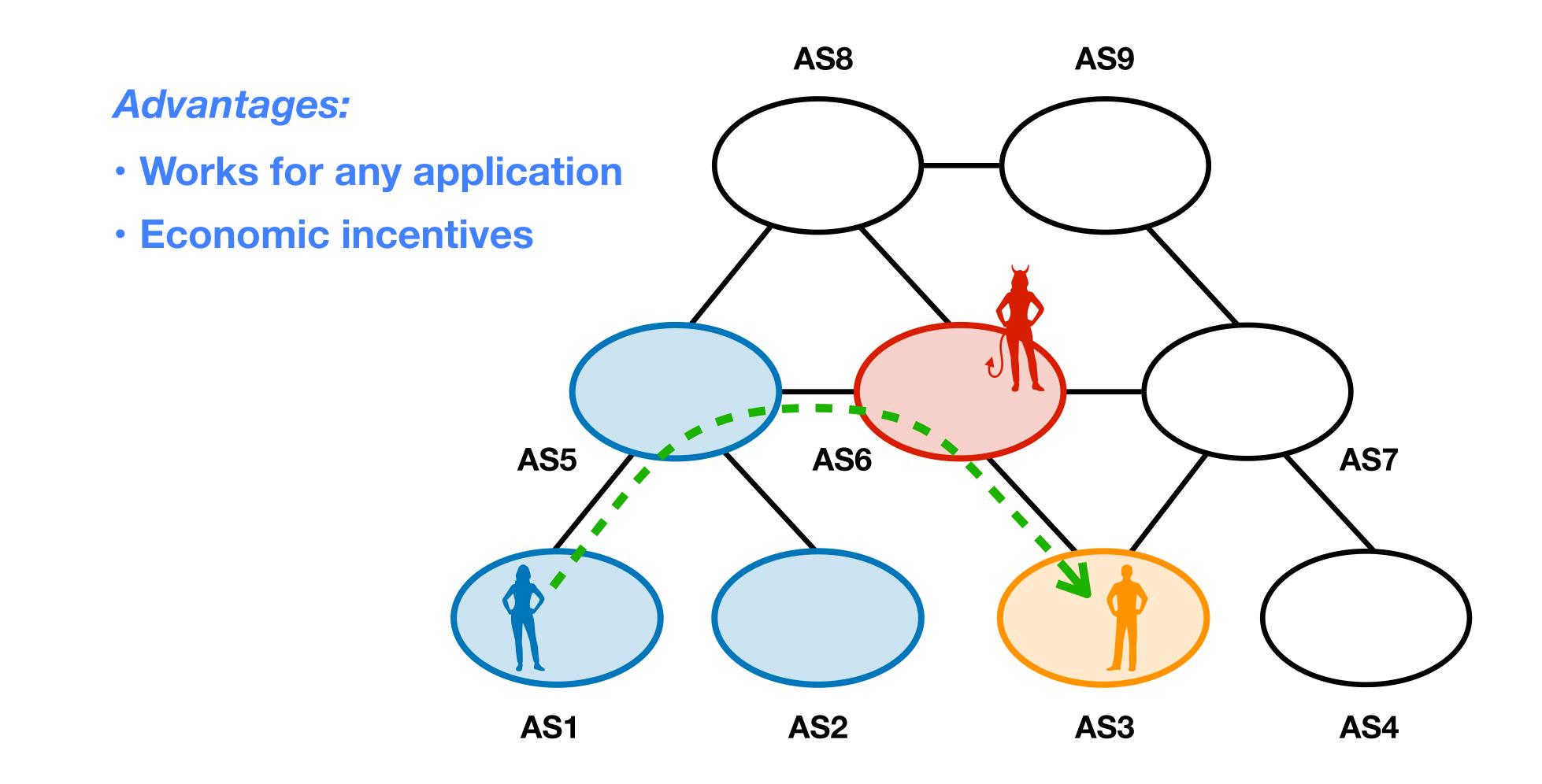






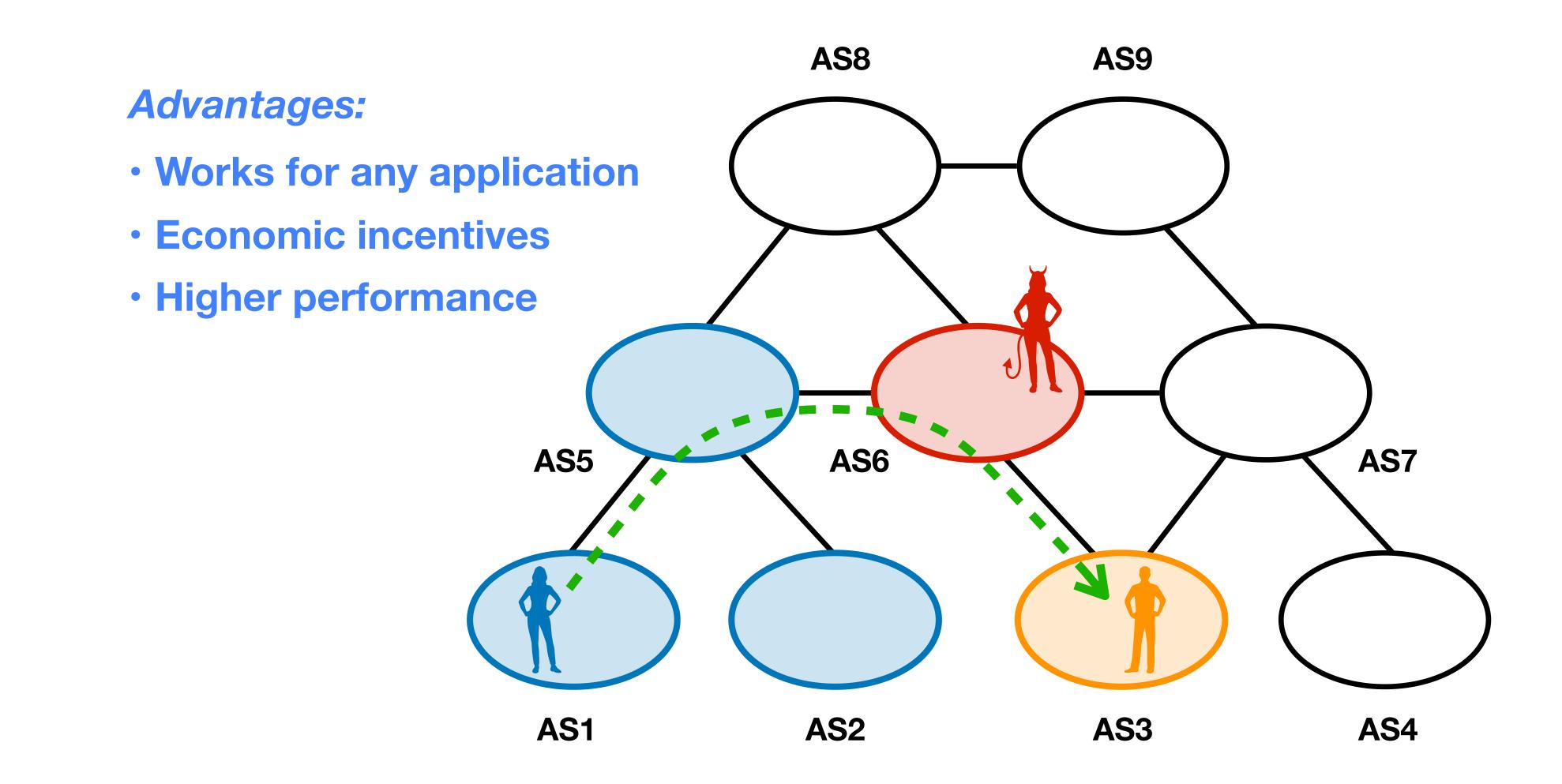








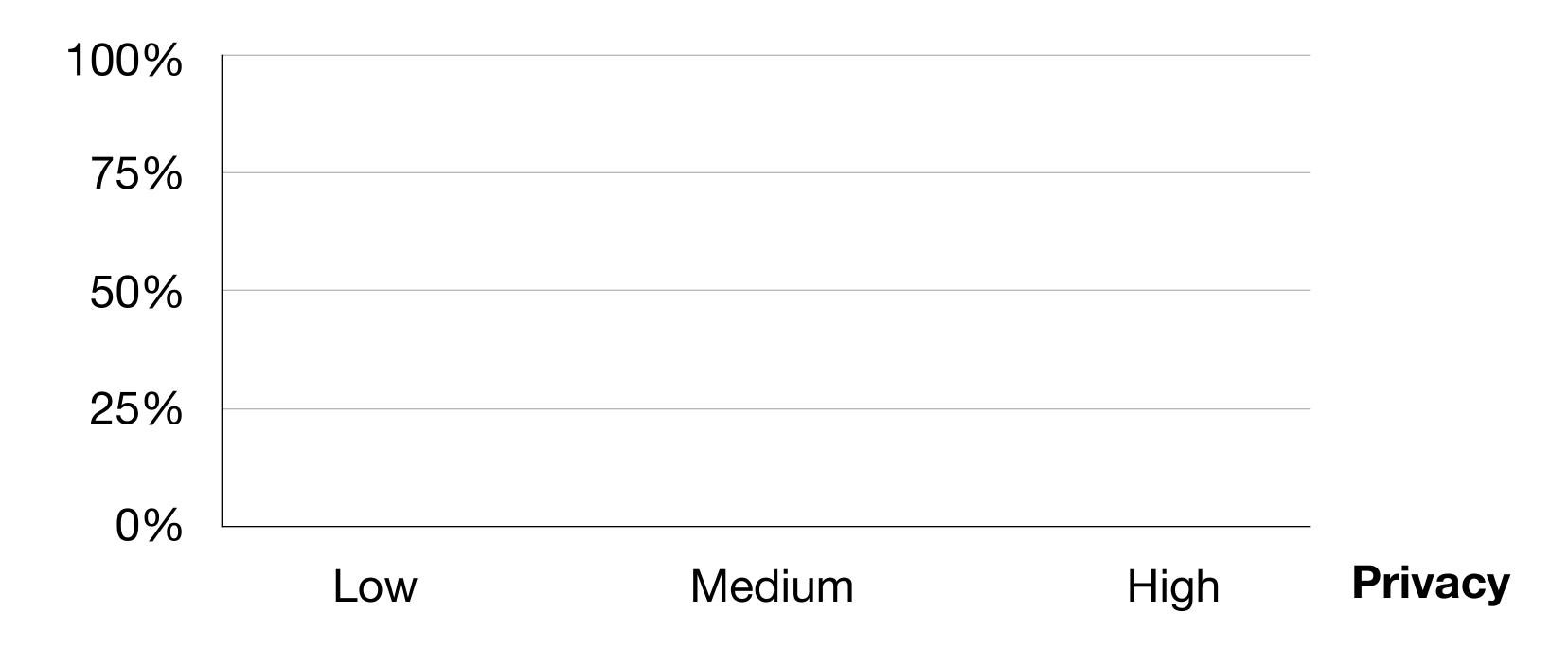








#### Performance

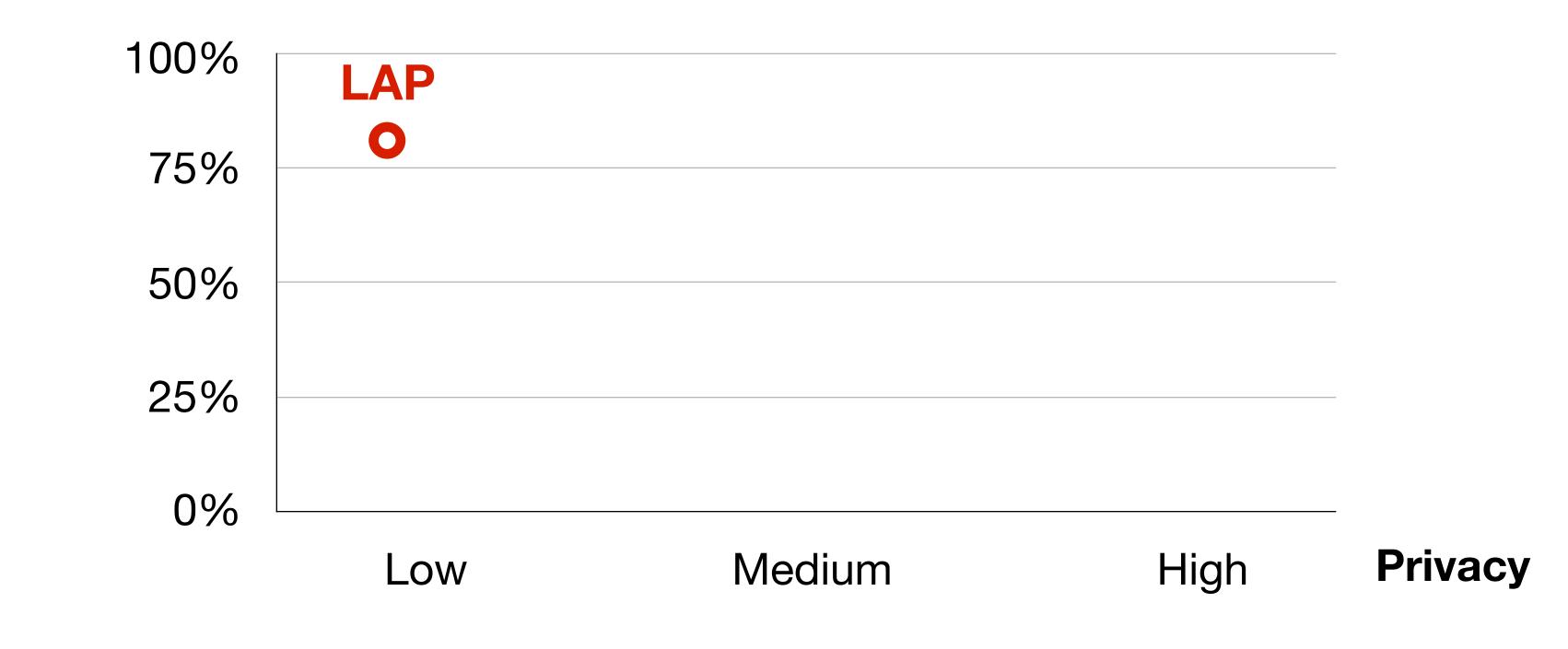








#### Performance



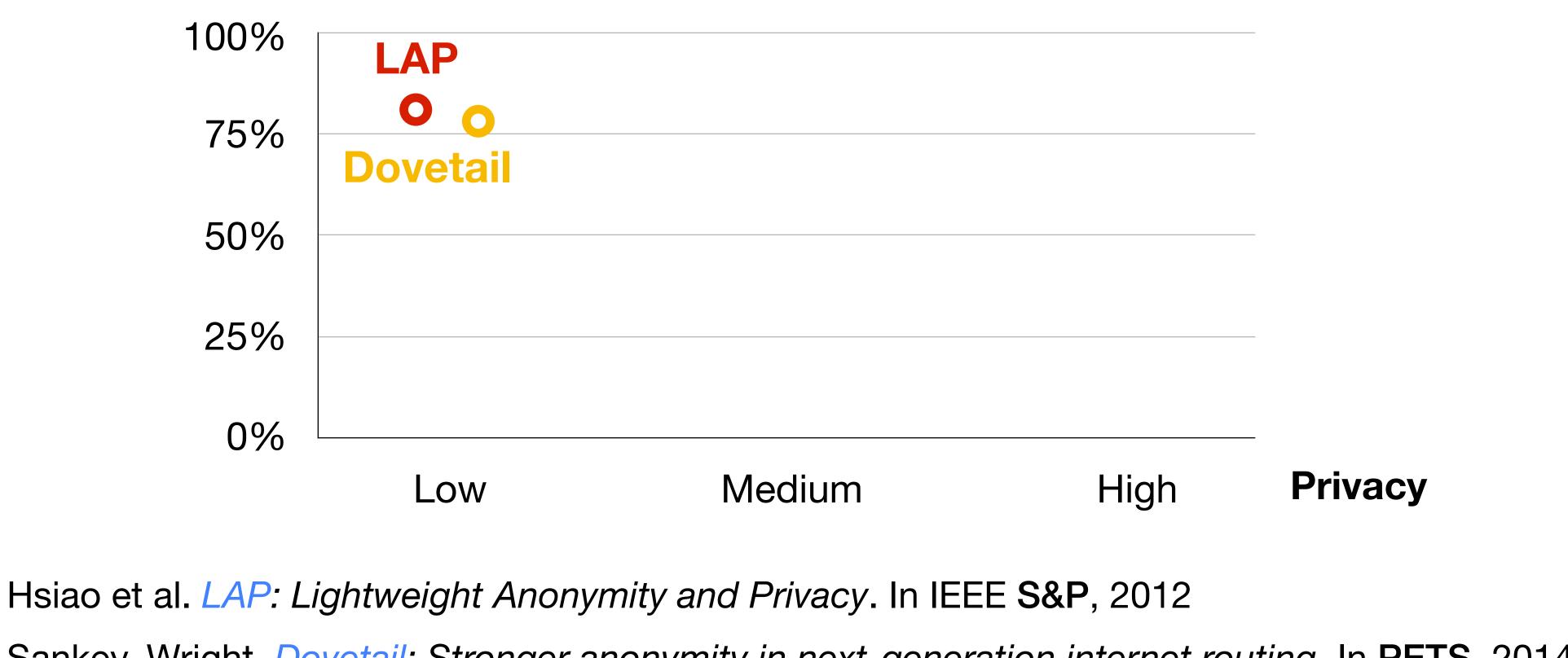
Hsiao et al. LAP: Lightweight Anonymity and Privacy. In IEEE S&P, 2012 ullet







#### Performance



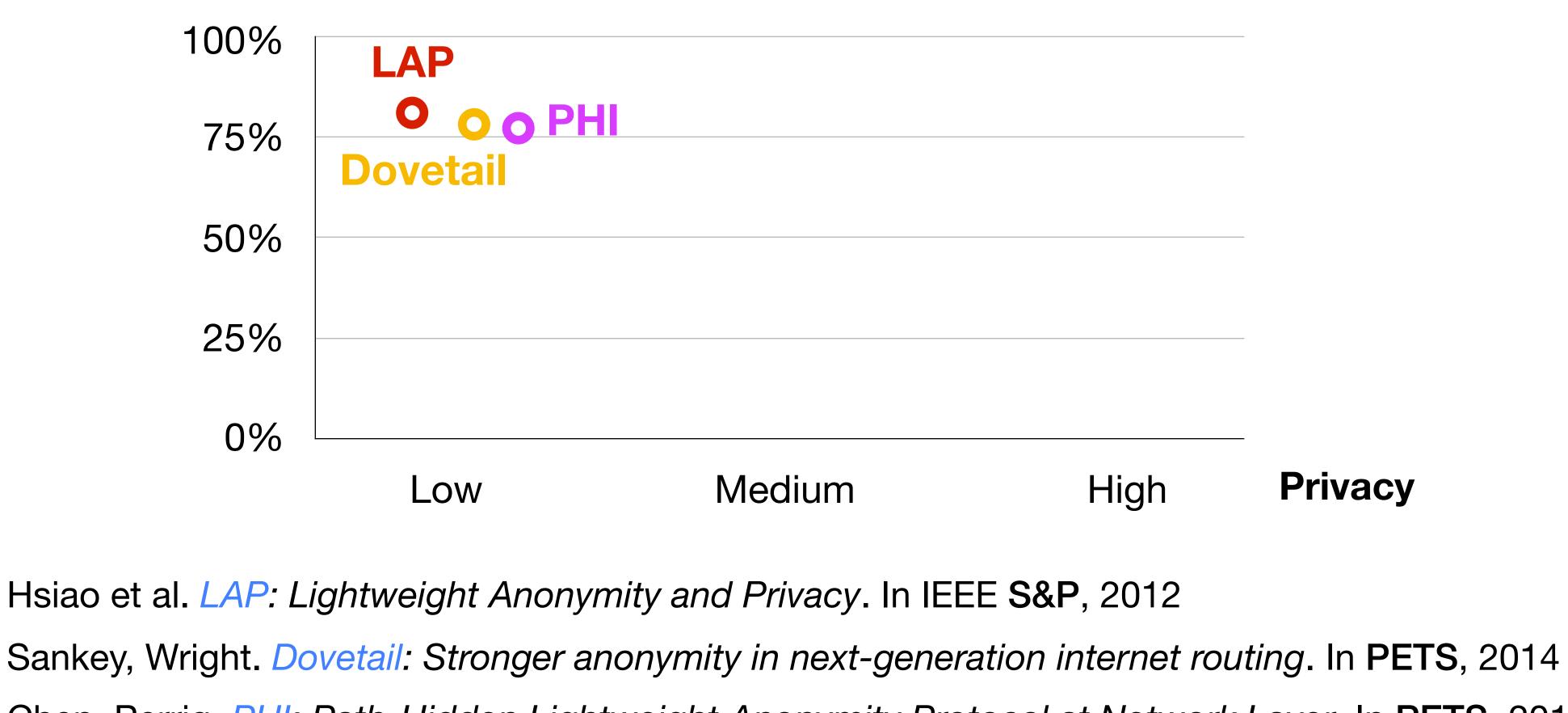
- $\bullet$
- Sankey, Wright. *Dovetail: Stronger anonymity in next-generation internet routing*. In PETS, 2014







### Performance



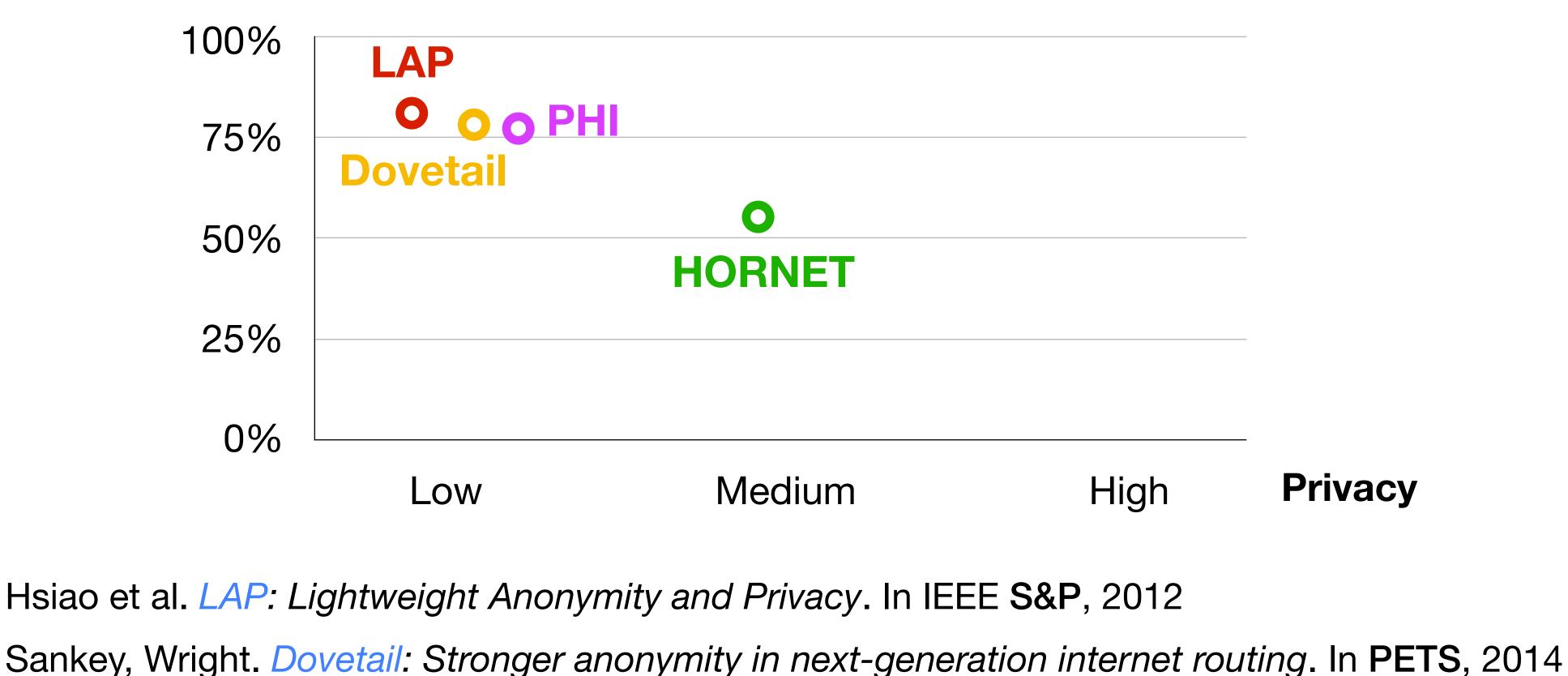
- $\bullet$
- Chen, Perrig. PHI: Path-Hidden Lightweight Anonymity Protocol at Network Layer. In PETS, 2017  $\bullet$







### Performance



- $\bullet$
- $\bullet$
- $\bullet$



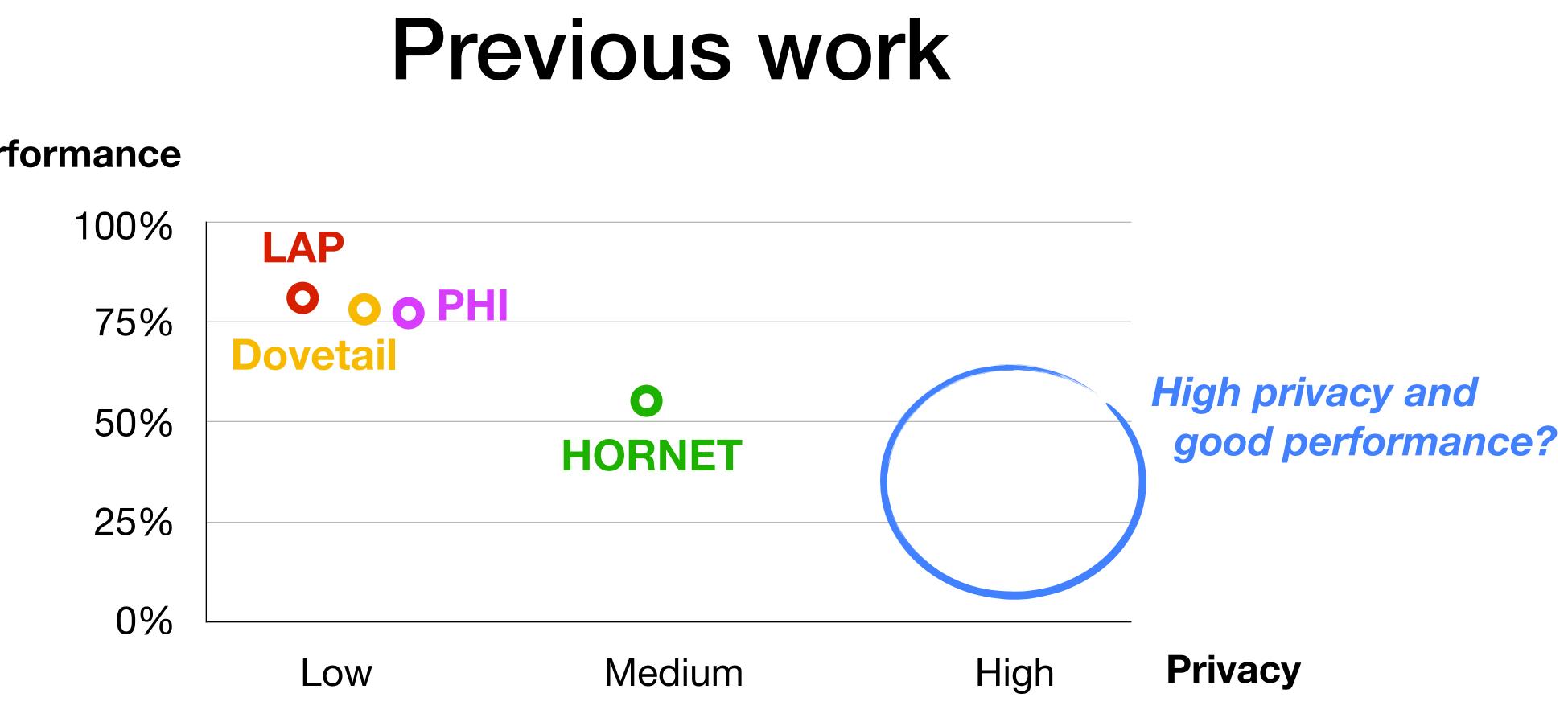
Chen, Perrig. PHI: Path-Hidden Lightweight Anonymity Protocol at Network Layer. In PETS, 2017

Chen et al. HORNET: High-speed Onion Routing at the Network Layer. In ACM CCS, 2015





### Performance



- Hsiao et al. LAP: Lightweight Anonymity and Privacy. In IEEE S&P, 2012  $\bullet$
- $\bullet$
- $\bullet$



Sankey, Wright. *Dovetail: Stronger anonymity in next-generation internet routing*. In PETS, 2014

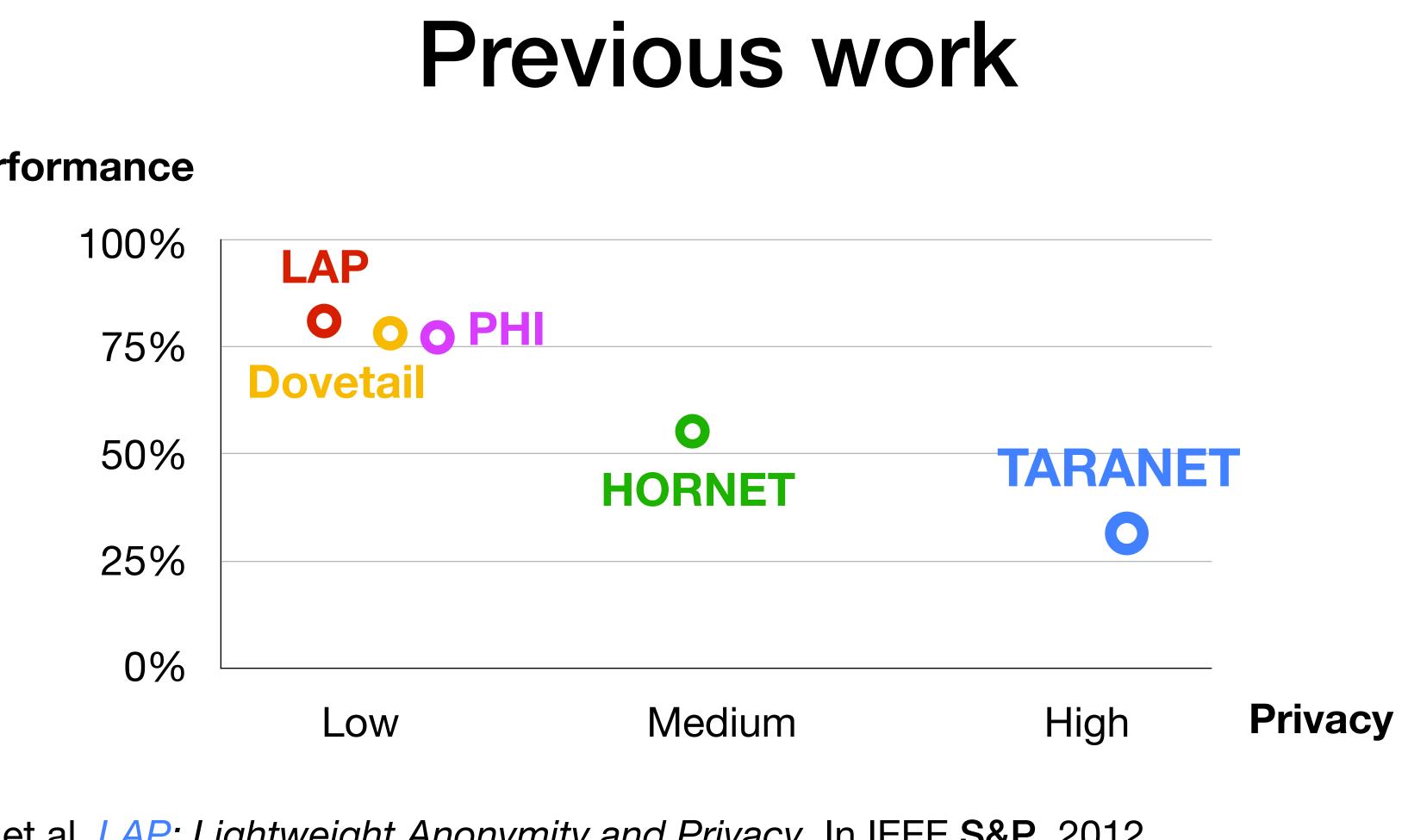
Chen, Perrig. PHI: Path-Hidden Lightweight Anonymity Protocol at Network Layer. In PETS, 2017

Chen et al. HORNET: High-speed Onion Routing at the Network Layer. In ACM CCS, 2015





### Performance



- Hsiao et al. LAP: Lightweight Anonymity and Privacy. In IEEE S&P, 2012  $\bullet$
- $\bullet$
- Chen et al. HORNET: High-speed Onion Routing at the Network Layer. In ACM CCS, 2015  $\bullet$

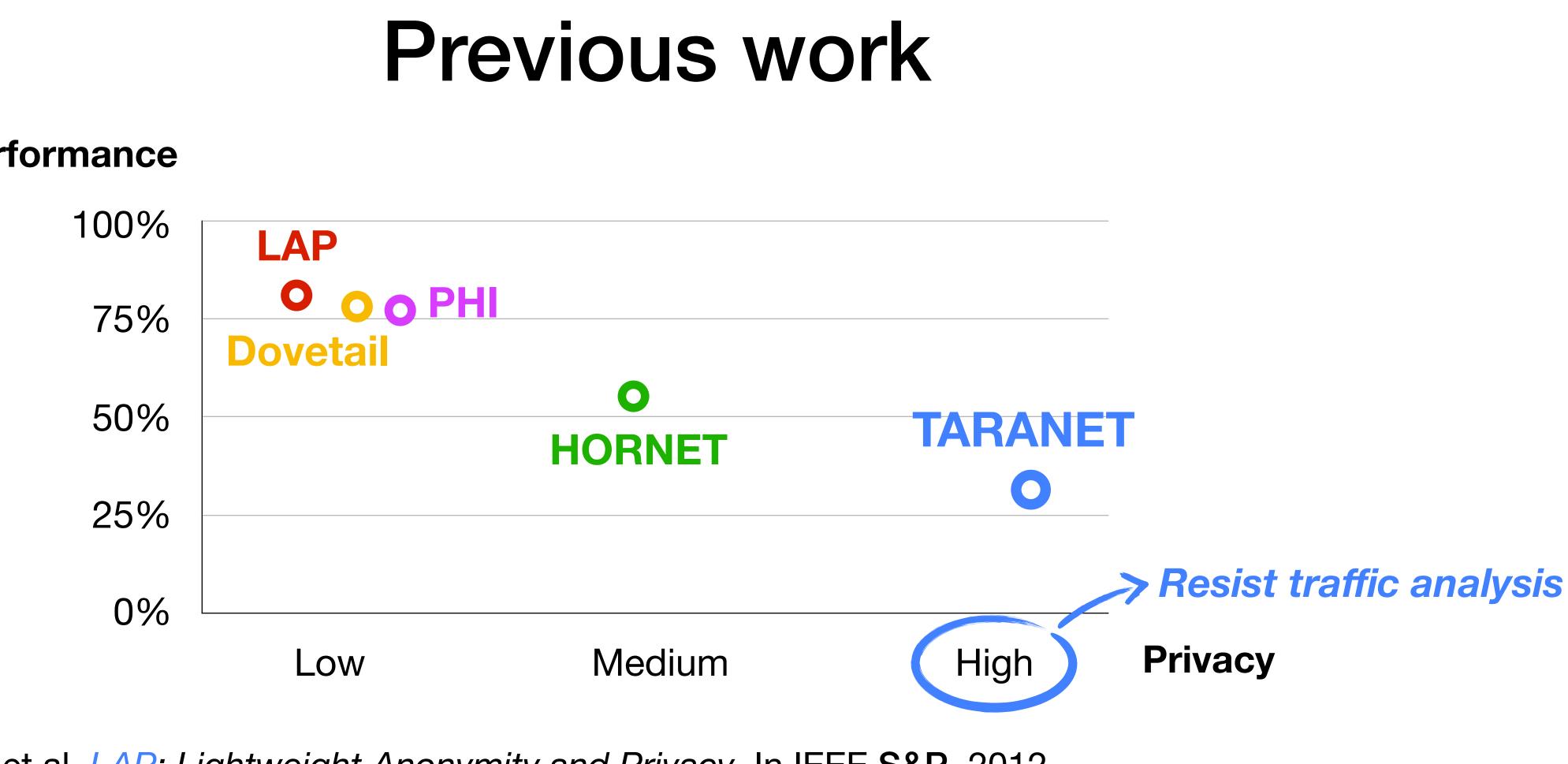


Sankey, Wright. *Dovetail: Stronger anonymity in next-generation internet routing*. In PETS, 2014 Chen, Perrig. PHI: Path-Hidden Lightweight Anonymity Protocol at Network Layer. In PETS, 2017





### Performance



- Hsiao et al. LAP: Lightweight Anonymity and Privacy. In IEEE S&P, 2012  $\bullet$
- $\bullet$
- $\bullet$



Sankey, Wright. *Dovetail: Stronger anonymity in next-generation internet routing*. In PETS, 2014

Chen, Perrig. PHI: Path-Hidden Lightweight Anonymity Protocol at Network Layer. In PETS, 2017

Chen et al. HORNET: High-speed Onion Routing at the Network Layer. In ACM CCS, 2015

### Traffic analysis Severe threat to anonymous communication



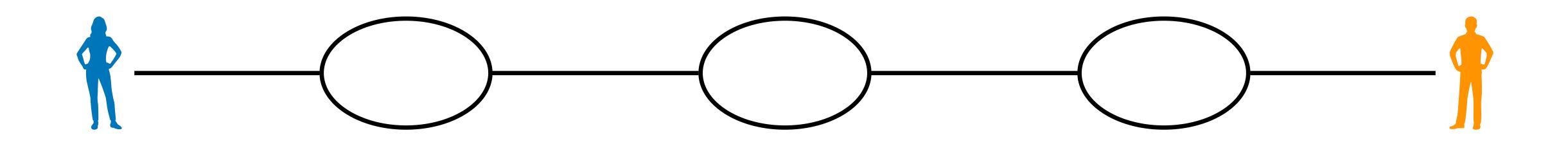
anonymity





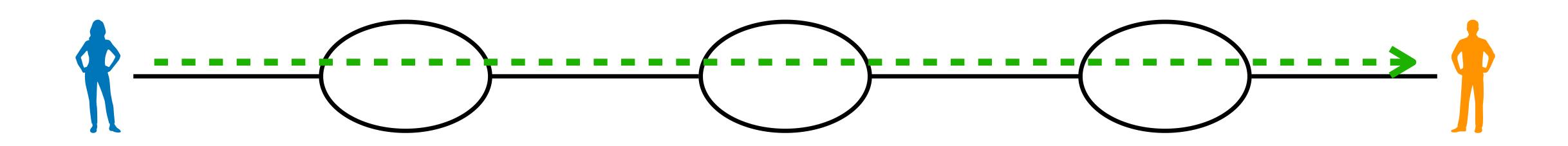






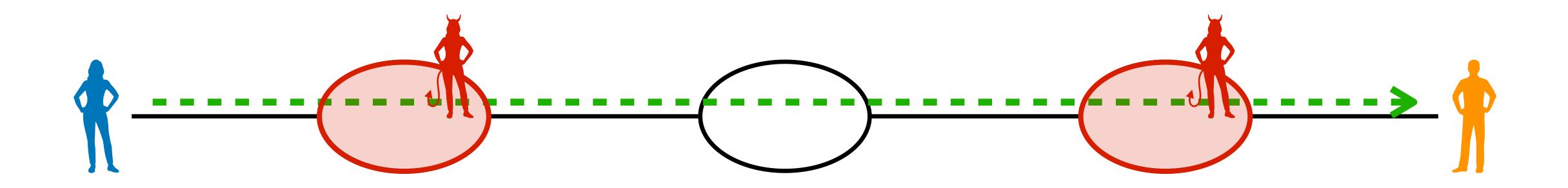










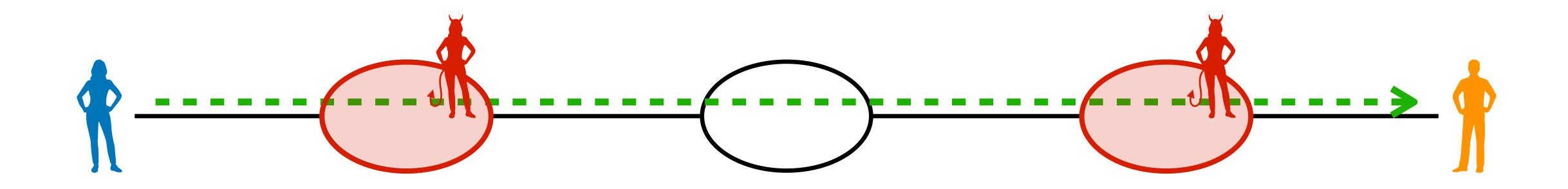






### **Starting point**

Layered encryption

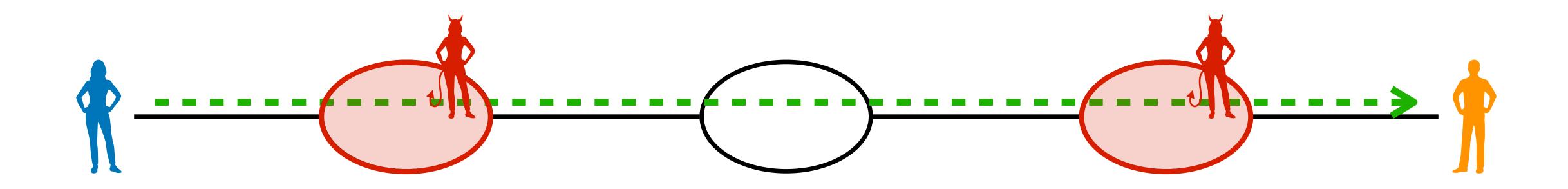






### **Starting point**

- Layered encryption
- Per-hop authentication

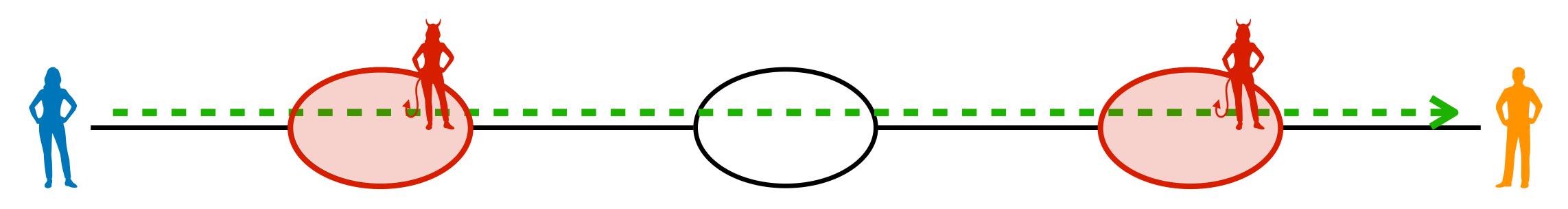






### **Starting point**

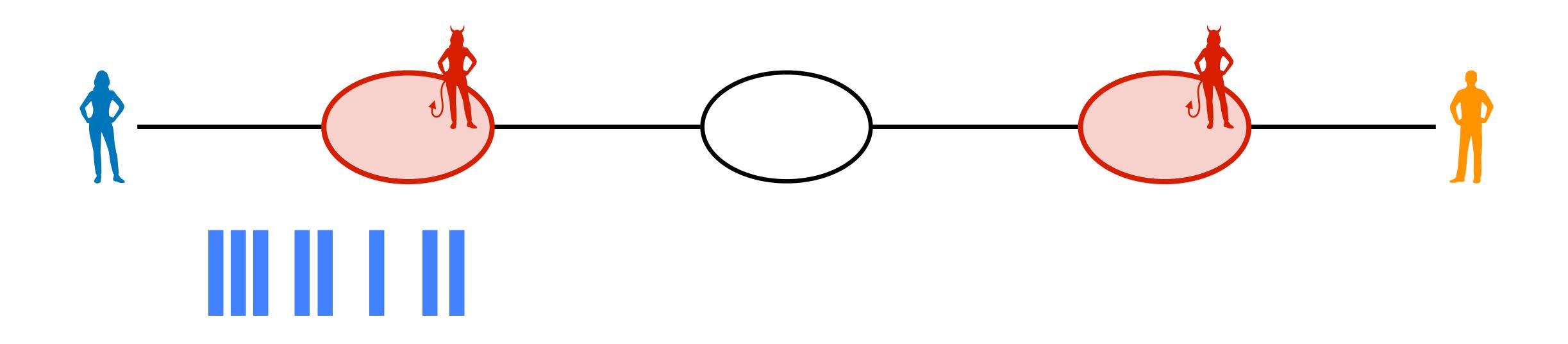
- Layered encryption
- Per-hop authentication
- Fixed packet length







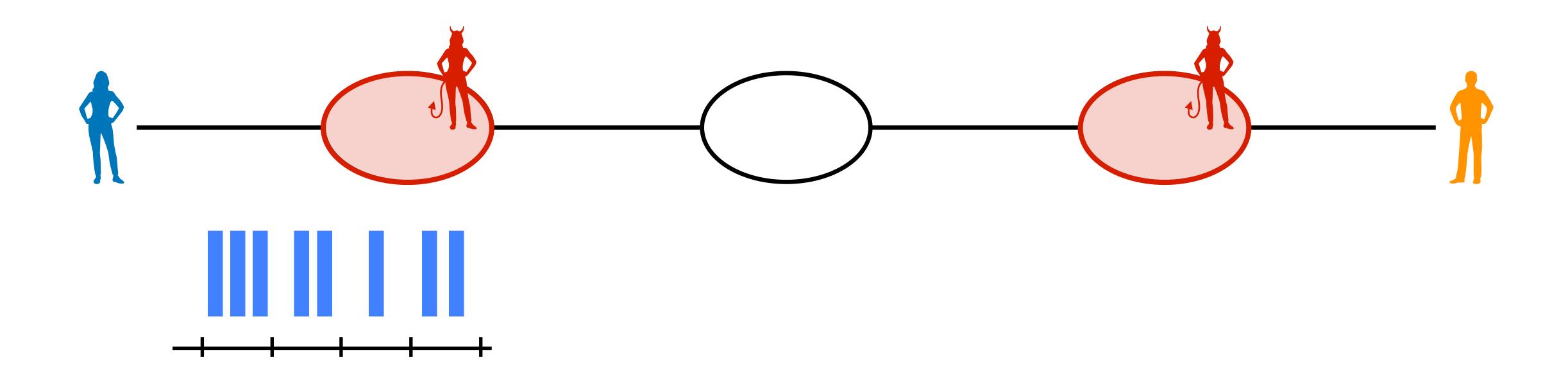






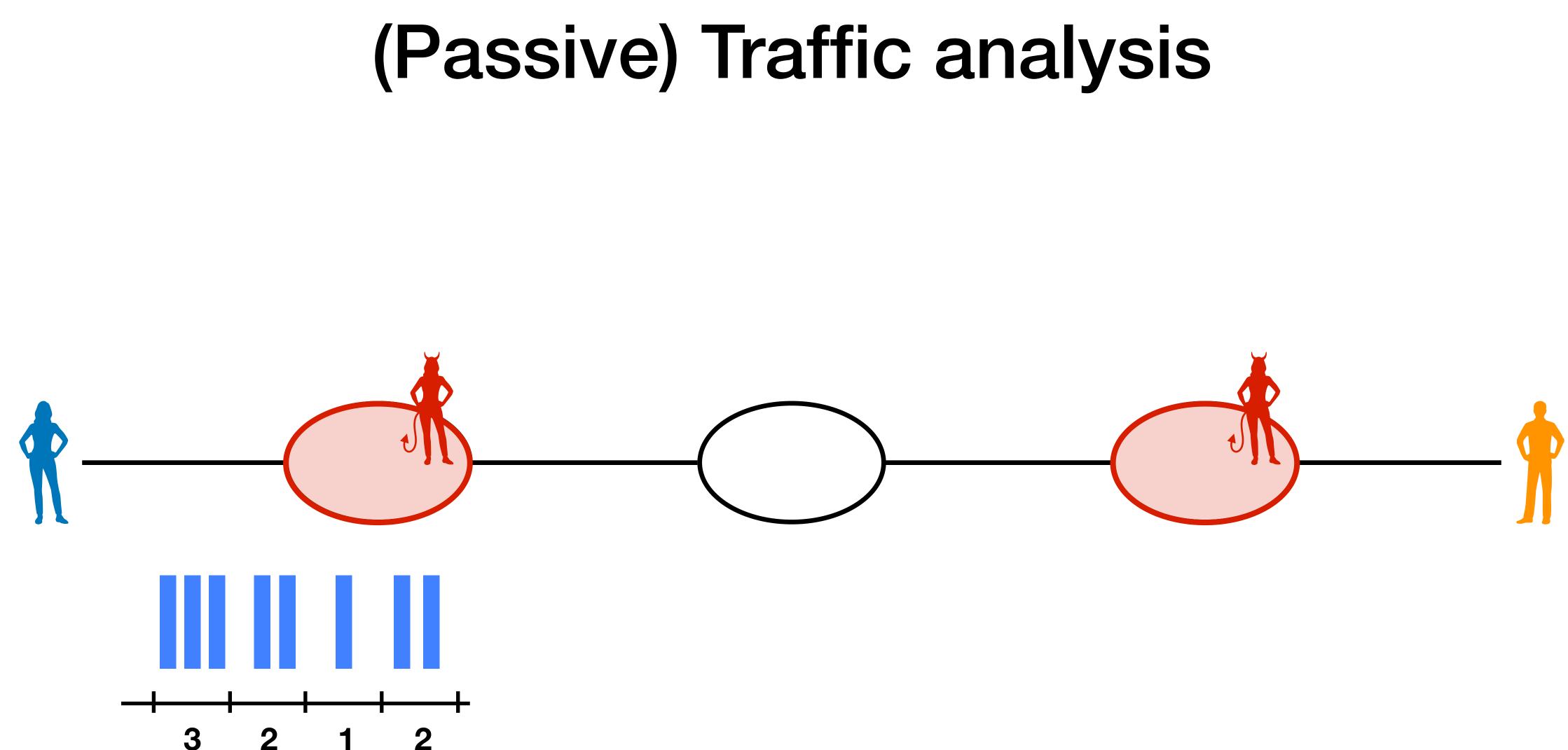






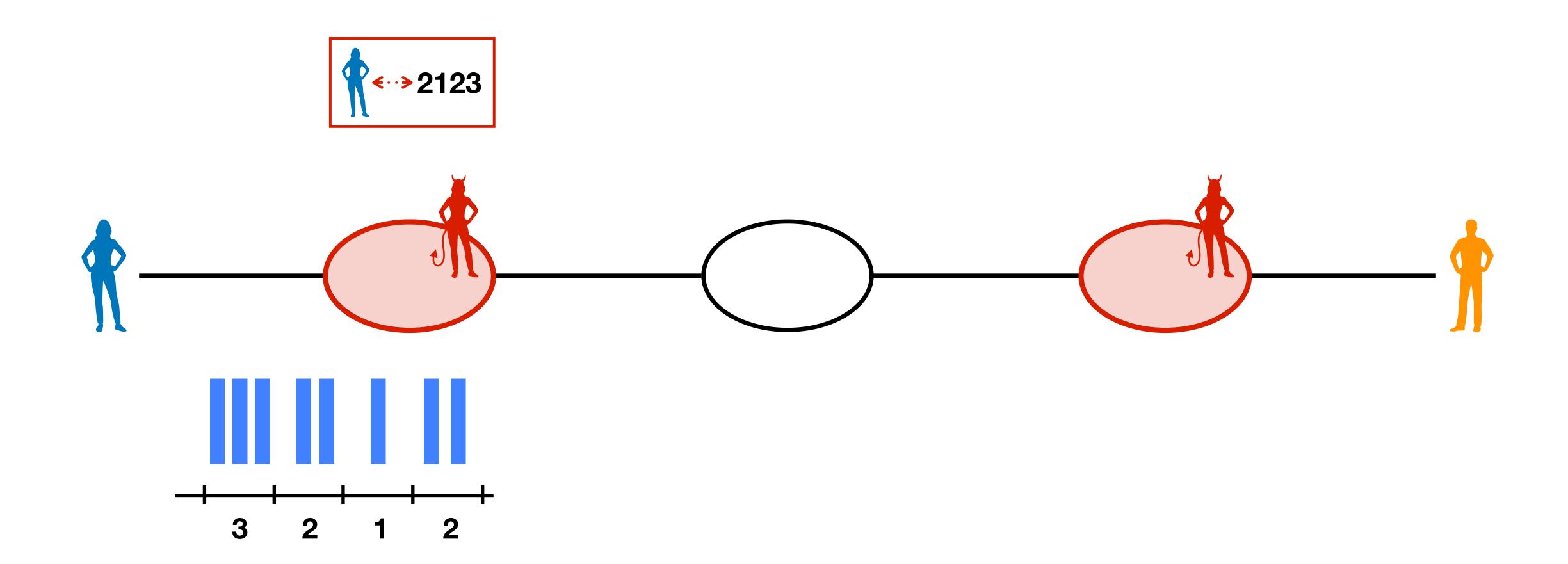






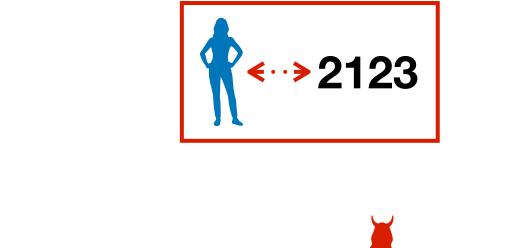


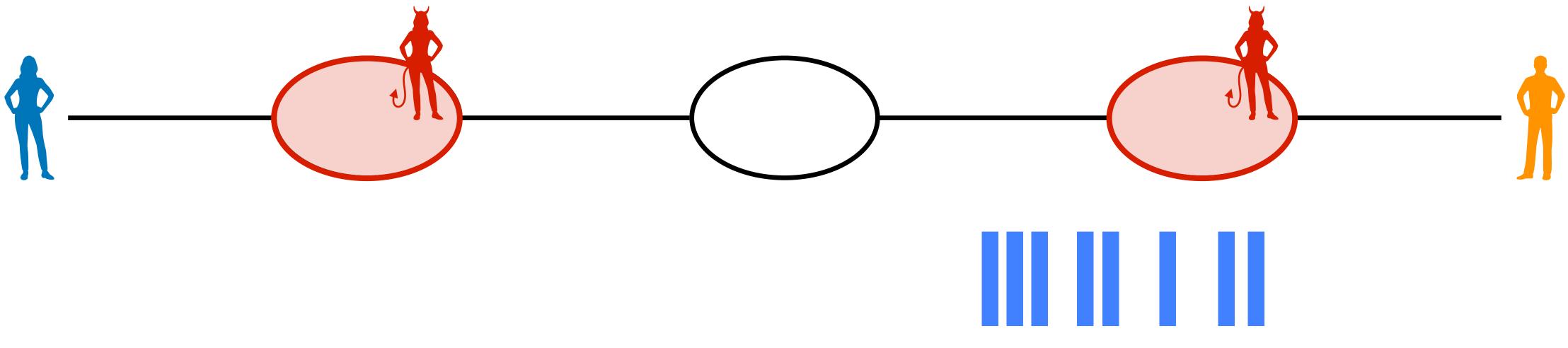






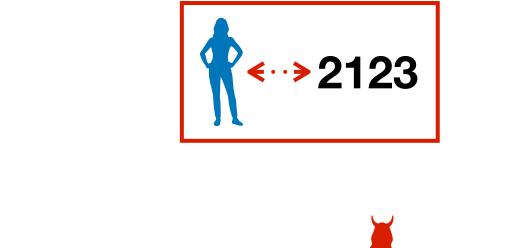


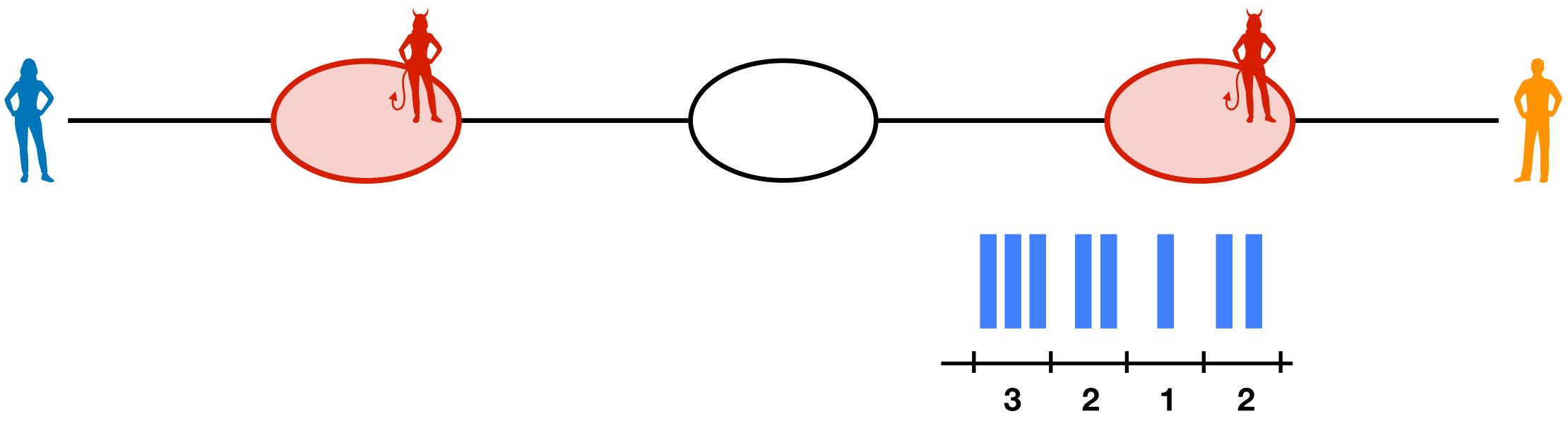






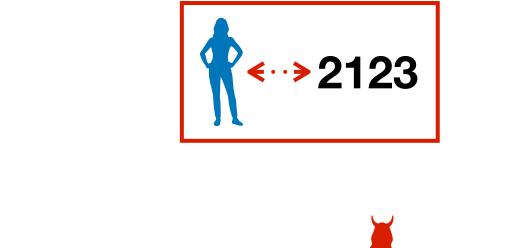


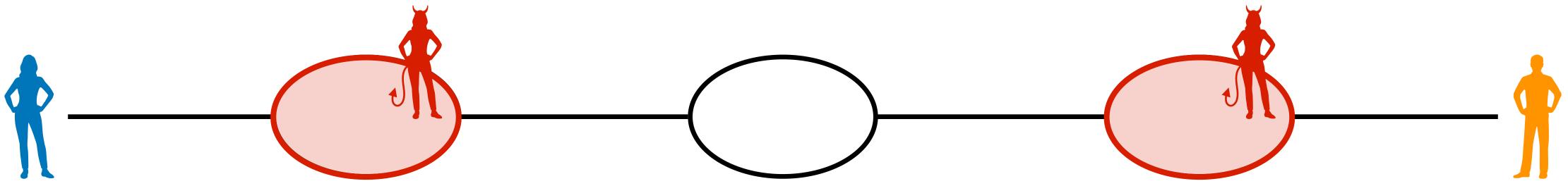








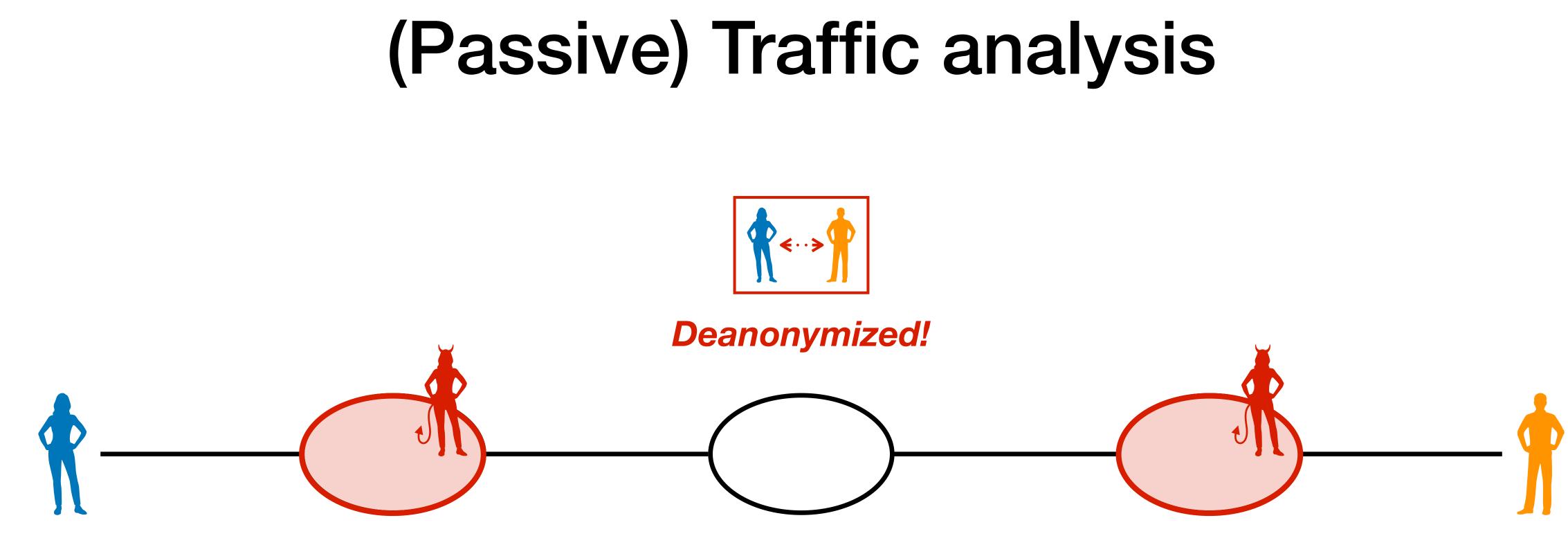






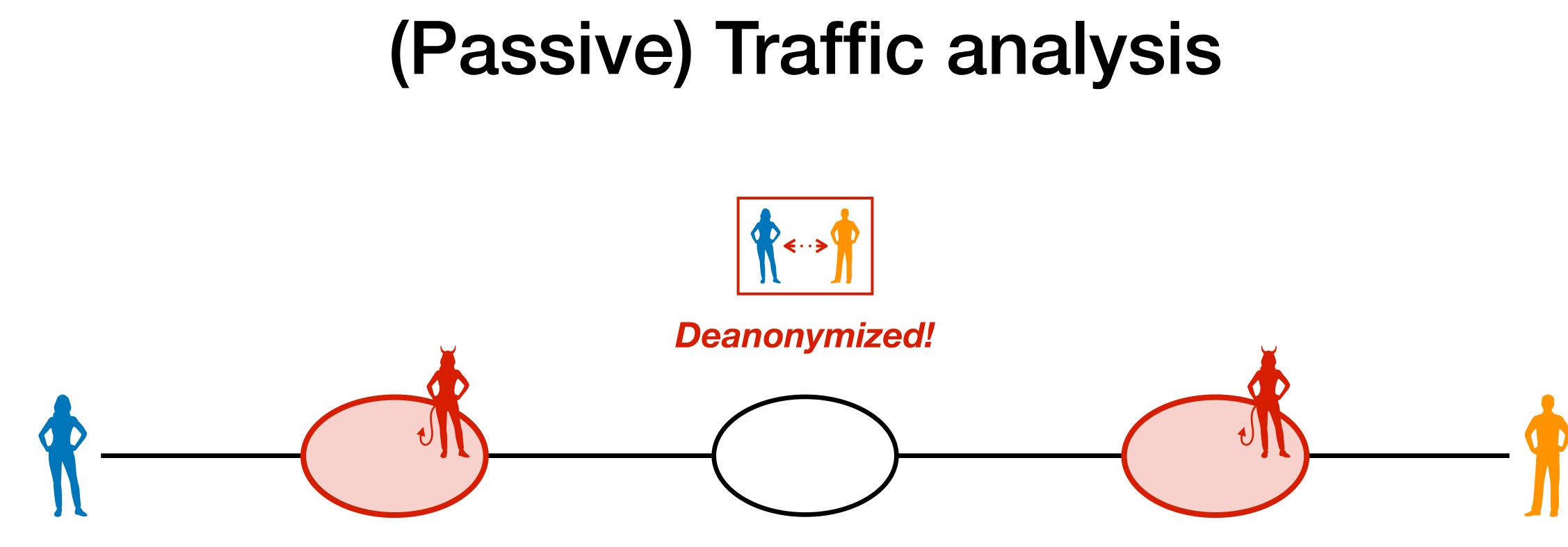








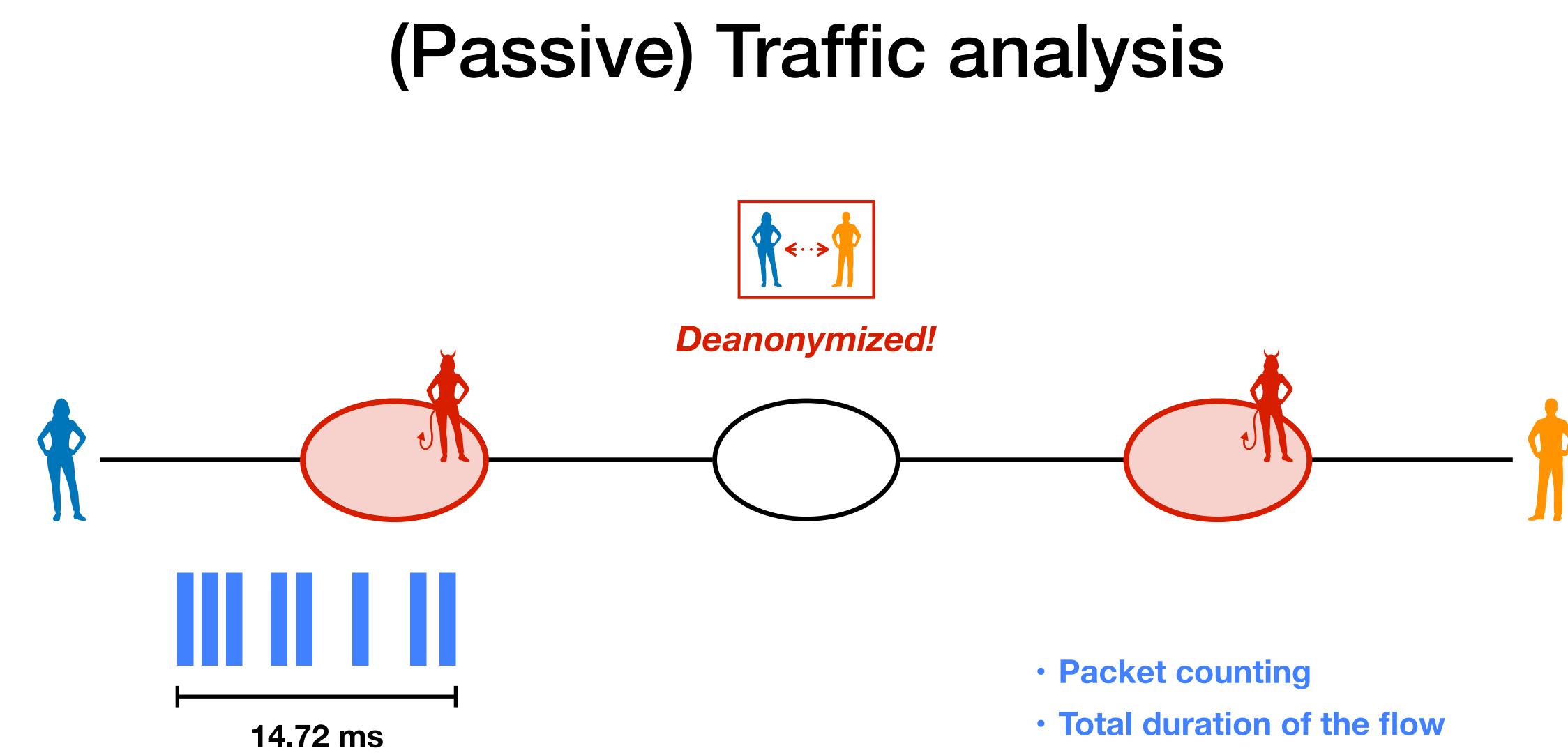






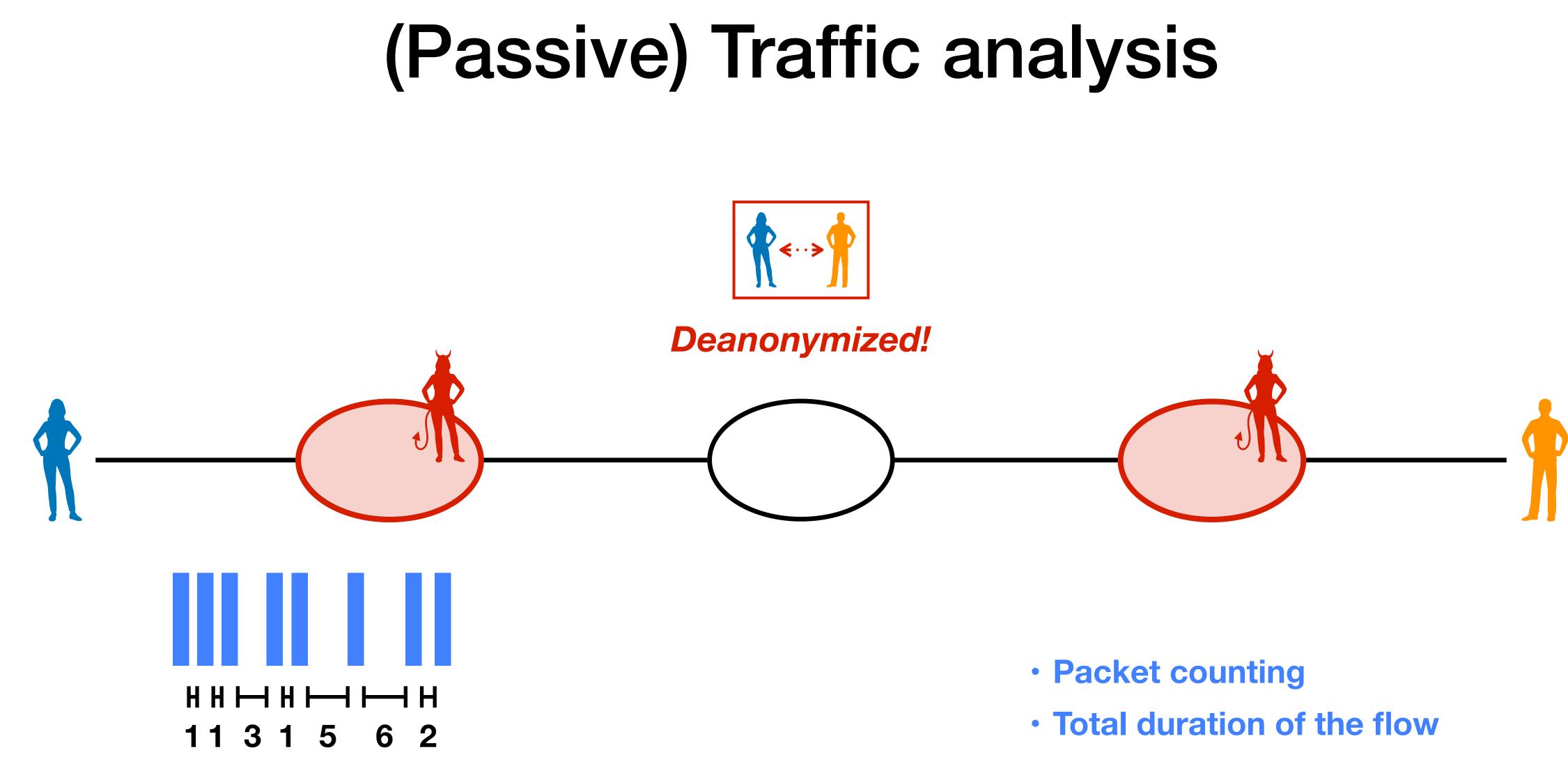
Packet counting







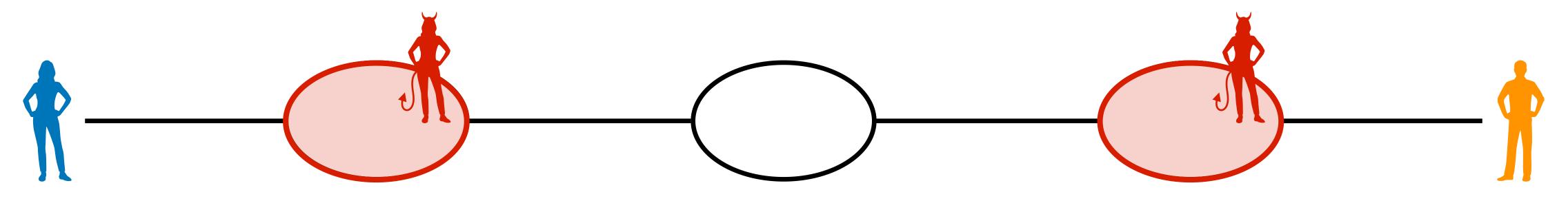




**TH**zürich

- Inter-packet timing

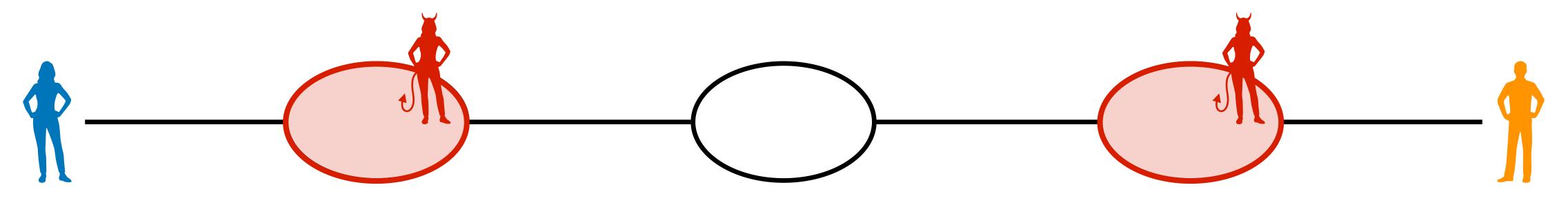








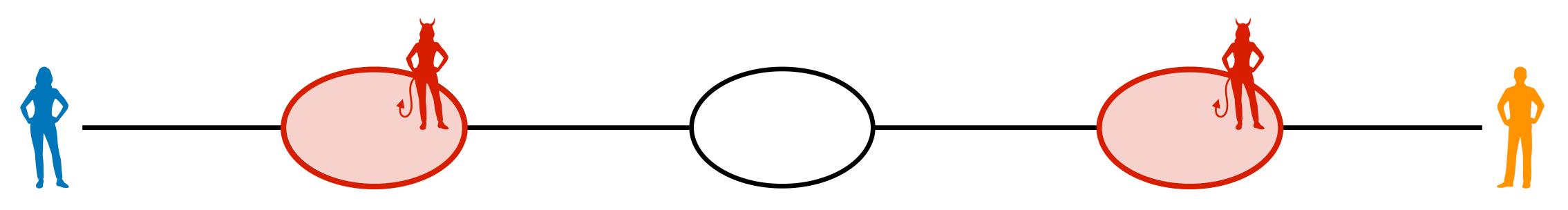








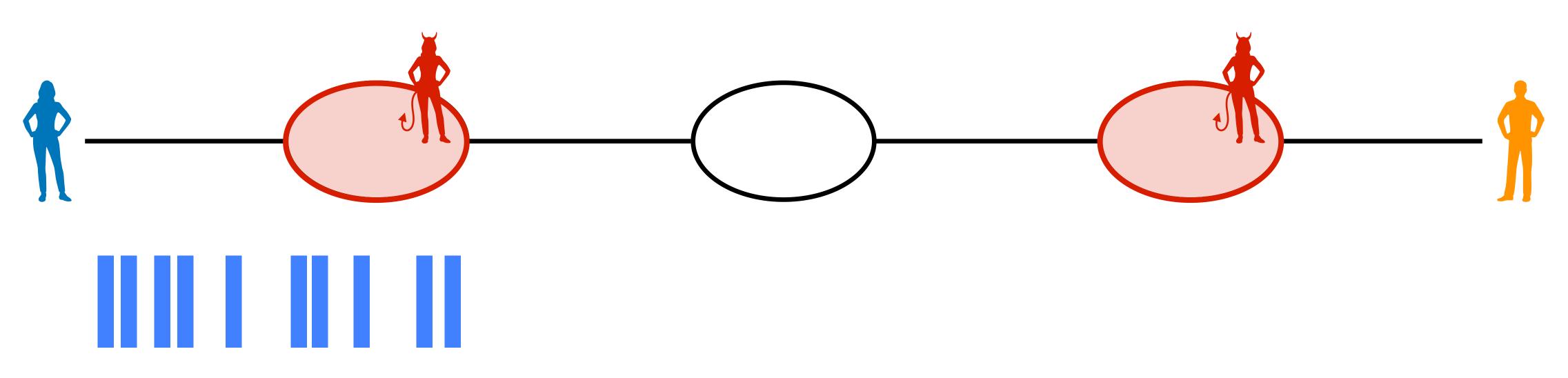








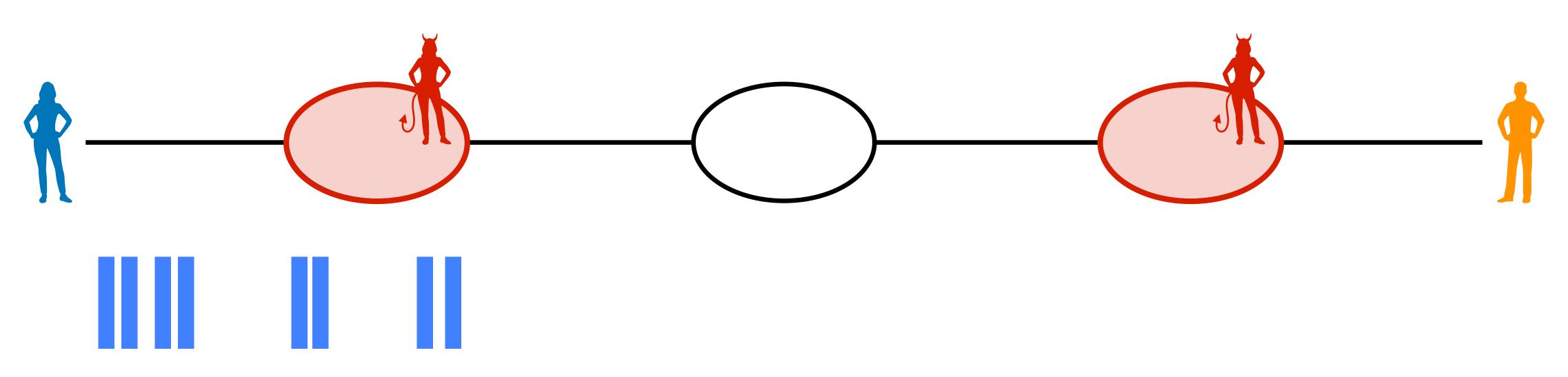








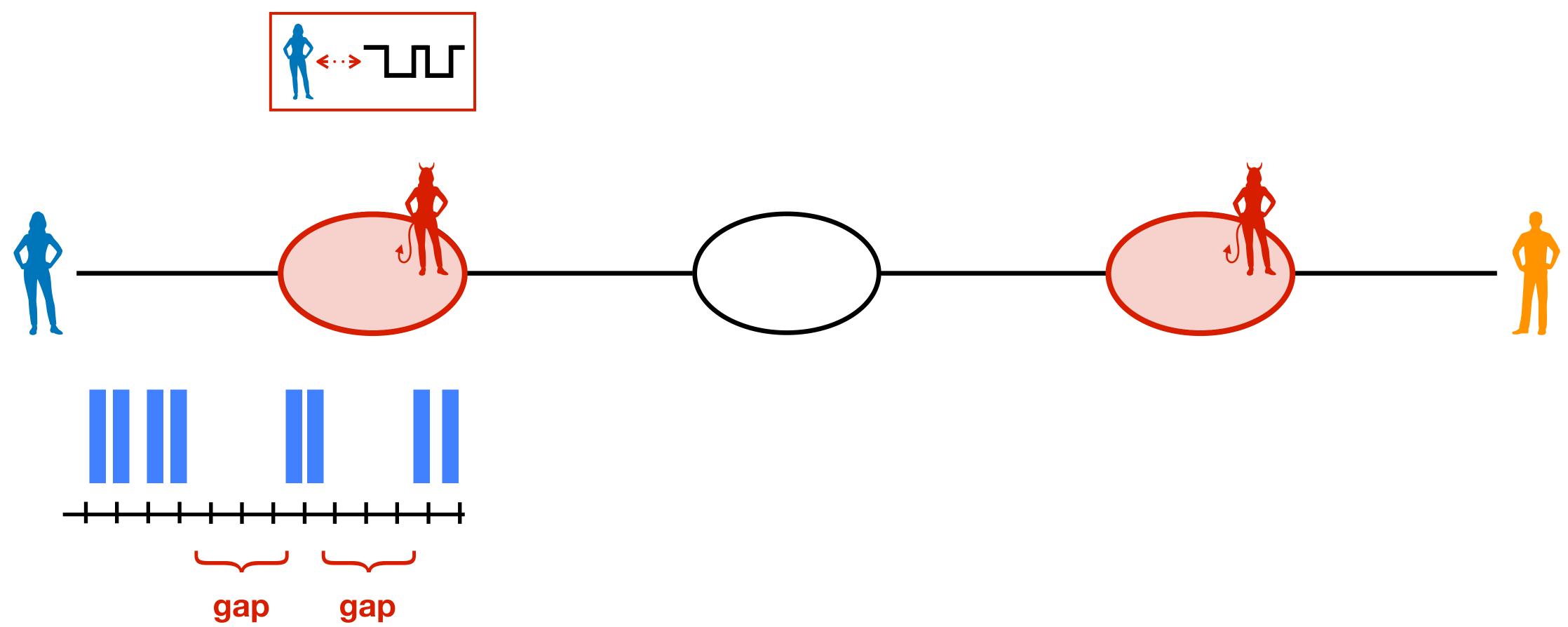








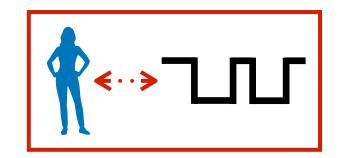


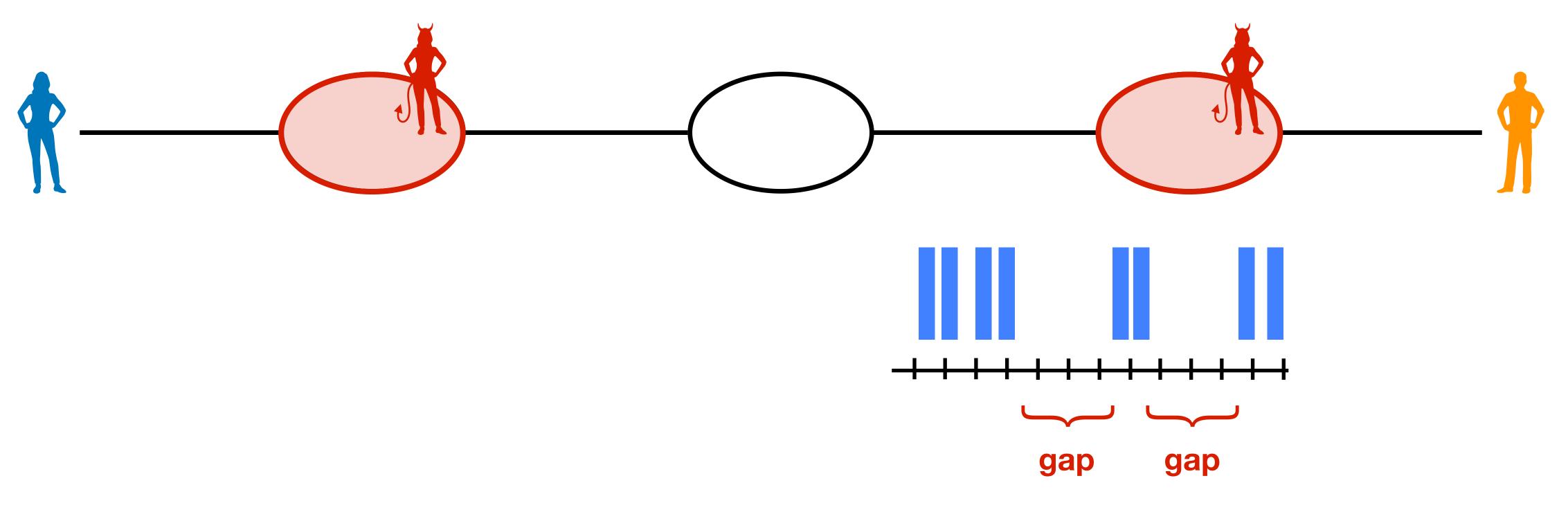








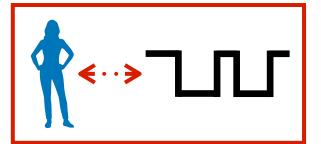


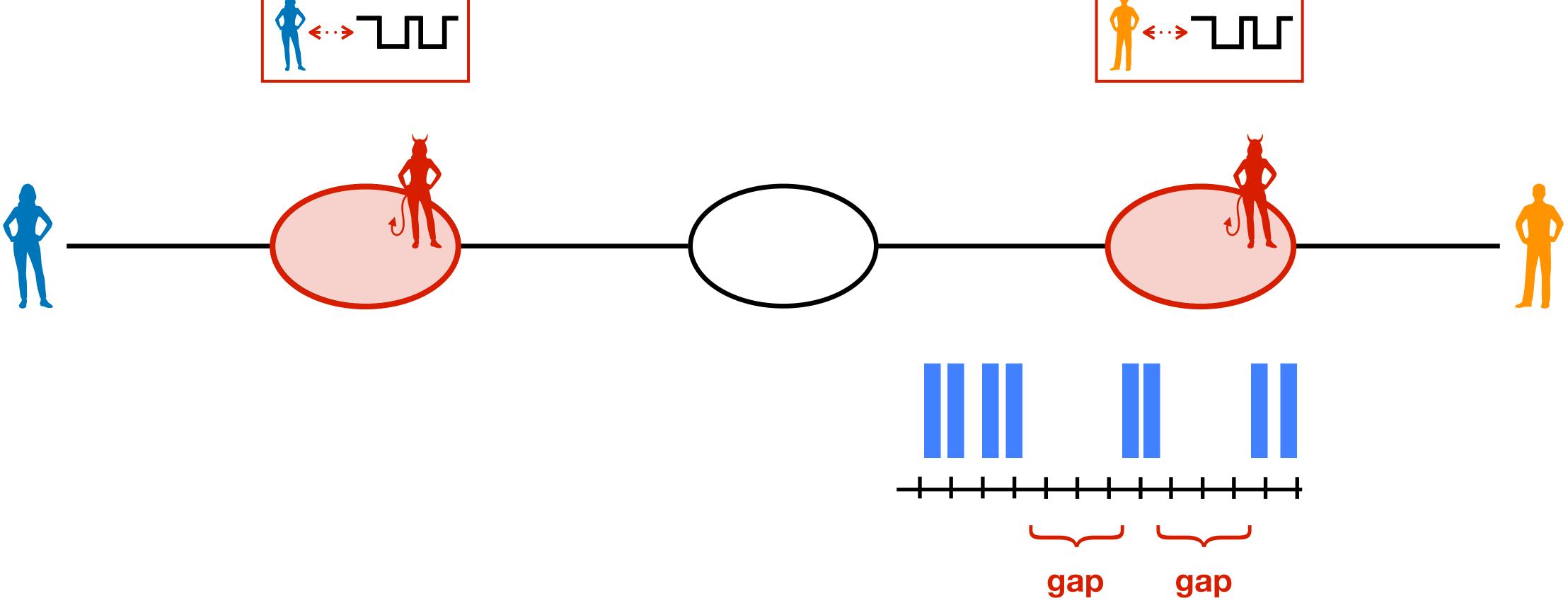




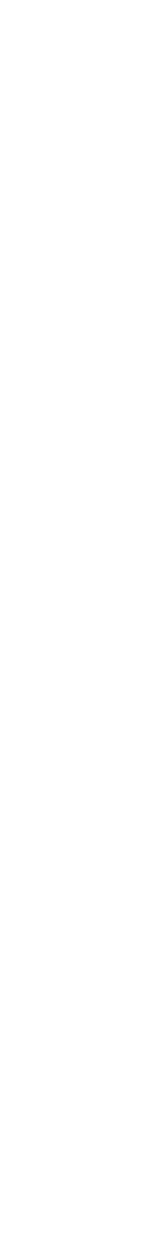




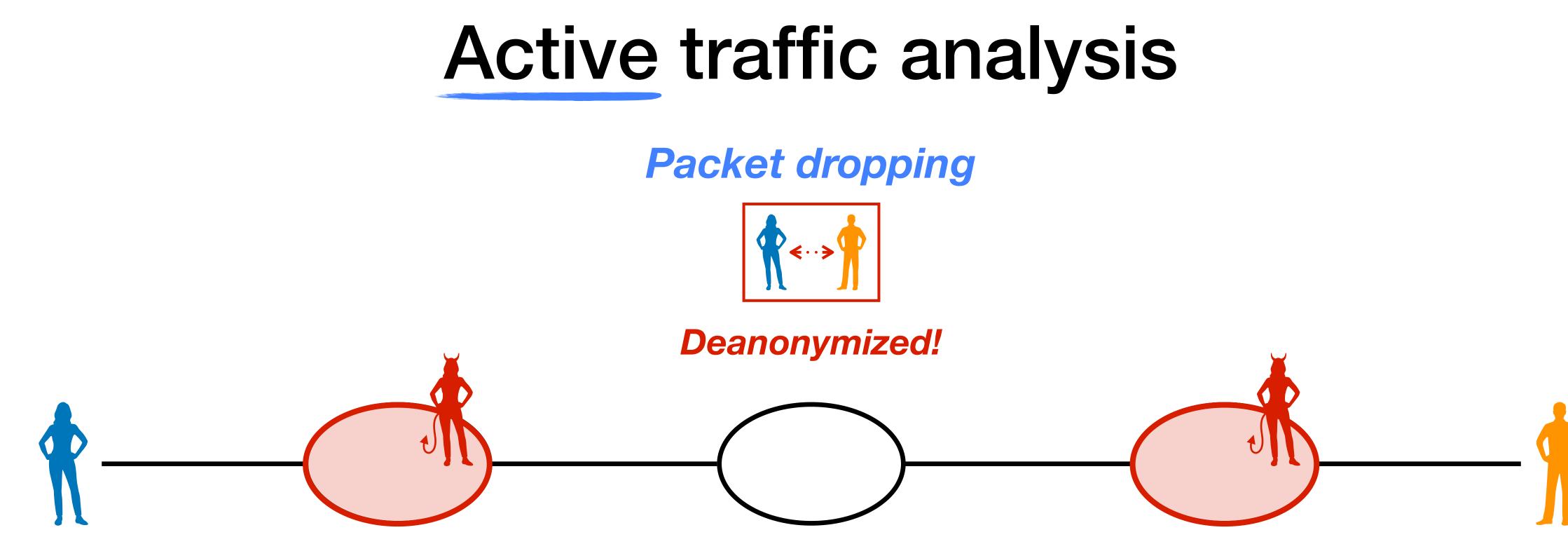


















#### **TARANET** Resisting traffic analysis attacks

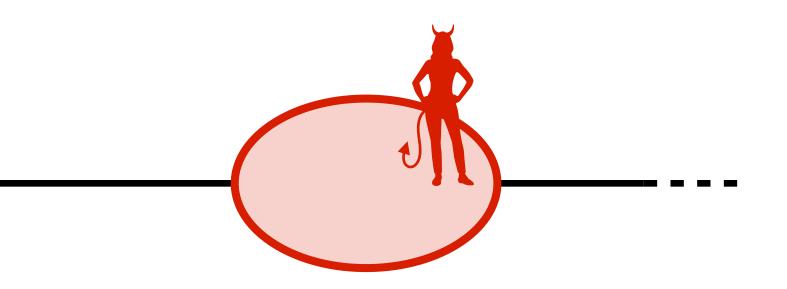








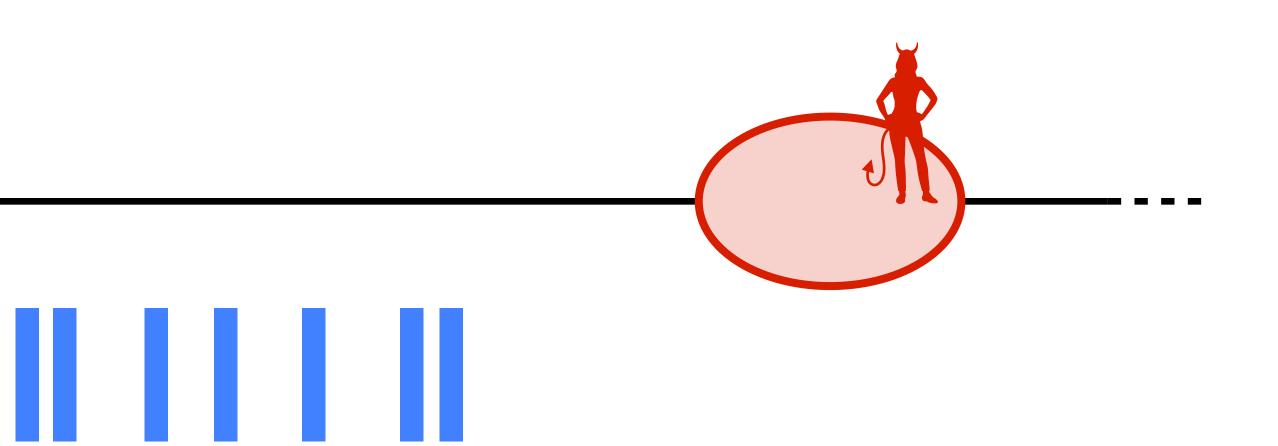










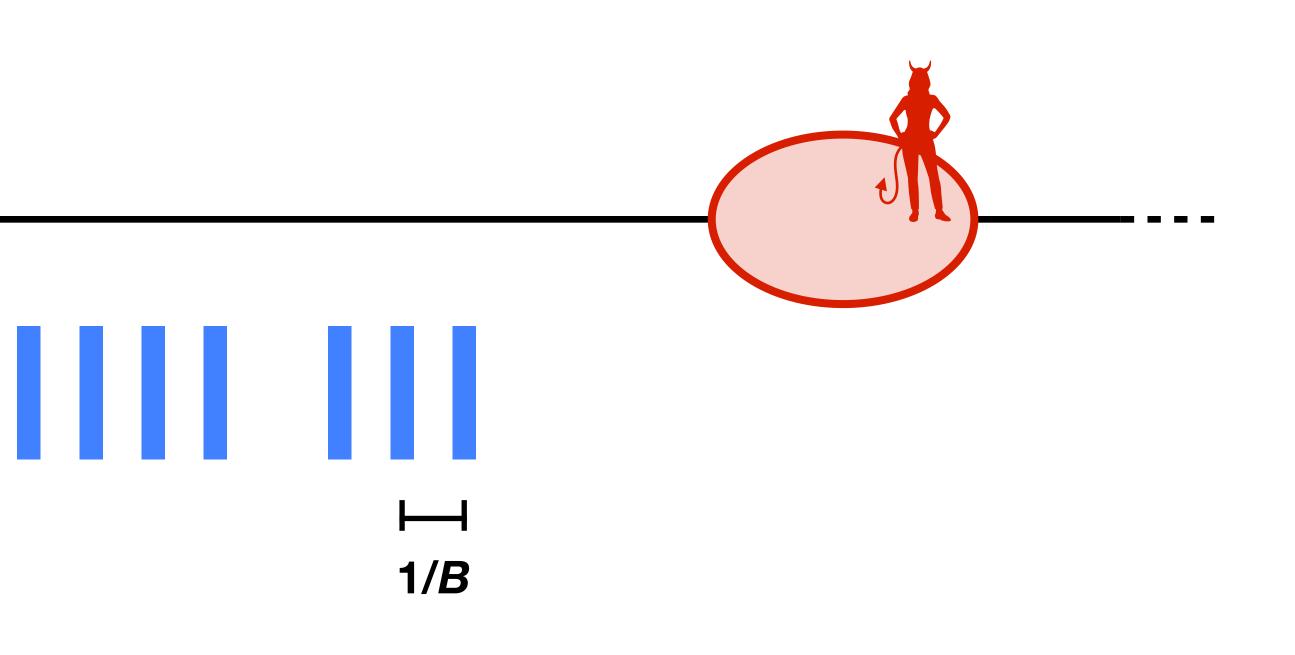




• Fixed rate 1/B





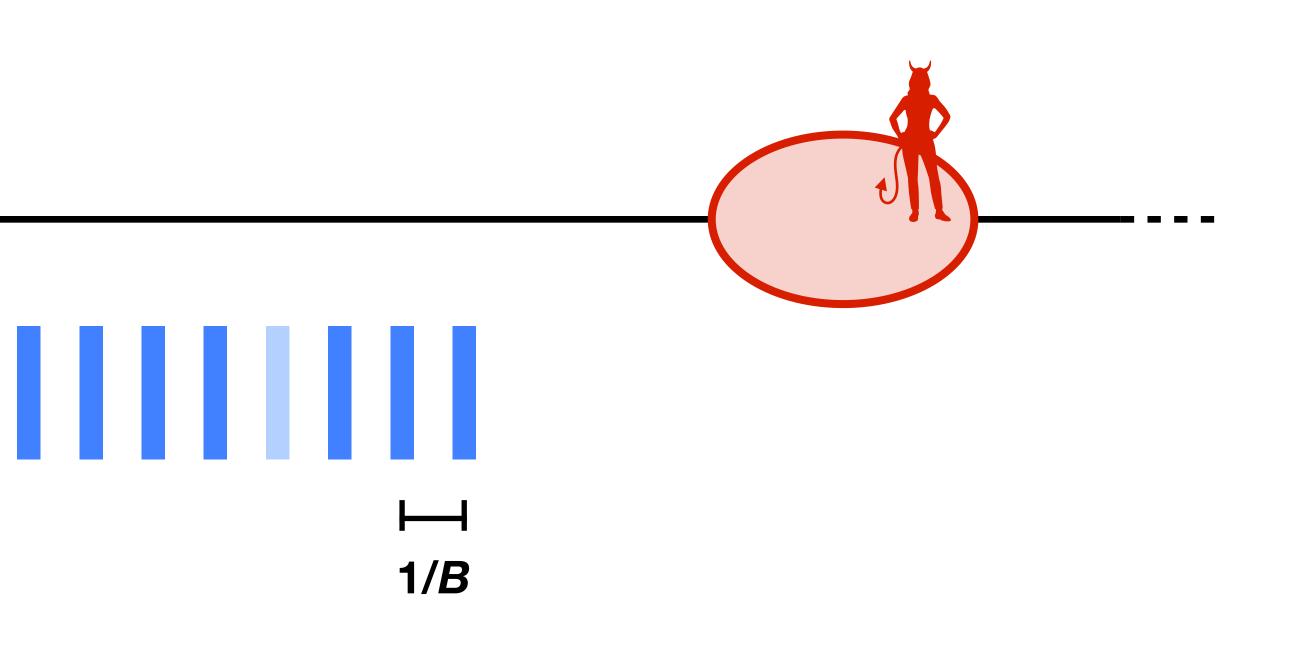




- Fixed rate 1/B
- Chaff packets





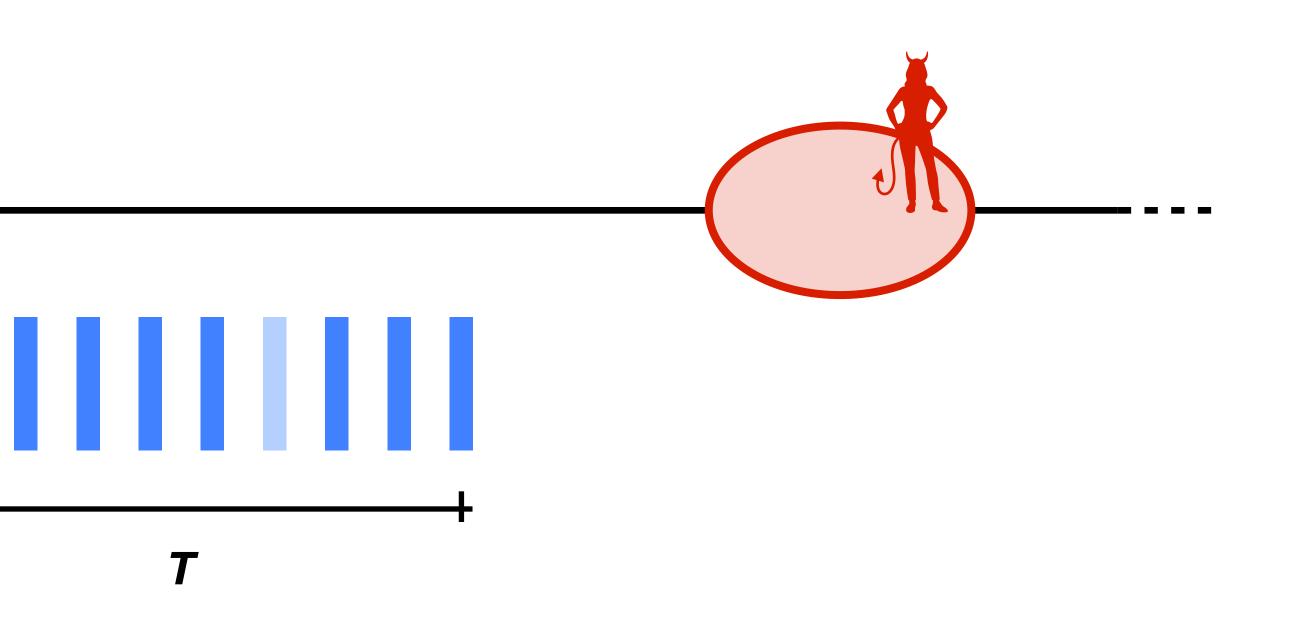




- Fixed rate 1/B
- Chaff packets
- Fixed duration T





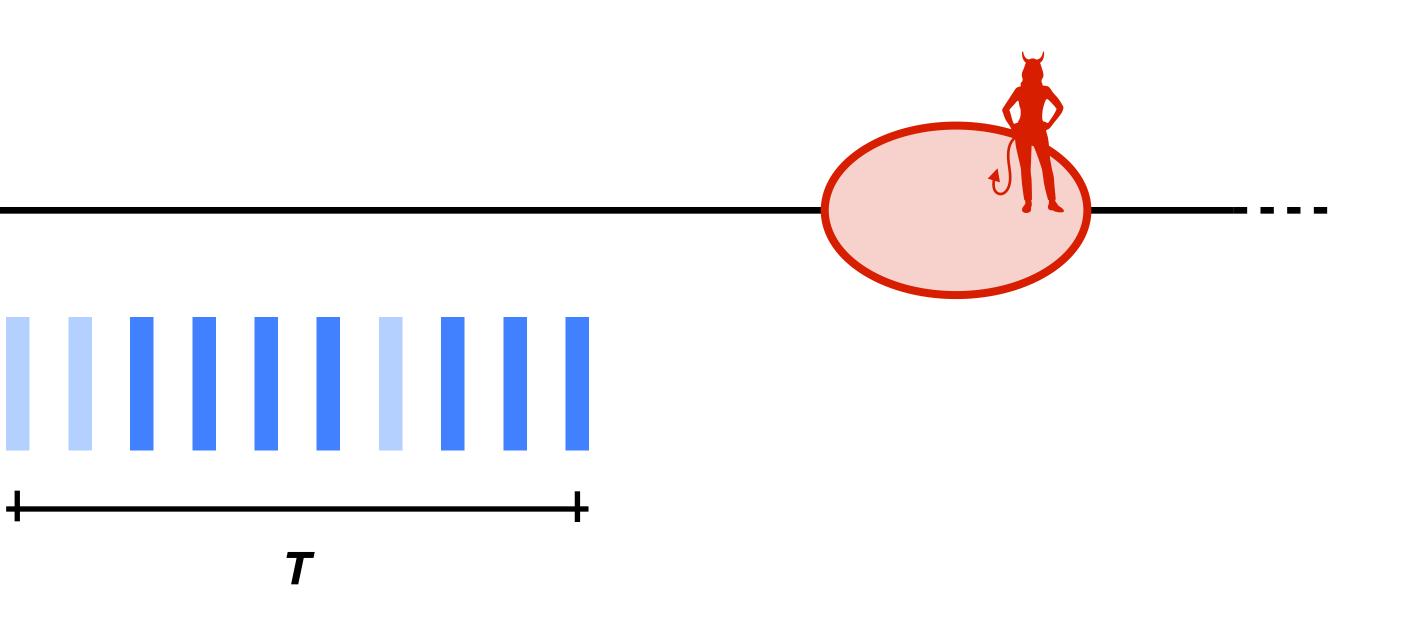




- Fixed rate 1/B
- Chaff packets
- Fixed duration T

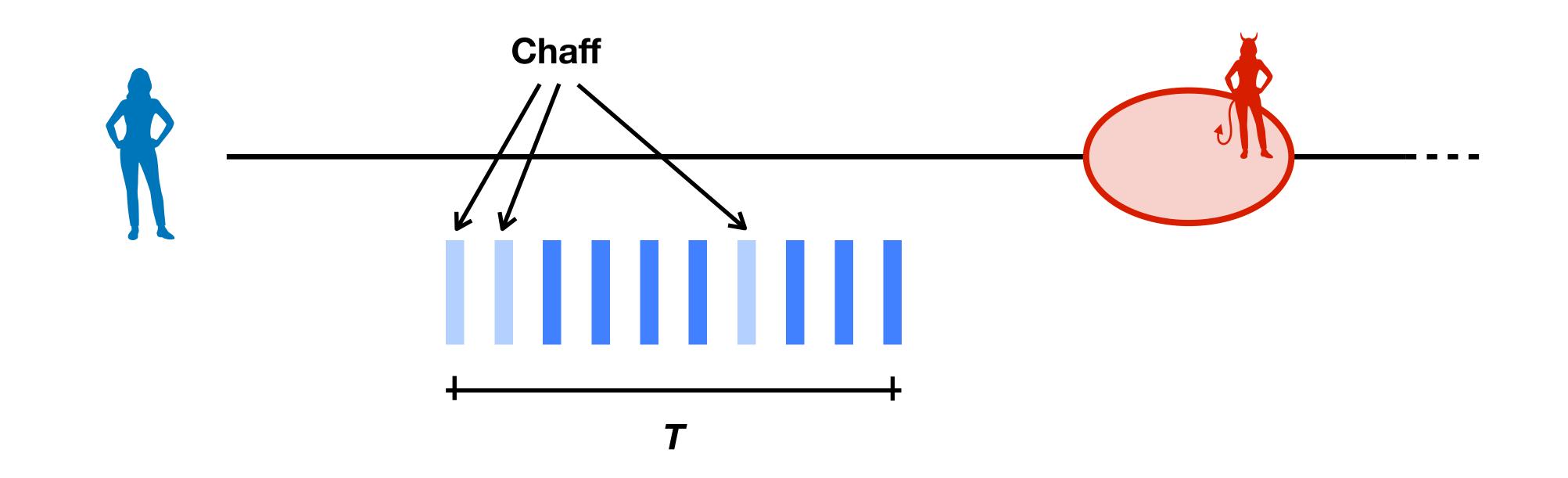








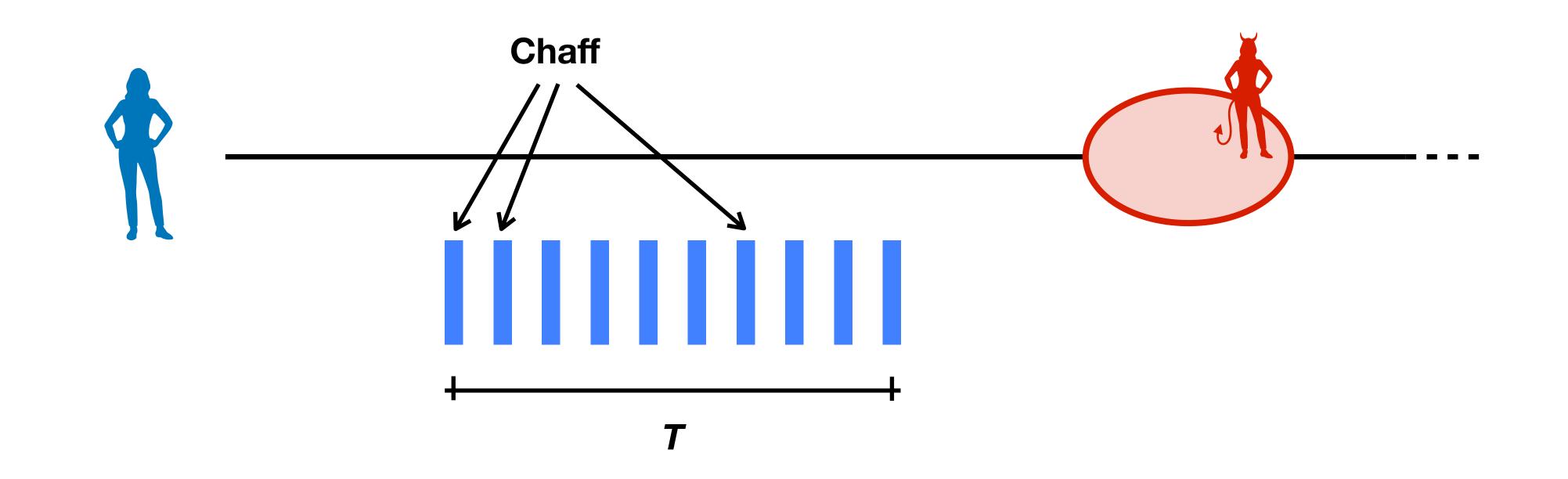
- Fixed rate 1/B
- Chaff packets
- Fixed duration T







- Fixed rate 1/B
- Chaff packets
- Fixed duration T



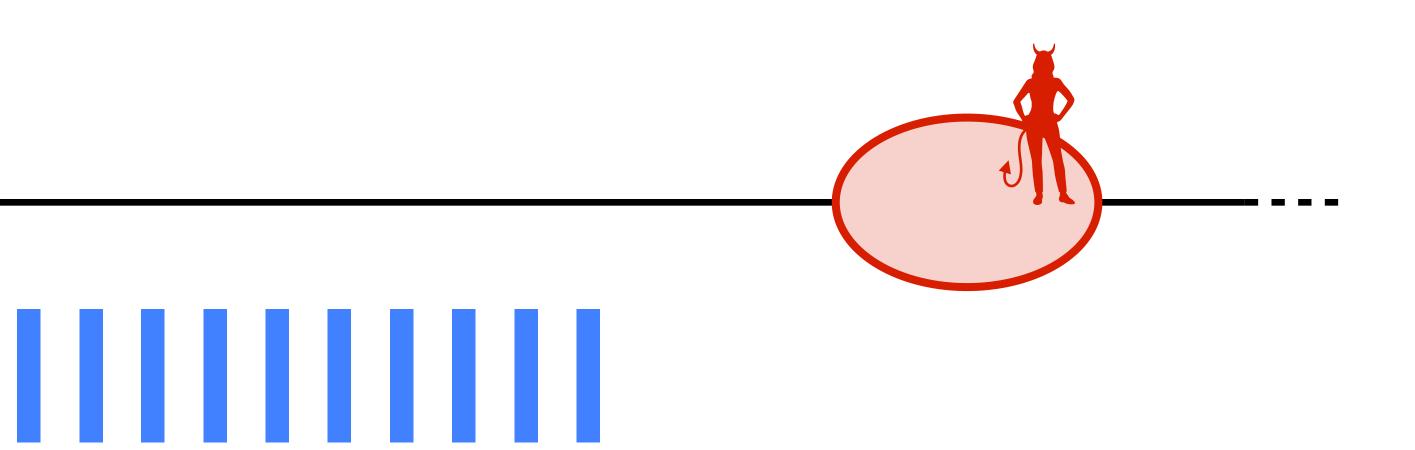




- Fixed rate 1/B
- Chaff packets
- Fixed duration T
- Same for everyone





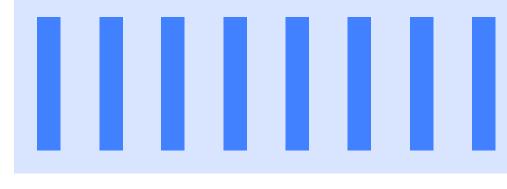




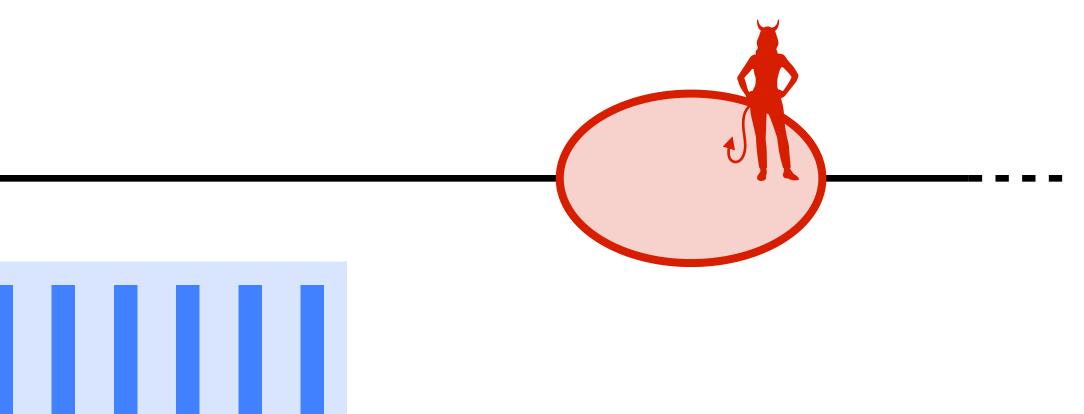
- Fixed rate 1/B
- Chaff packets
- Fixed duration T
- Same for everyone











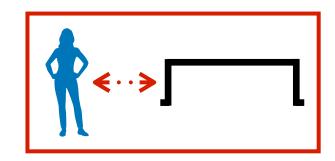


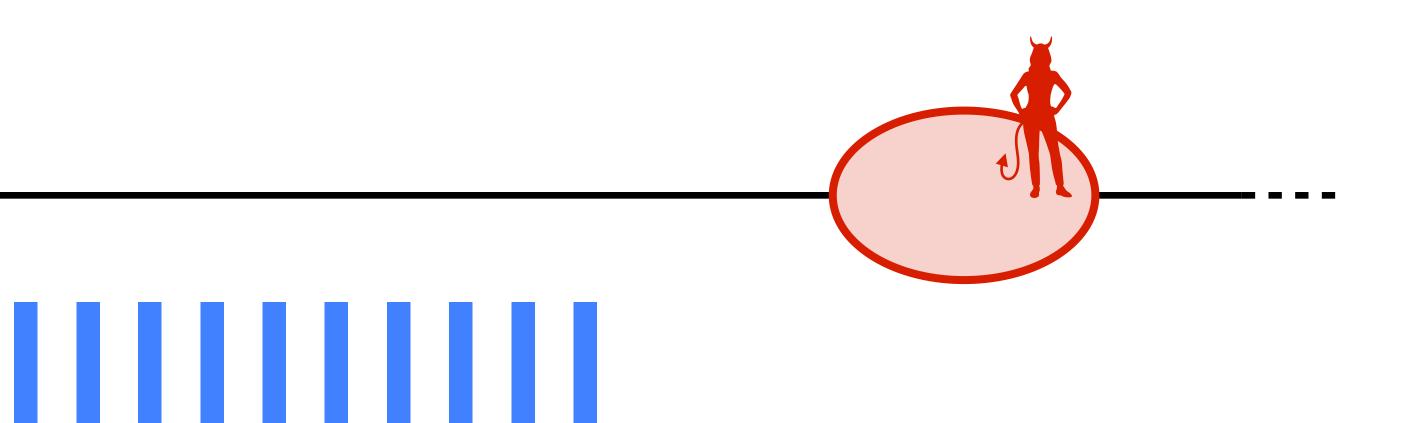
- Fixed rate 1/B
- Chaff packets
- Fixed duration T
- Same for everyone



**Flowlet** 







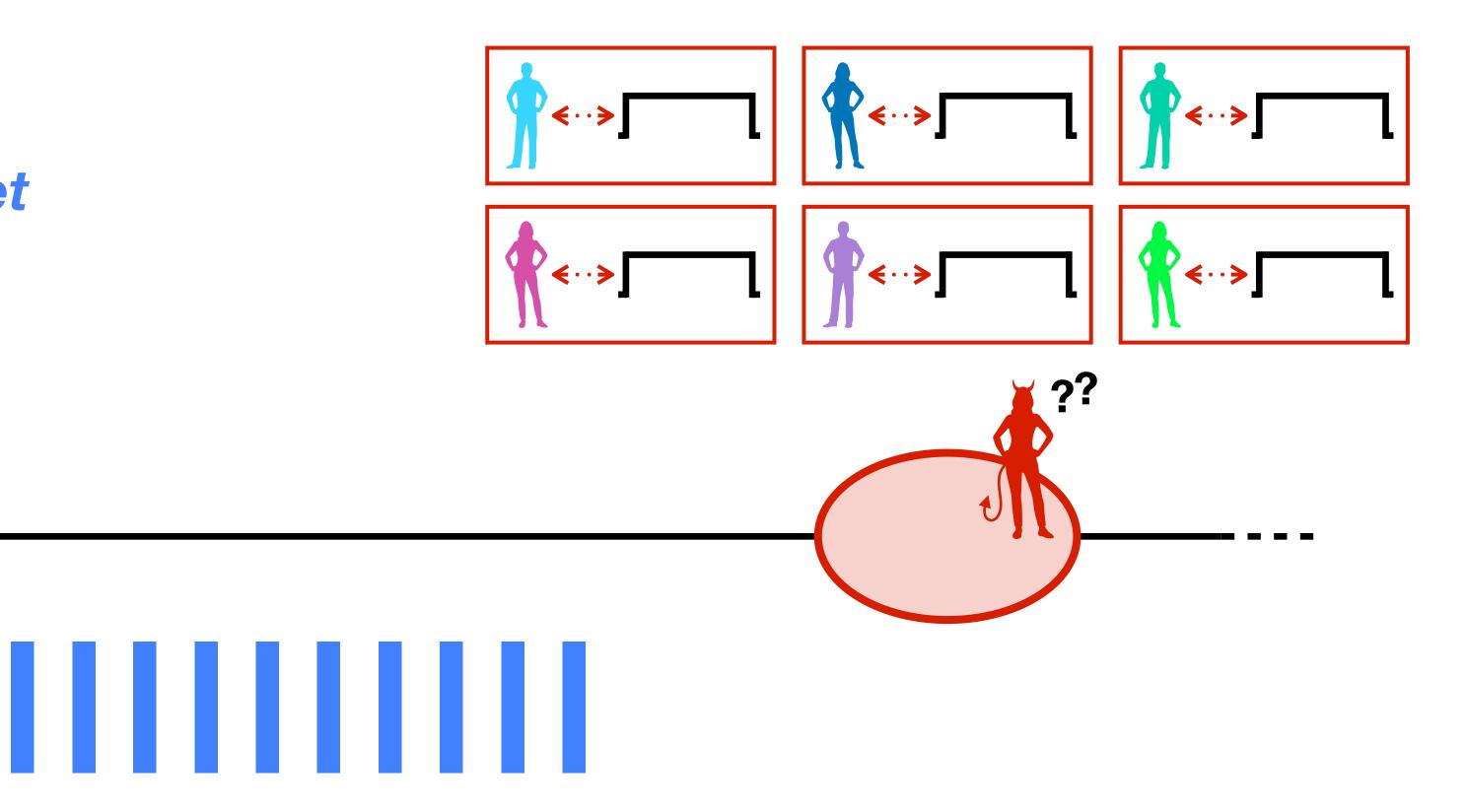


- Fixed rate 1/B
- Chaff packets
- Fixed duration T
- Same for everyone



**Flowlet** 





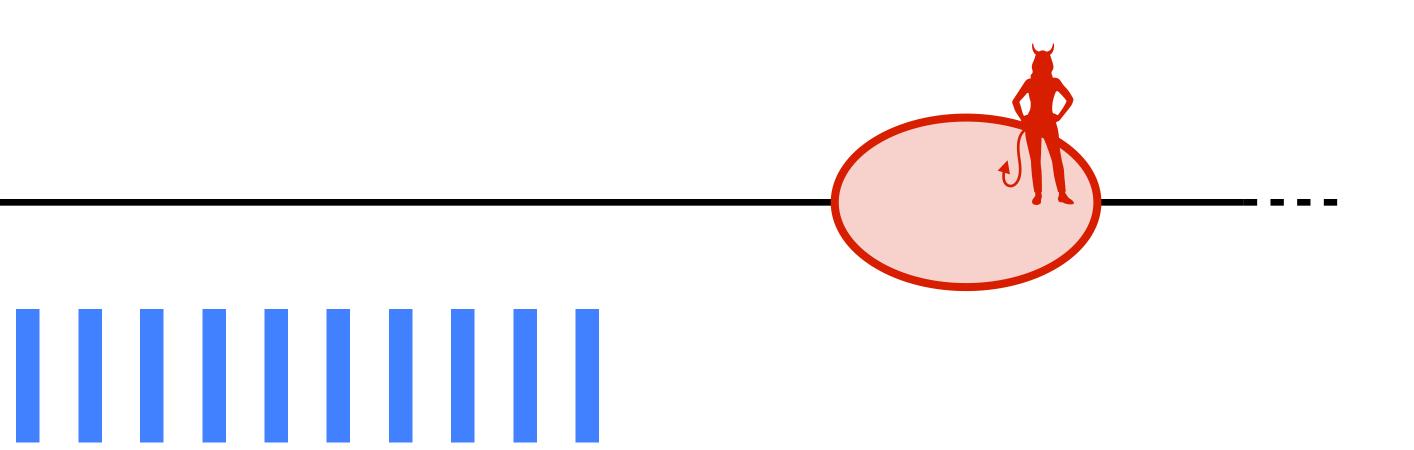


- Fixed rate 1/B
- Chaff packets
- Fixed duration T
- Same for everyone



**Flowlet** 





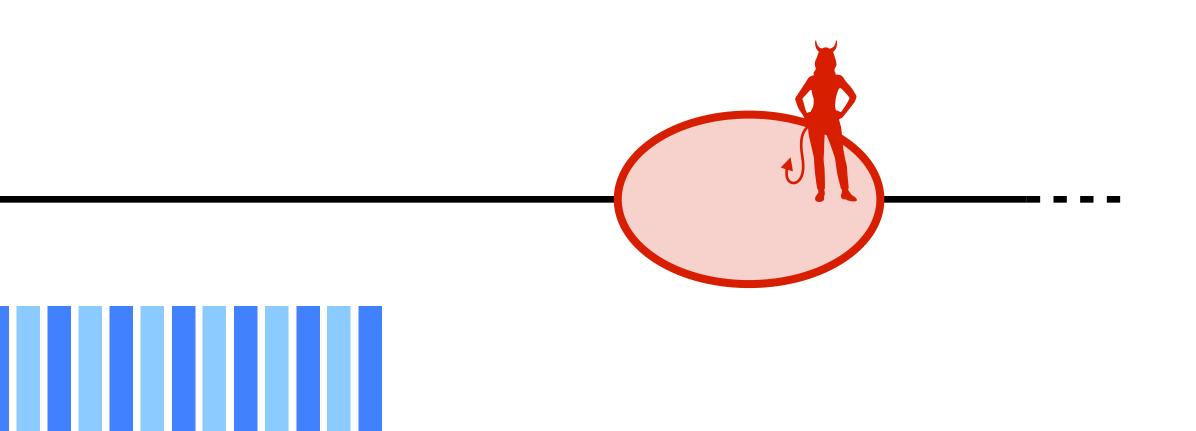


- Fixed rate 1/B
- Chaff packets
- Fixed duration T
- Same for everyone



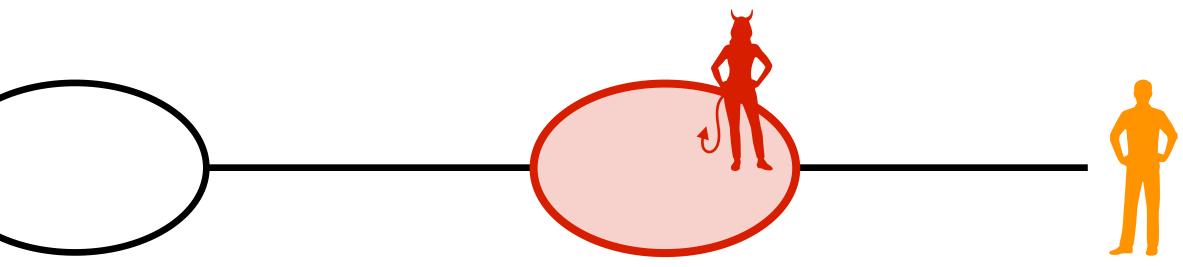






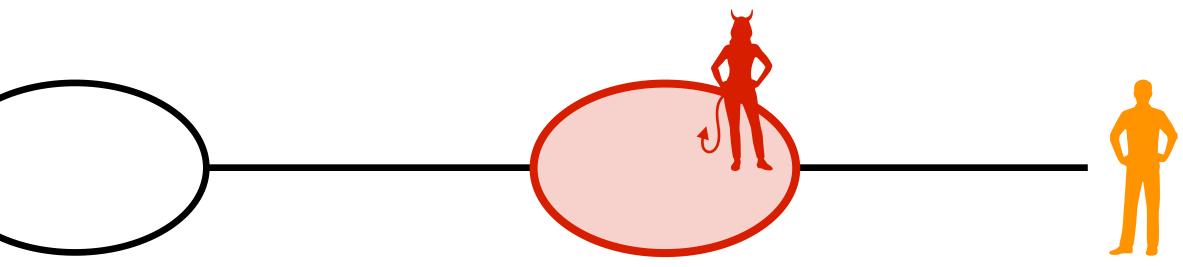




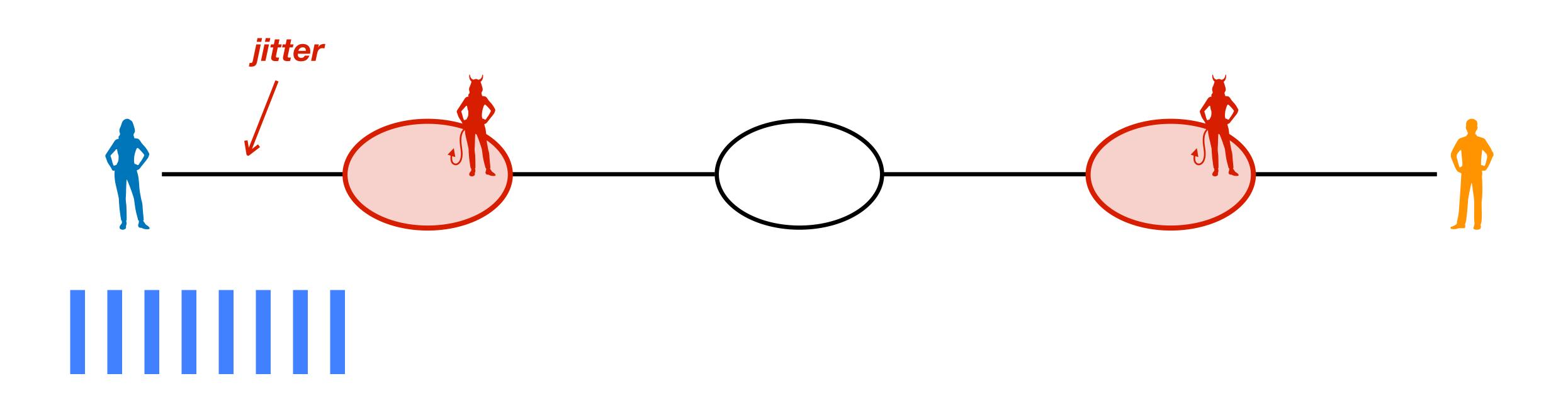






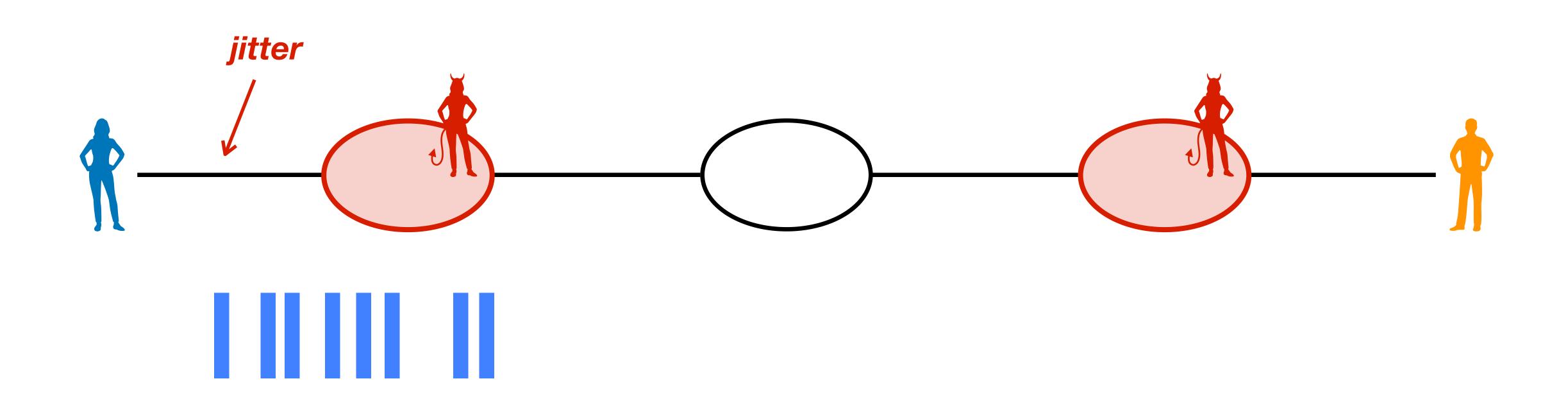










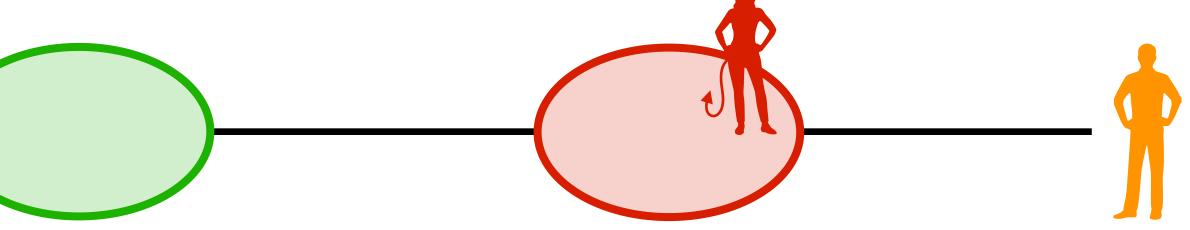






# Resisting passive traffic analysis *re-enforce* flowlet schedule

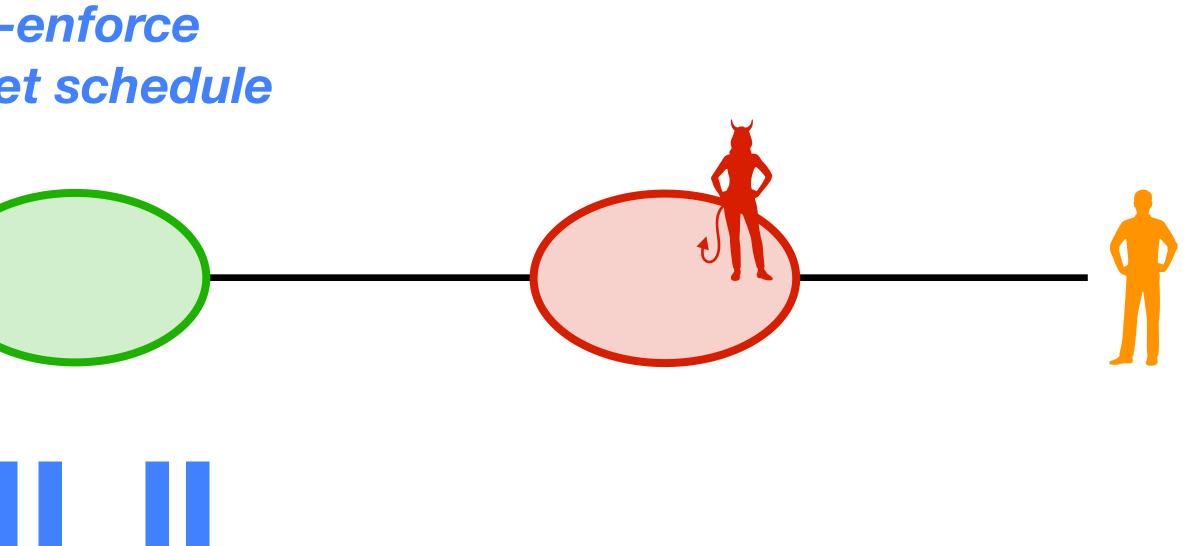






# Resisting passive traffic analysis *re-enforce* flowlet schedule

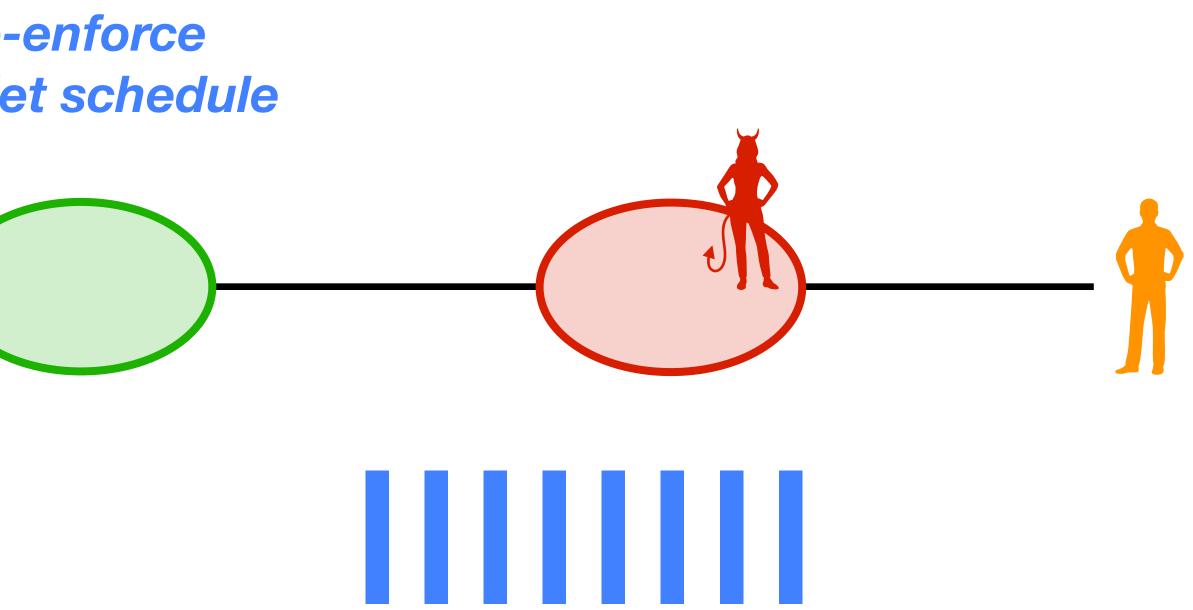






## Resisting passive traffic analysis *re-enforce* flowlet schedule U

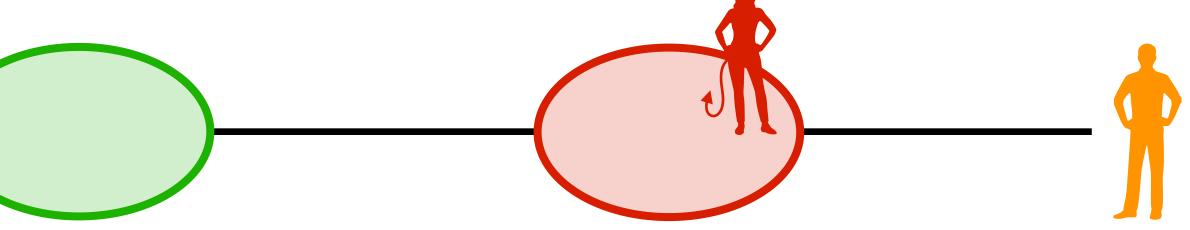




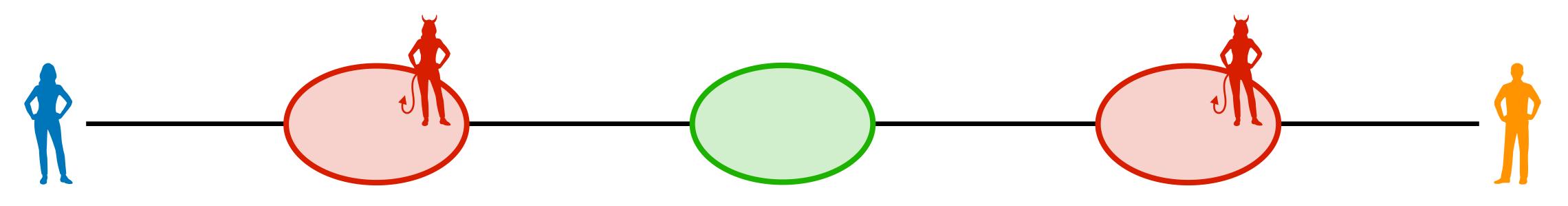


# Resisting passive traffic analysis *re-enforce* flowlet schedule





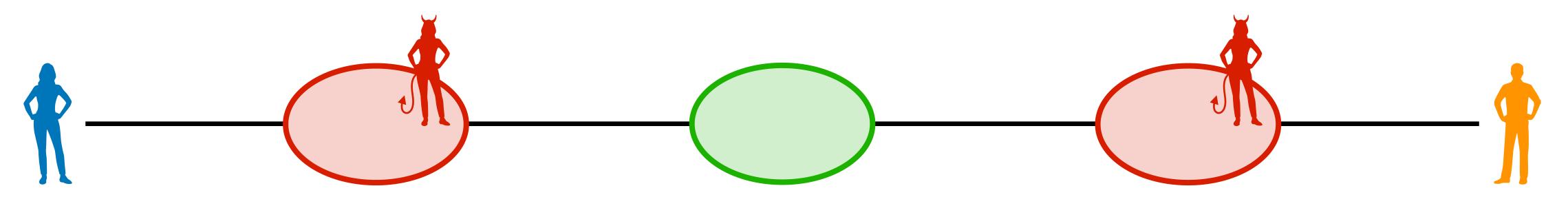












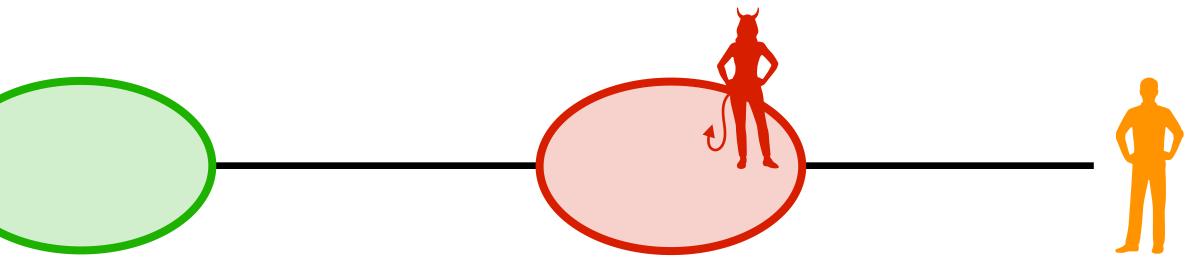






# Resisting active traffic analysis **Packet dropping** U



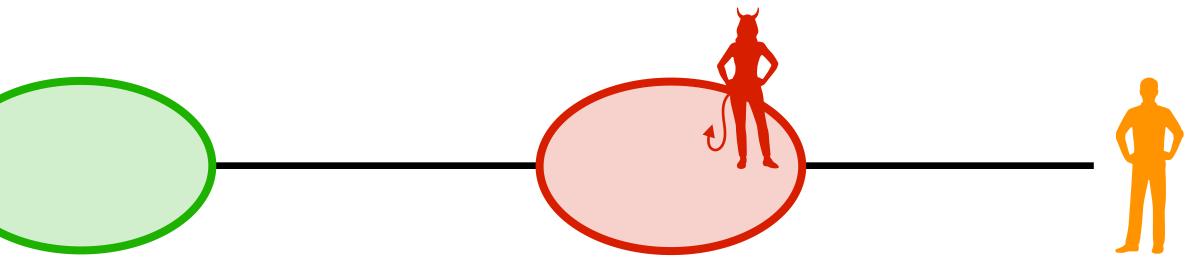






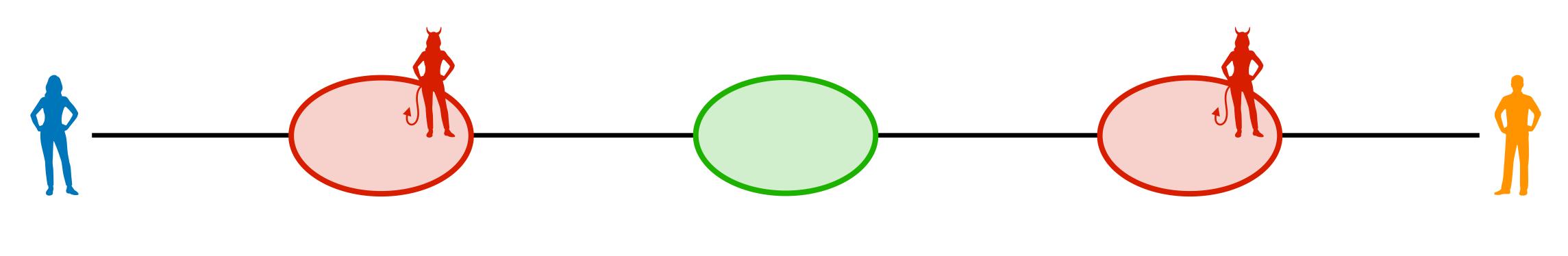
# Resisting active traffic analysis **Packet dropping** U







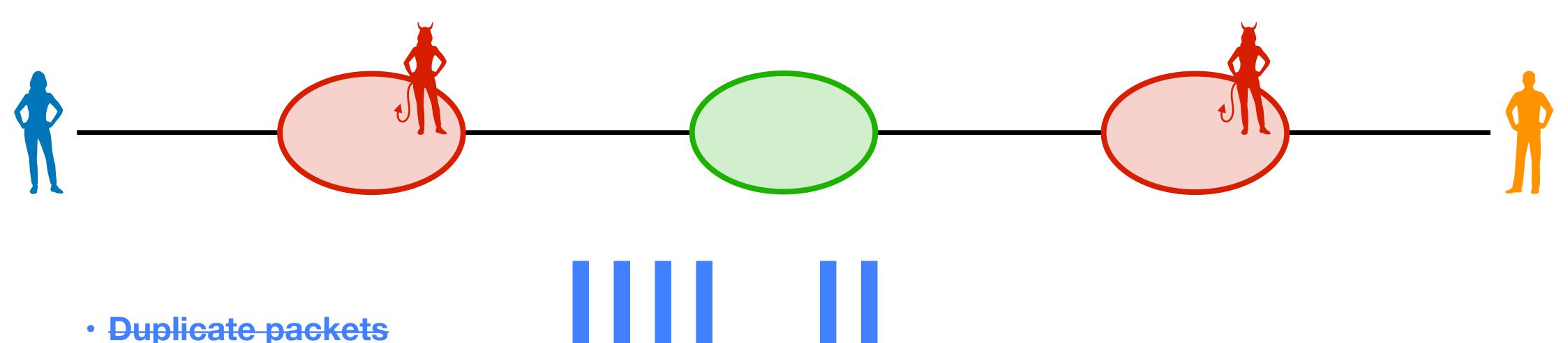










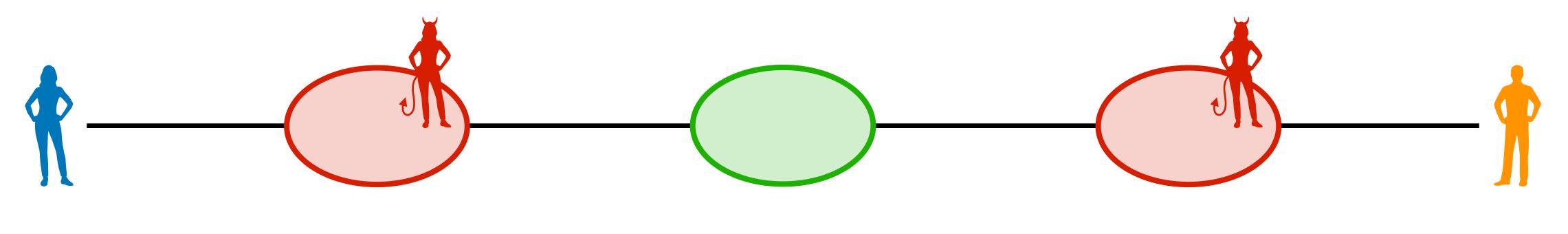


Duplicate packets



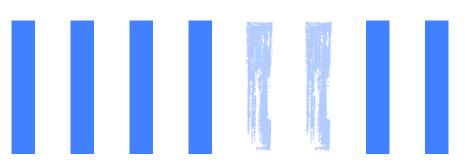






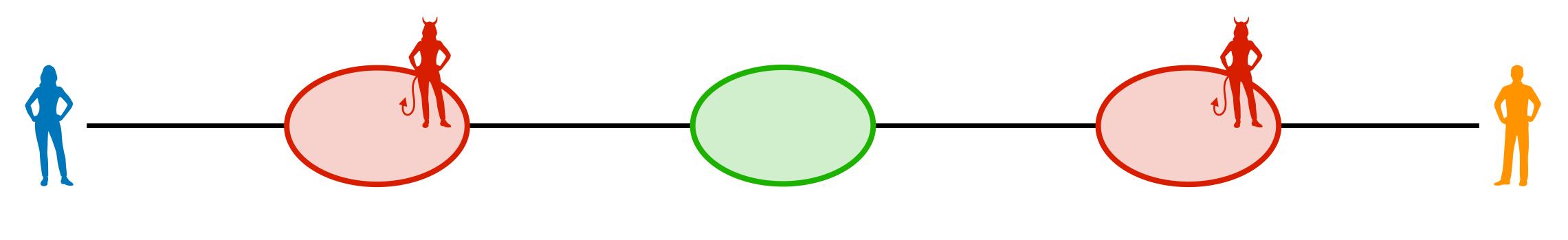
- Duplicate packets
- Create new packets











- Duplicate packets
- Create new packets
- Sender send more?

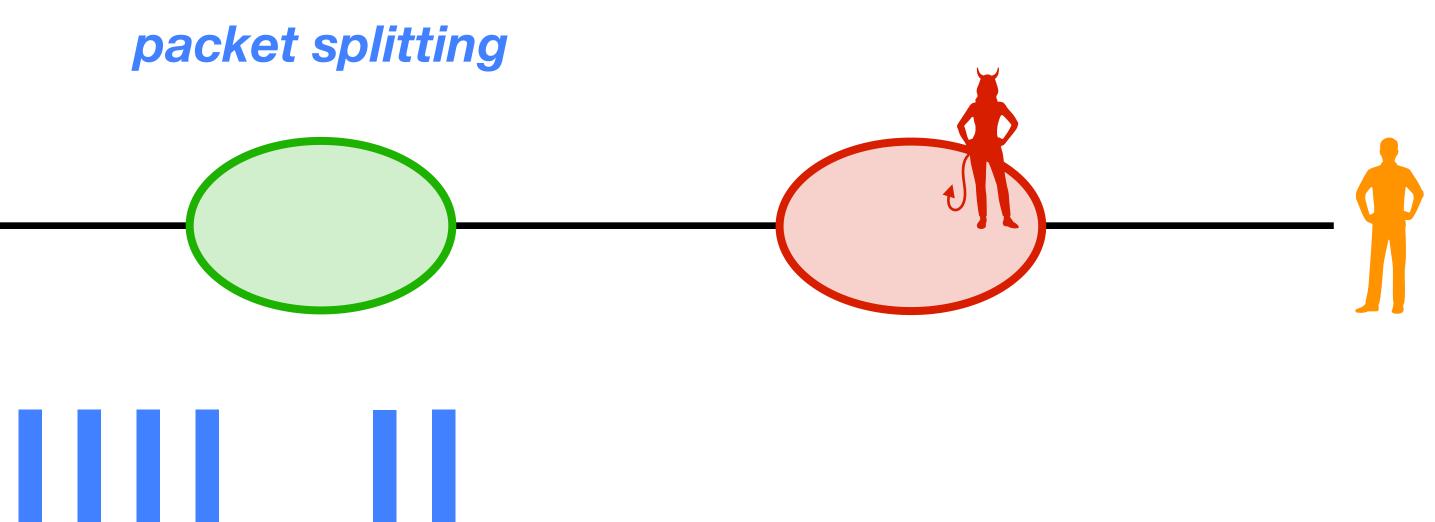










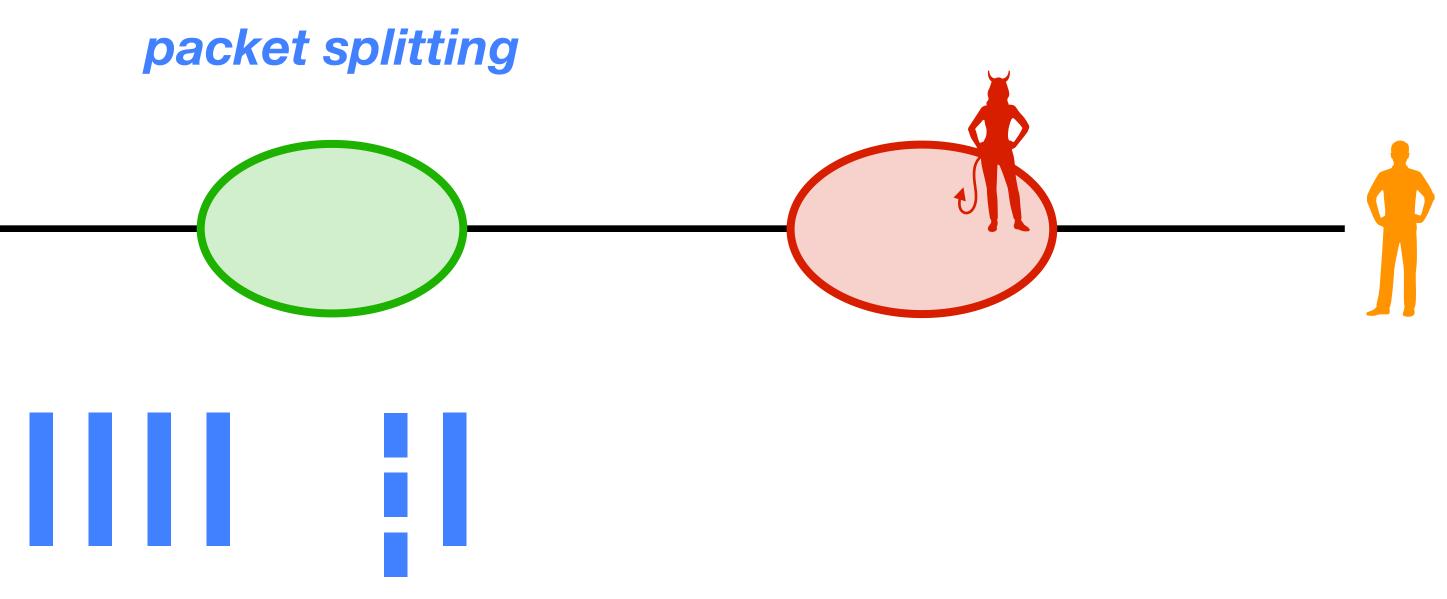










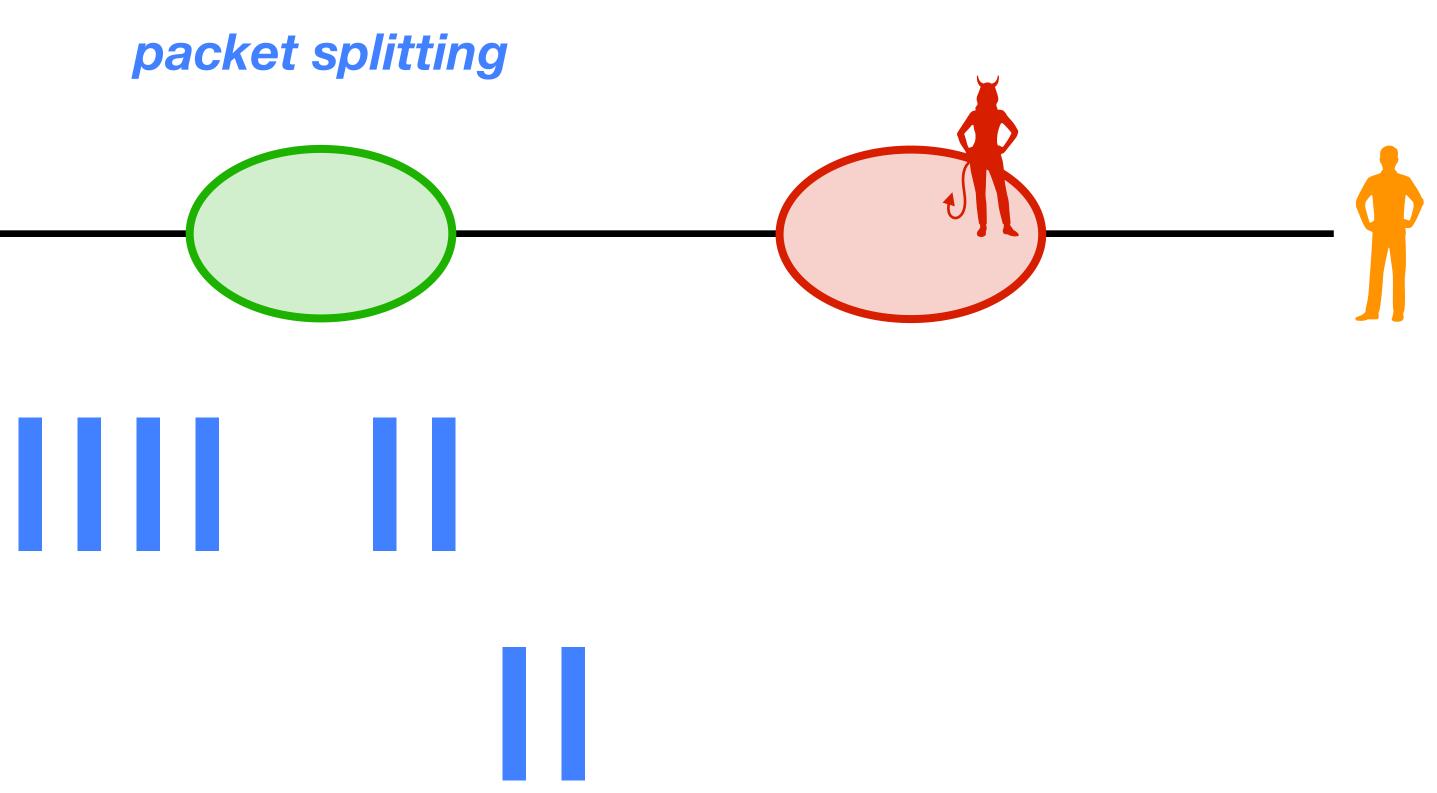










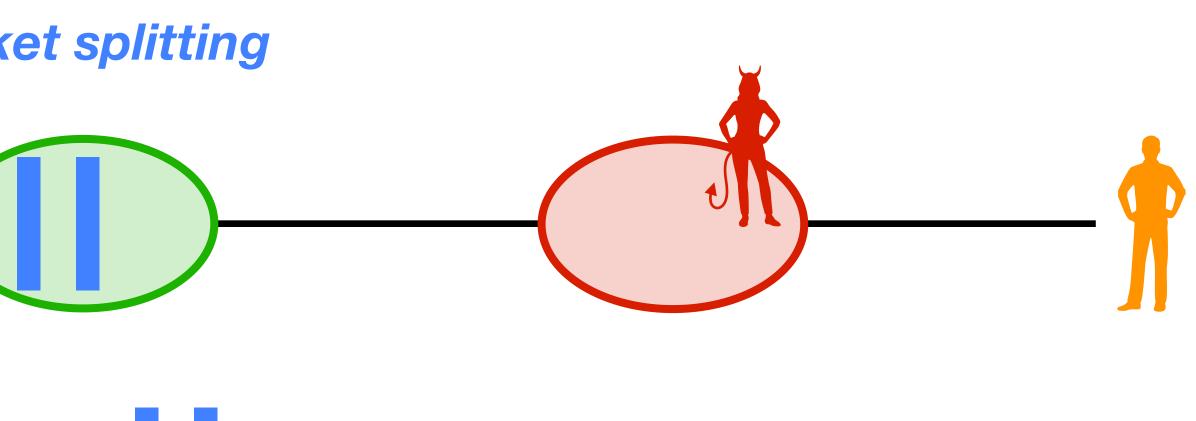










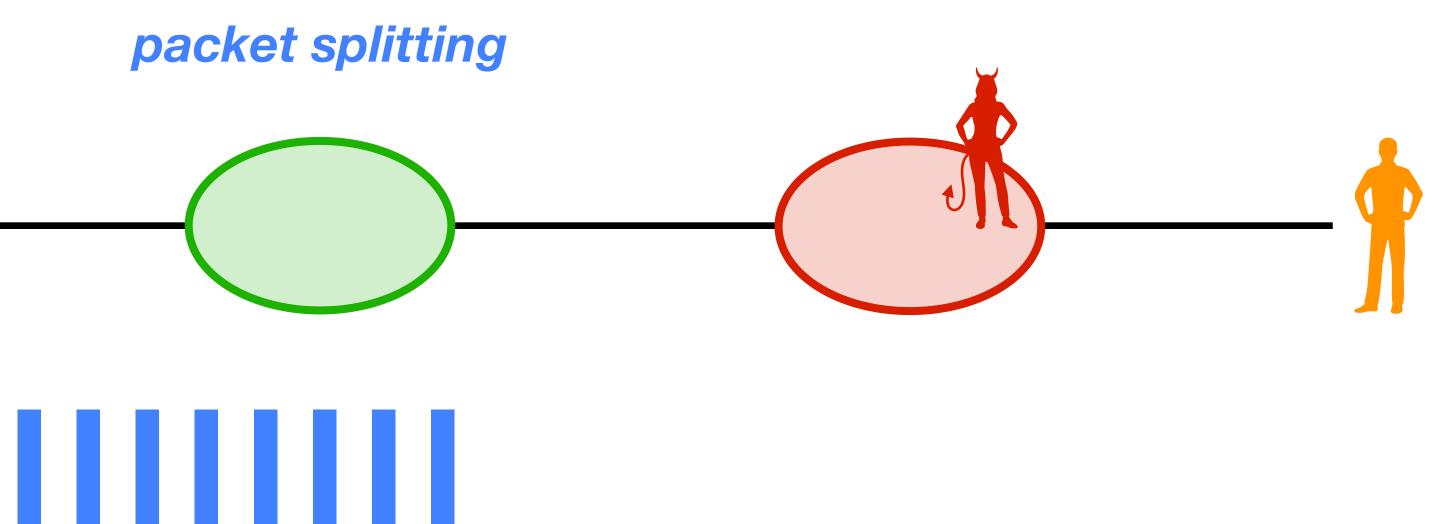








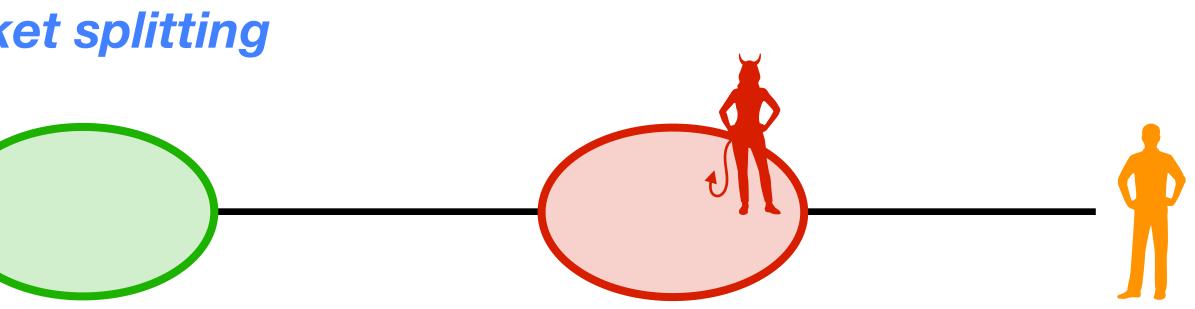








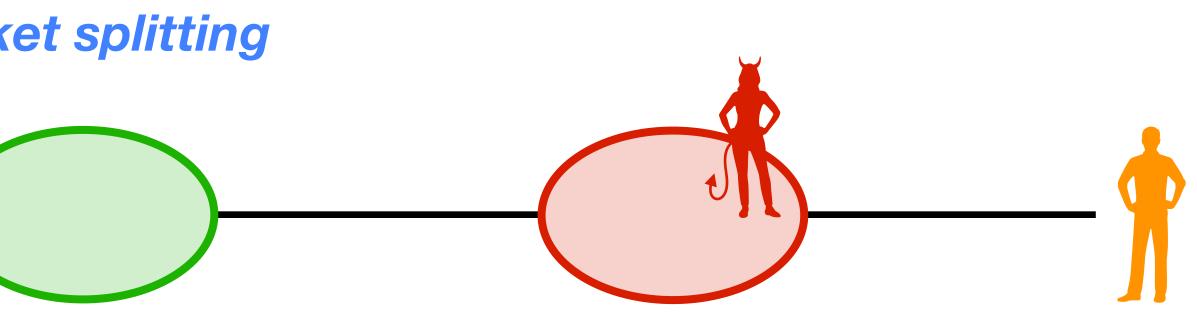










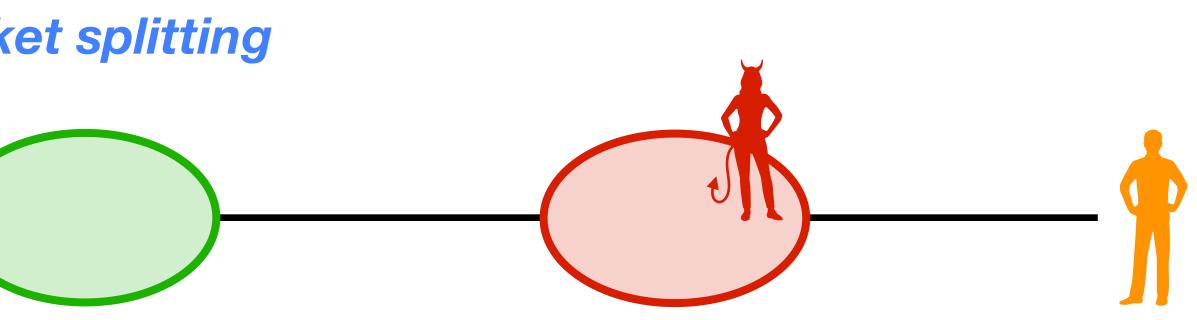


#### When is packet splitting done?









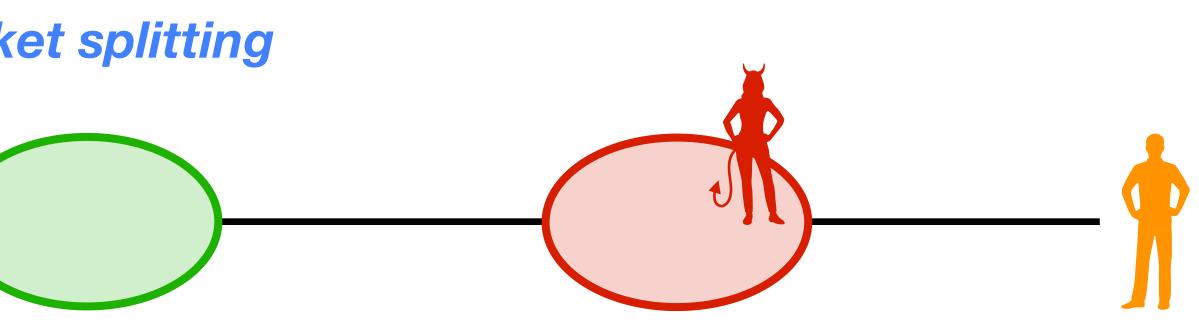
#### When is packet splitting done?

Sender includes splittable packets









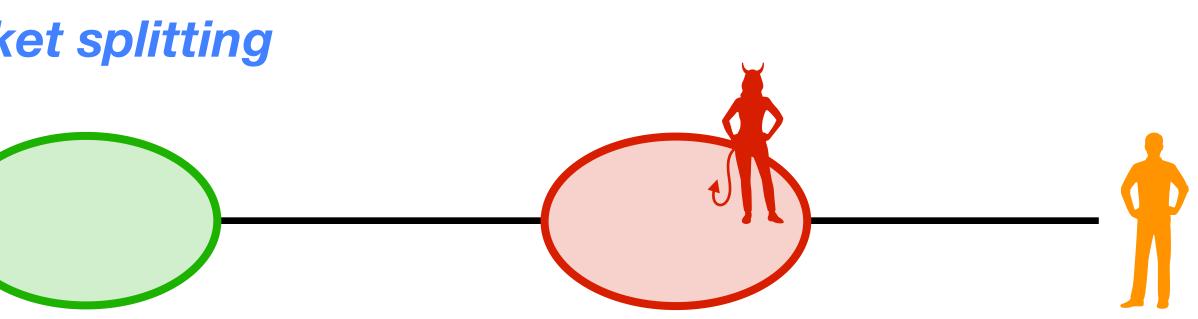
When is packet splitting done?

- Sender includes splittable packets
- ... for each AS on the path









#### When is packet splitting done?

- Sender includes splittable packets
- ... for each AS on the path
- ... at random intervals







#### How does splitting work concretely?



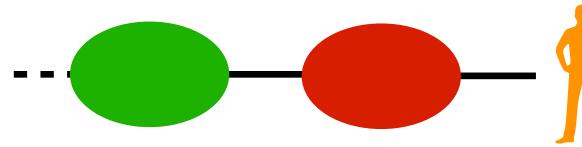


## 



#### How does splitting work concretely?

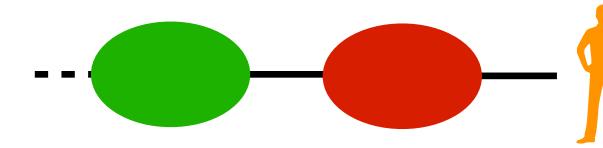






R	MAC	R	MAC	MAC	

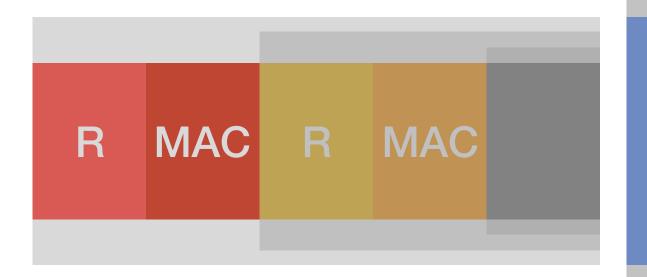




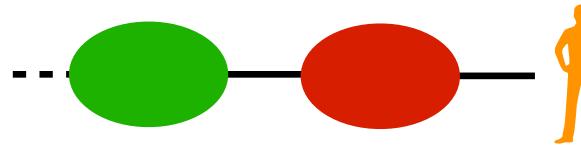
#### Payload









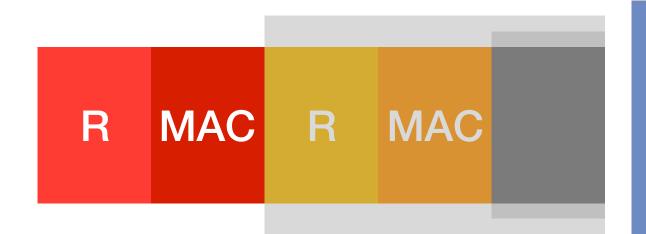


#### Payload

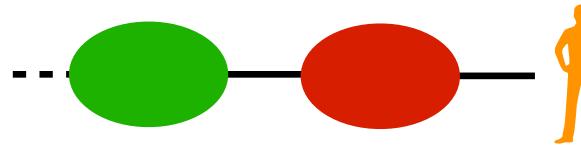








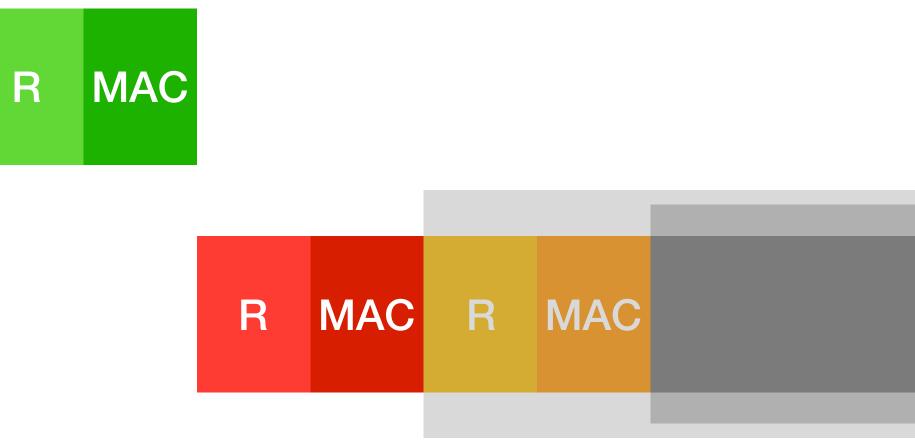




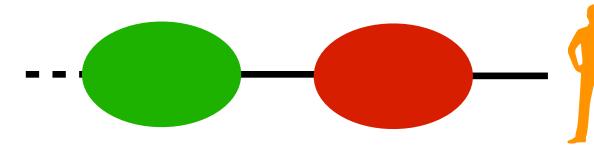
#### Payload











#### Payload

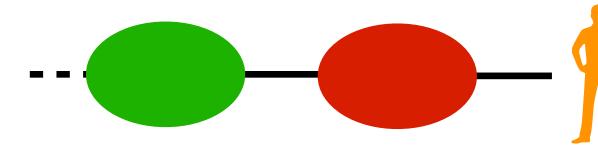






MAC
-----





#### Payload







	R	MAC	R	MAC	
--	---	-----	---	-----	--

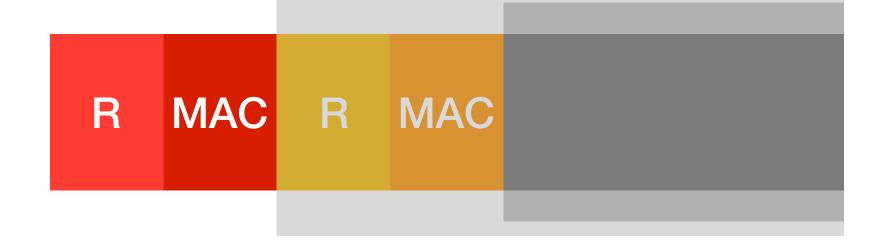




	R	MAC	R	MAC	
--	---	-----	---	-----	--

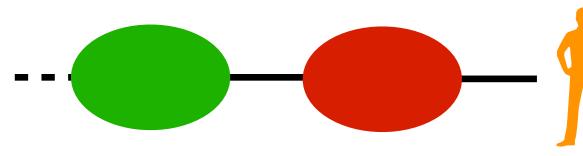






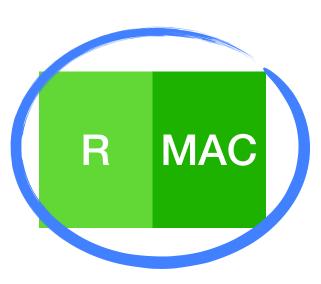


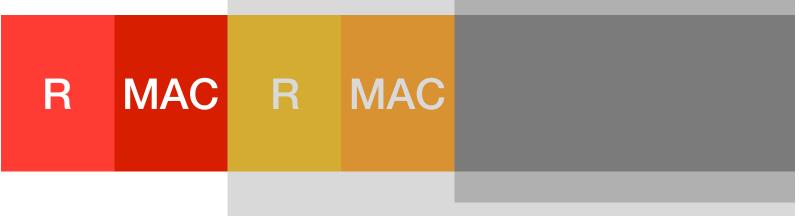






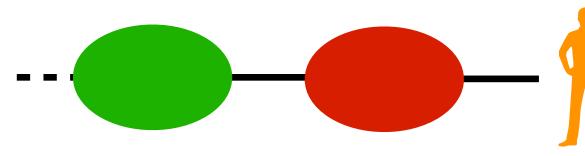






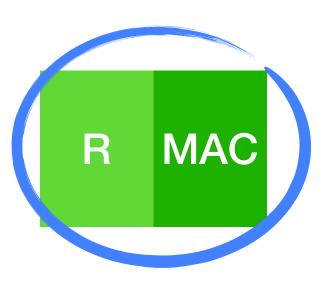


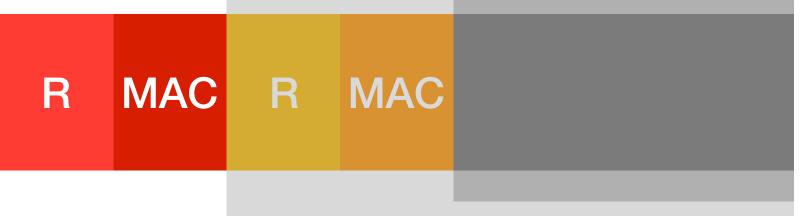






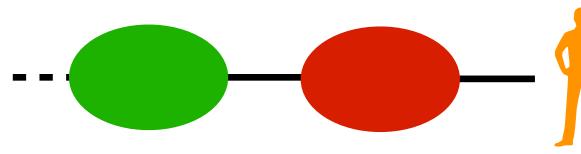










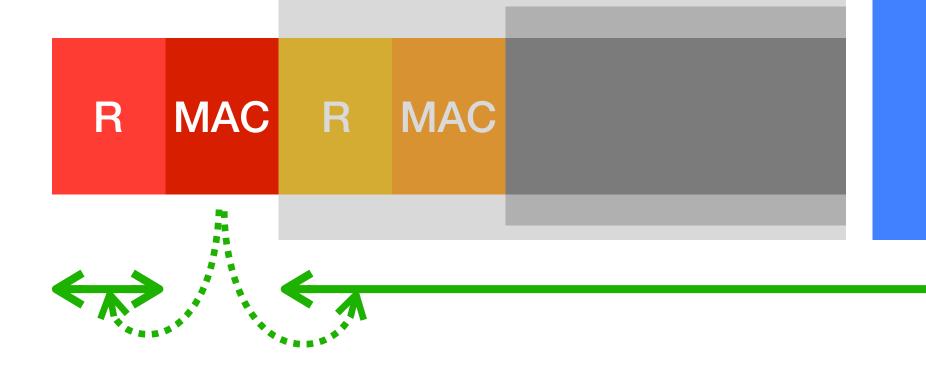


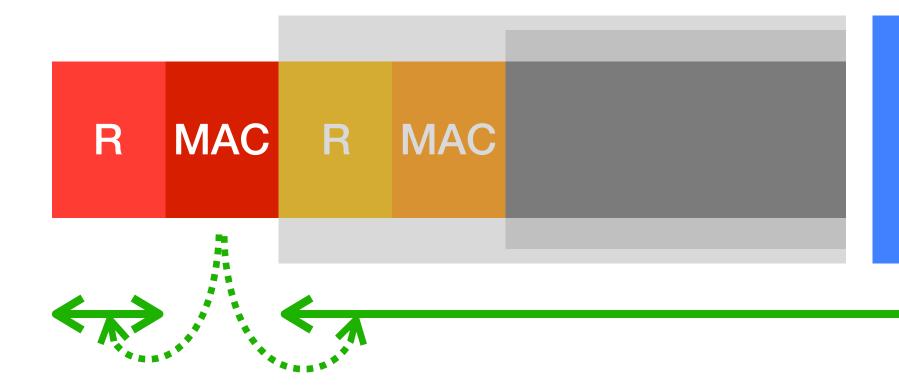
#### **Pseudorandom Payload**

#### **Pseudorandom Payload**

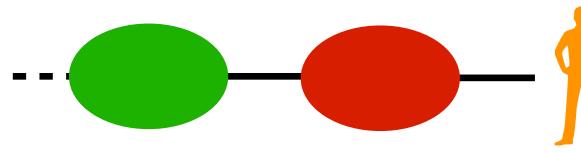












#### **Pseudorandom Payload**

#### **Pseudorandom Payload**

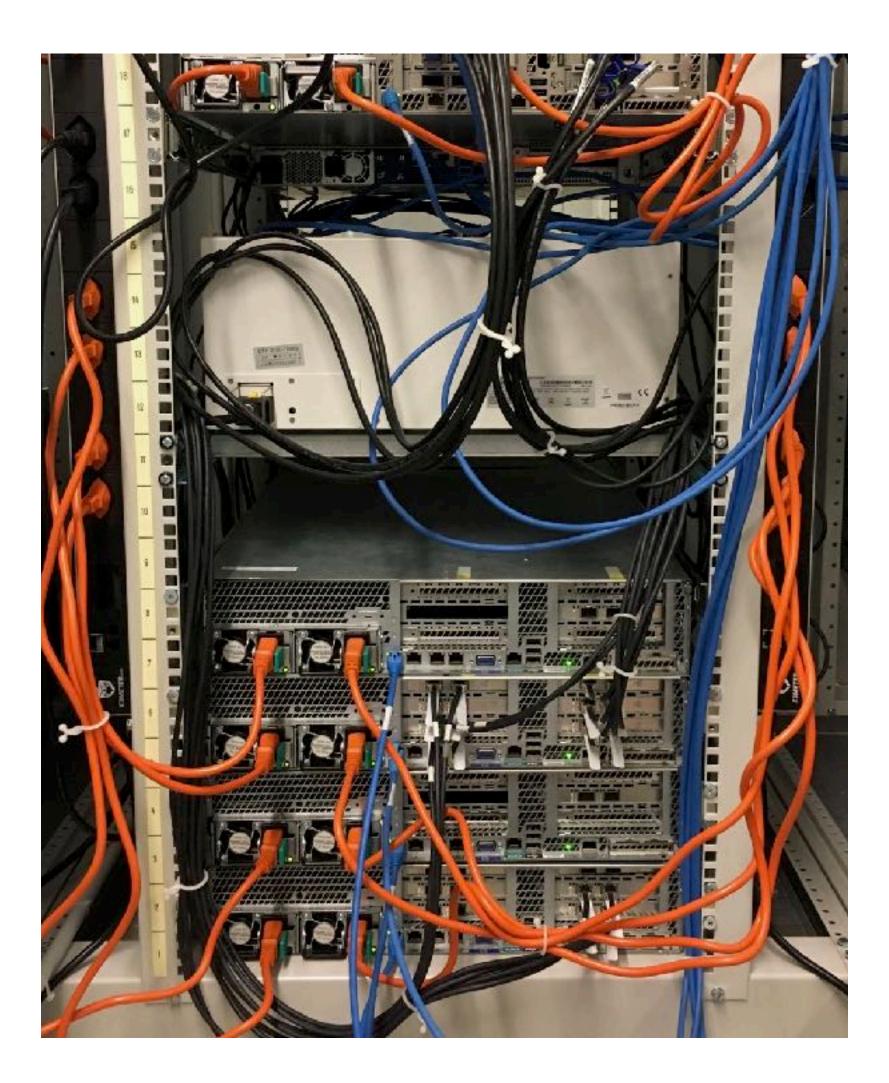


## **TARANET Performance**Evaluation setup, Throughput, Latency





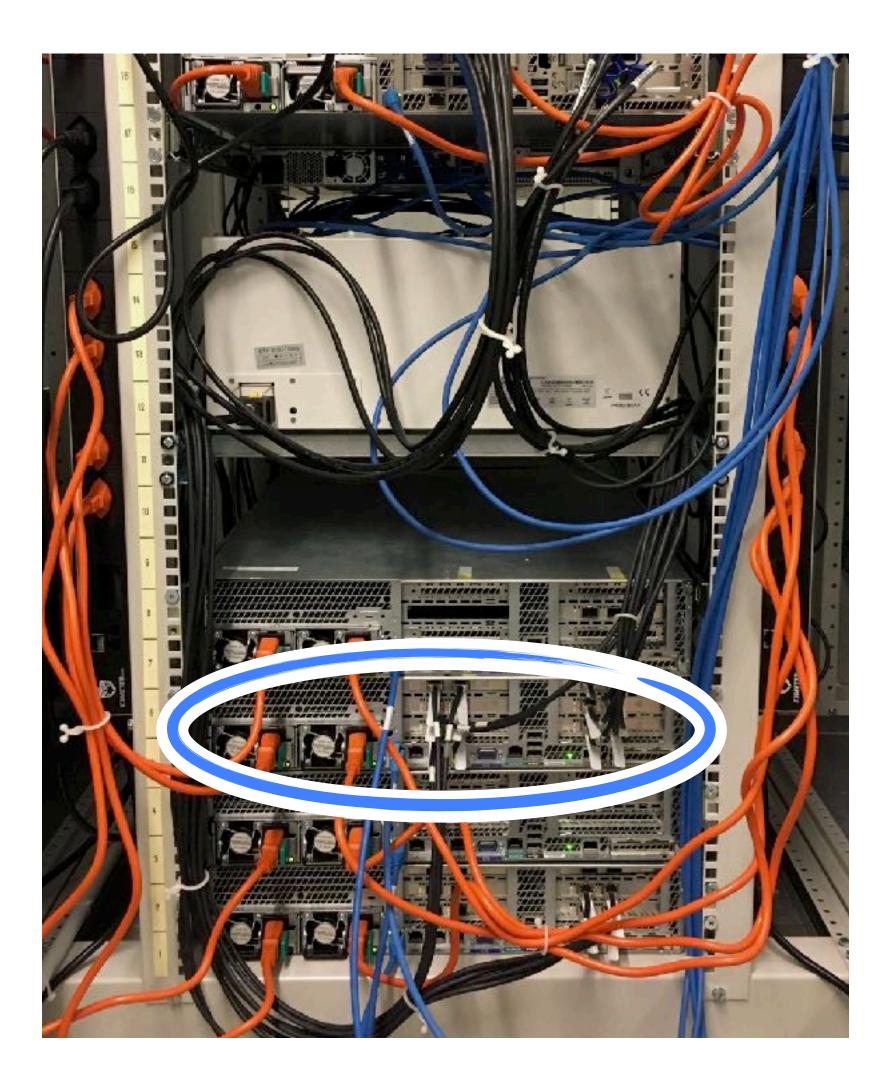
### **Evaluation setup**





- Prototype implementation
  - Data-Plane Development Kit (DPDK)

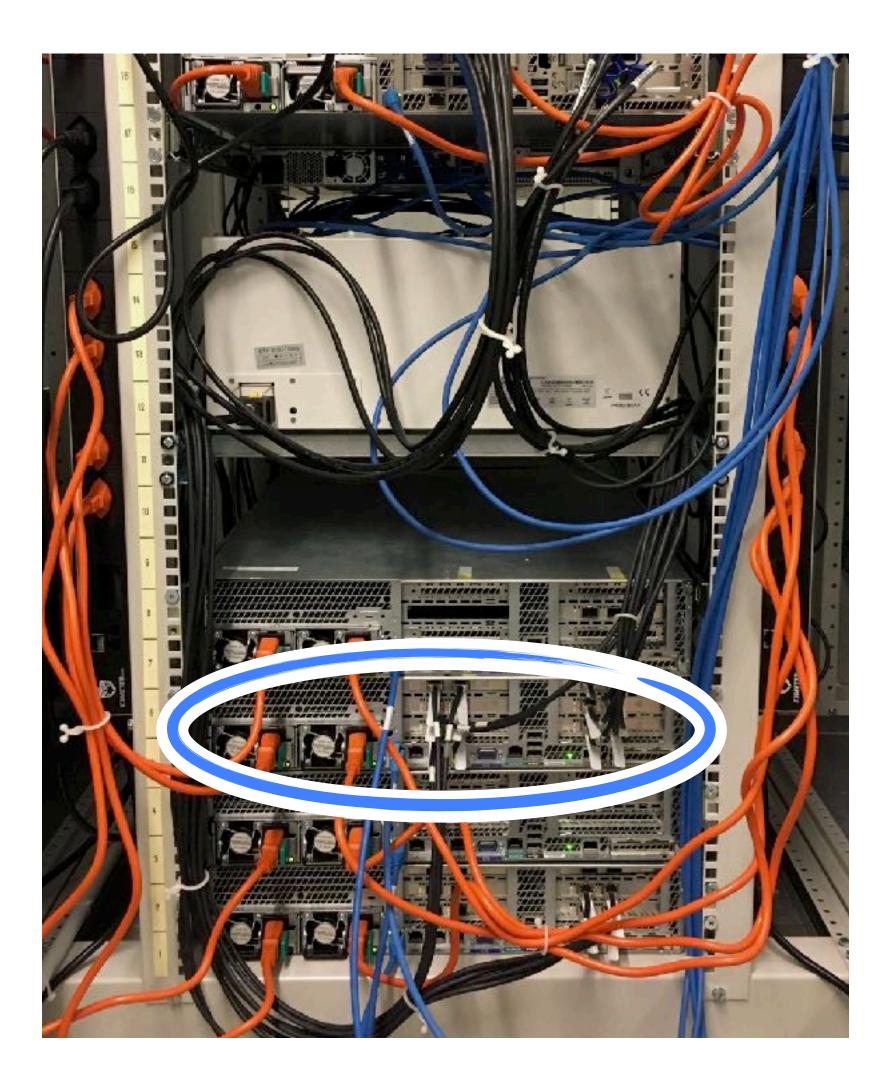






- Prototype implementation
  - Data-Plane Development Kit (DPDK)
- Software router
  - 12x 10 GbE NICs
  - Intel Xeon 2.7 GHz (2x 8 cores)

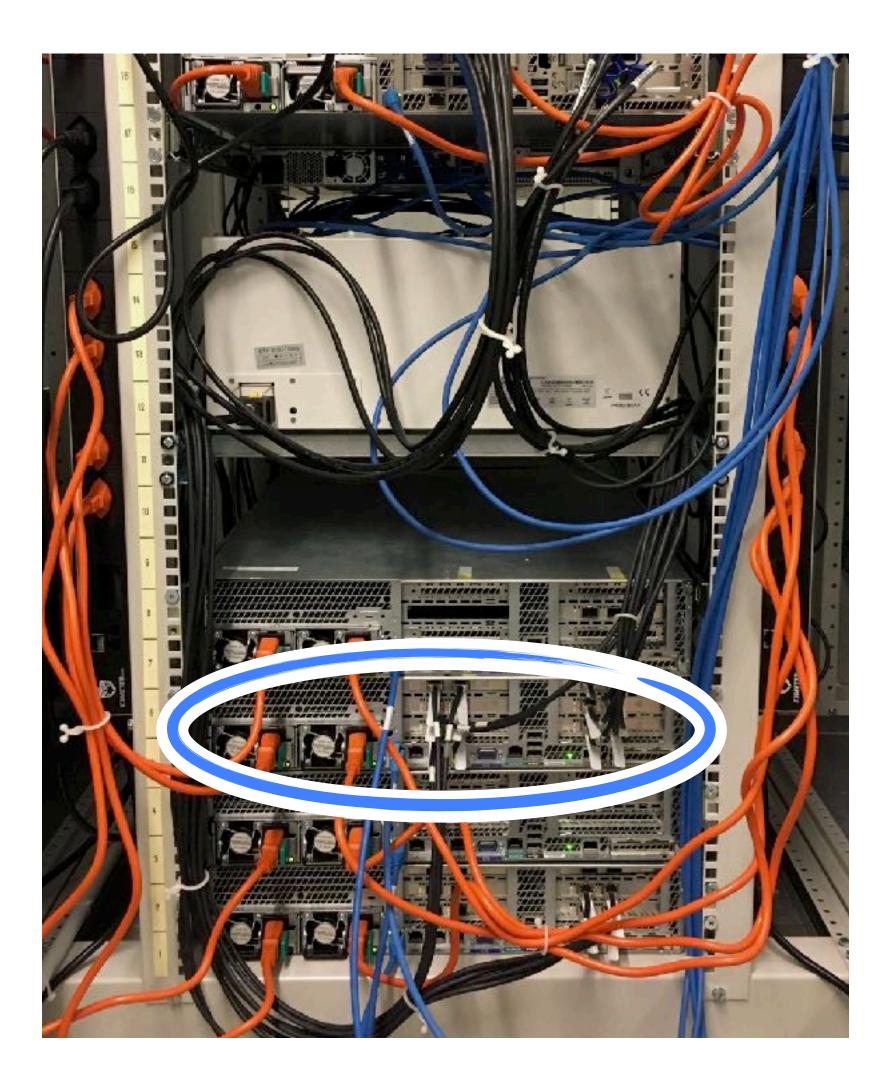






- Prototype implementation
  - Data-Plane Development Kit (DPDK)
- Software router
  - 12x 10 GbE NICs
  - Intel Xeon 2.7 GHz (2x 8 cores)
- Results in this presentation

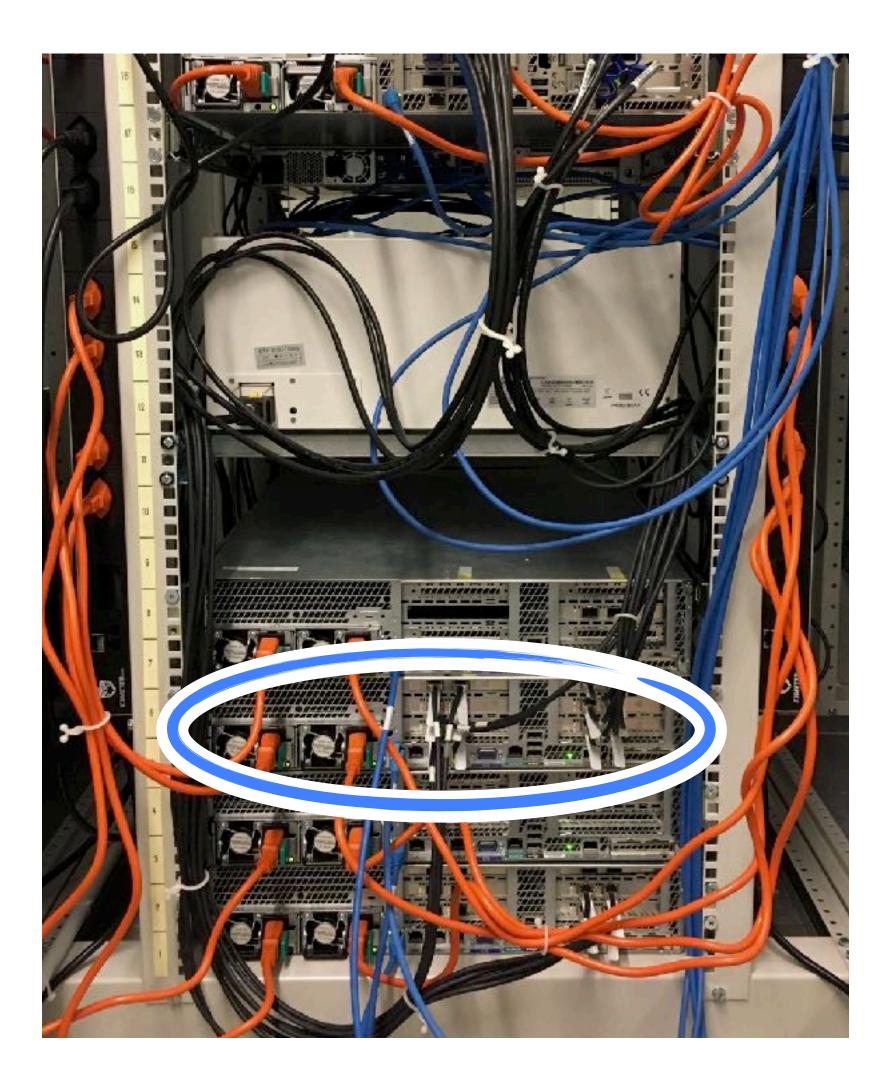






- Prototype implementation
  - Data-Plane Development Kit (DPDK)
- Software router
  - 12x 10 GbE NICs
  - Intel Xeon 2.7 GHz (2x 8 cores)
- Results in this presentation
  - Performance of one node

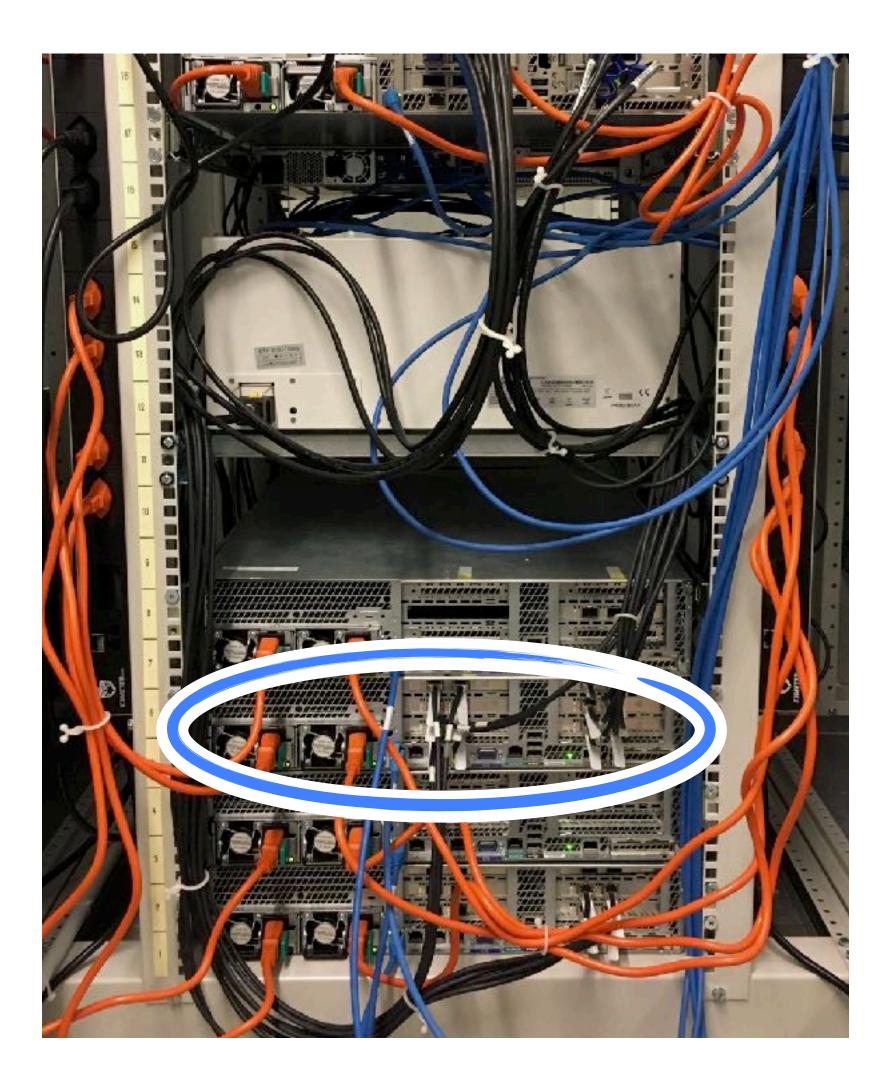






- Prototype implementation
  - Data-Plane Development Kit (DPDK)
- Software router
  - 12x 10 GbE NICs
  - Intel Xeon 2.7 GHz (2x 8 cores)
- Results in this presentation
  - Performance of one node
  - Single 10 GbE interface



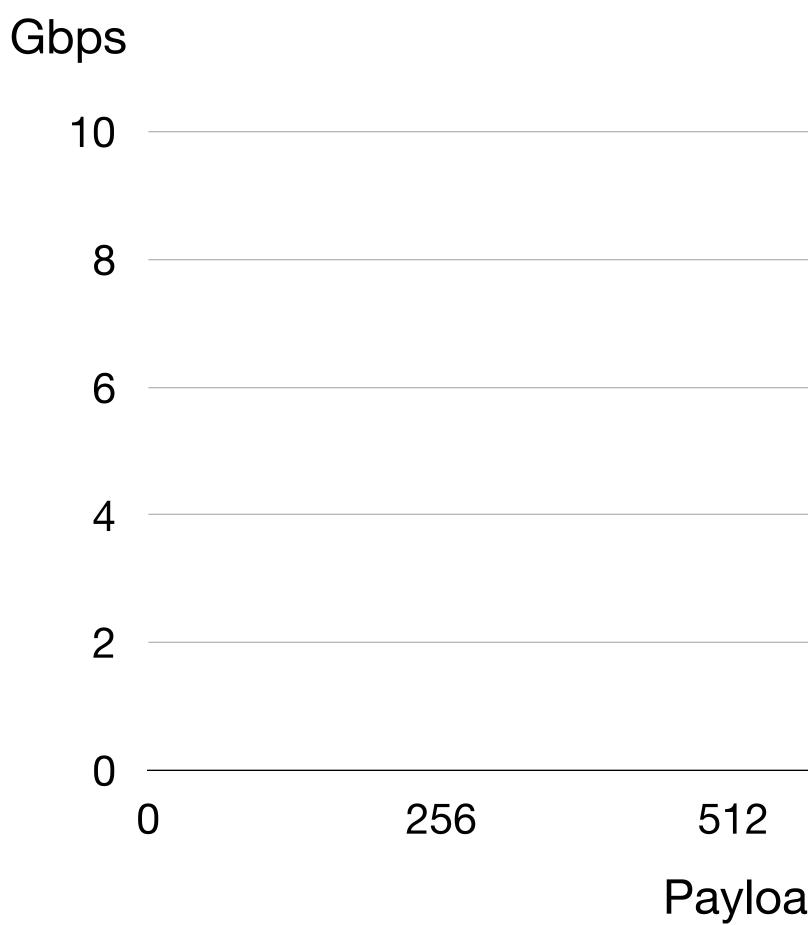




- Prototype implementation
  - Data-Plane Development Kit (DPDK)
- Software router
  - 12x 10 GbE NICs
  - Intel Xeon 2.7 GHz (2x 8 cores)
- Results in this presentation
  - Performance of one node
  - Single 10 GbE interface
  - Single processing core



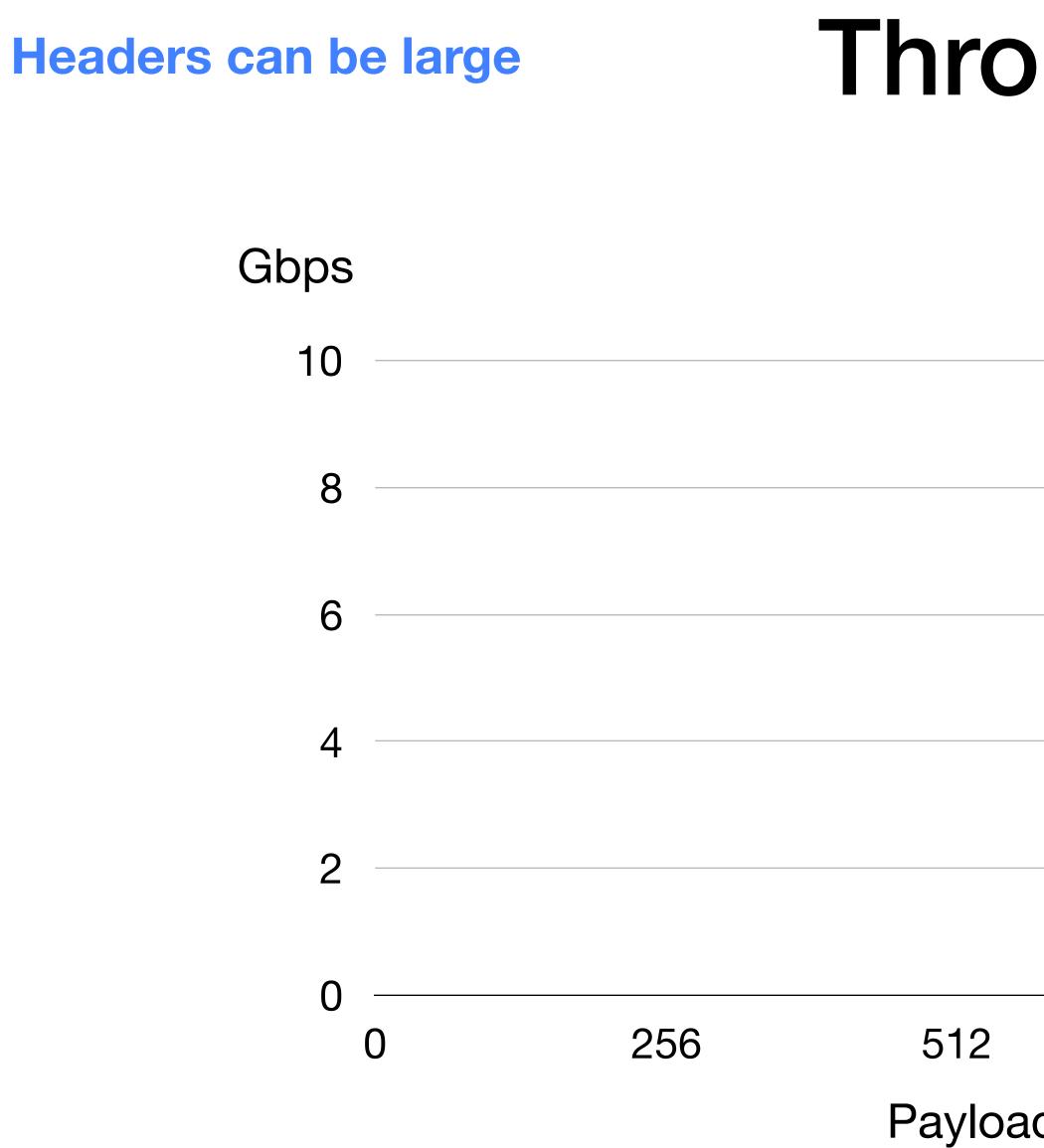
# Throughput





768	1024	1280
ad Size (bytes)		





**ETH** zürich

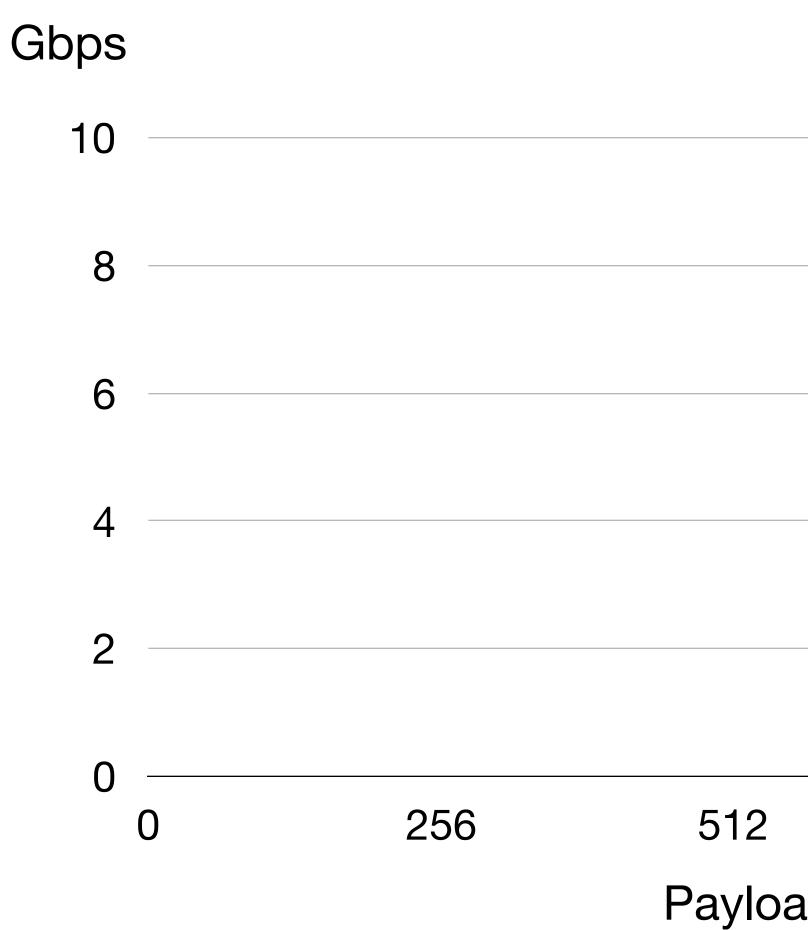
# Throughput

768	1024	1280
ad Size (bytes)		









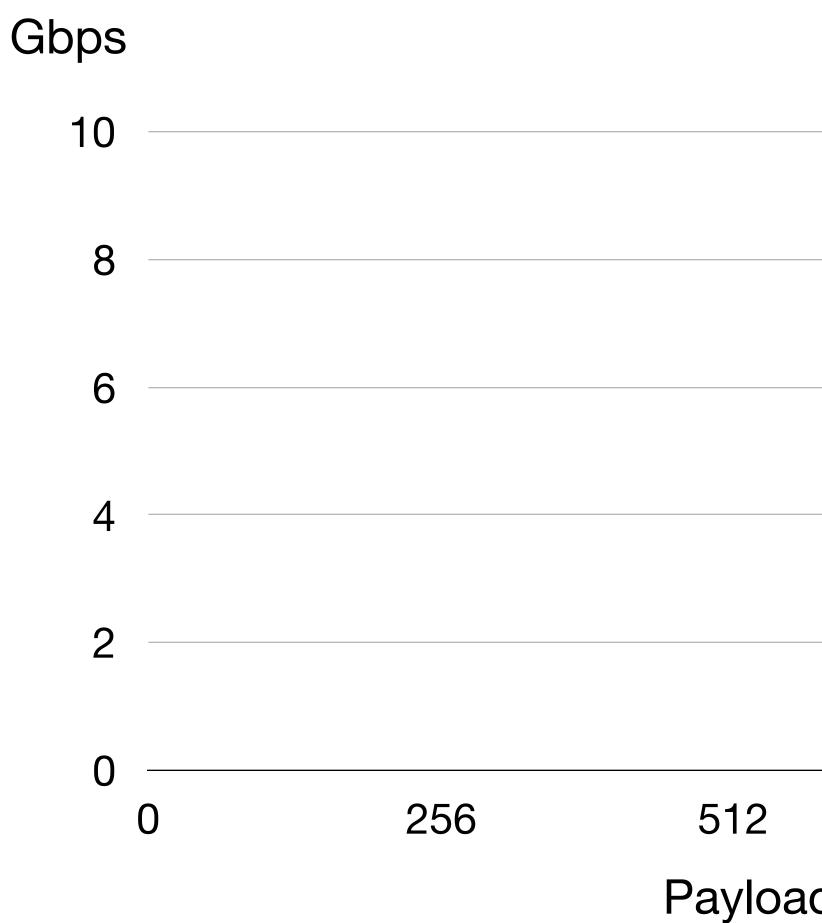


768	1024	1280
ad Size (bytes)		





#### Path length: 7 nodes



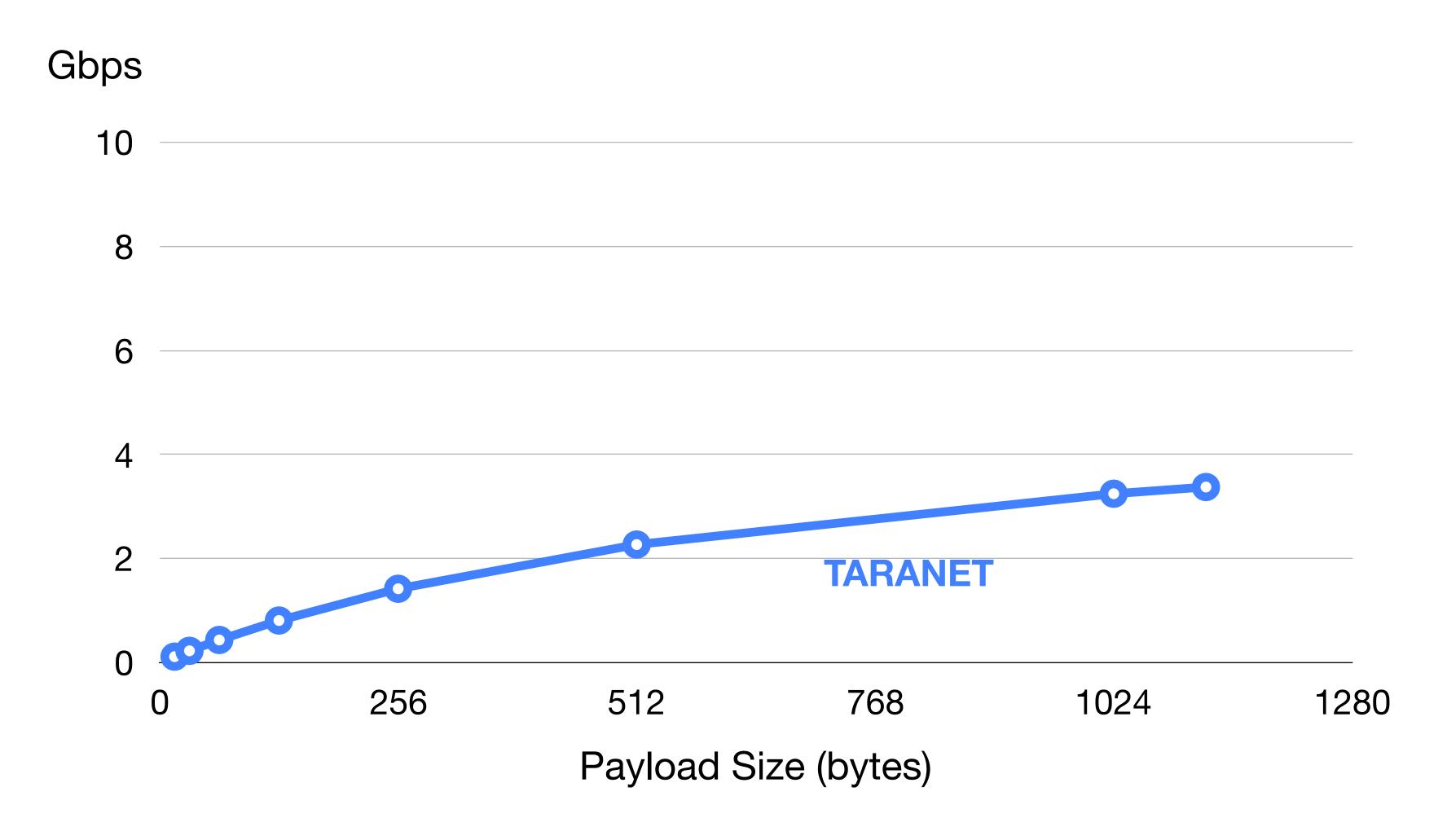


768	1024	1280
ad Size (bytes)		





#### Path length: 7 nodes

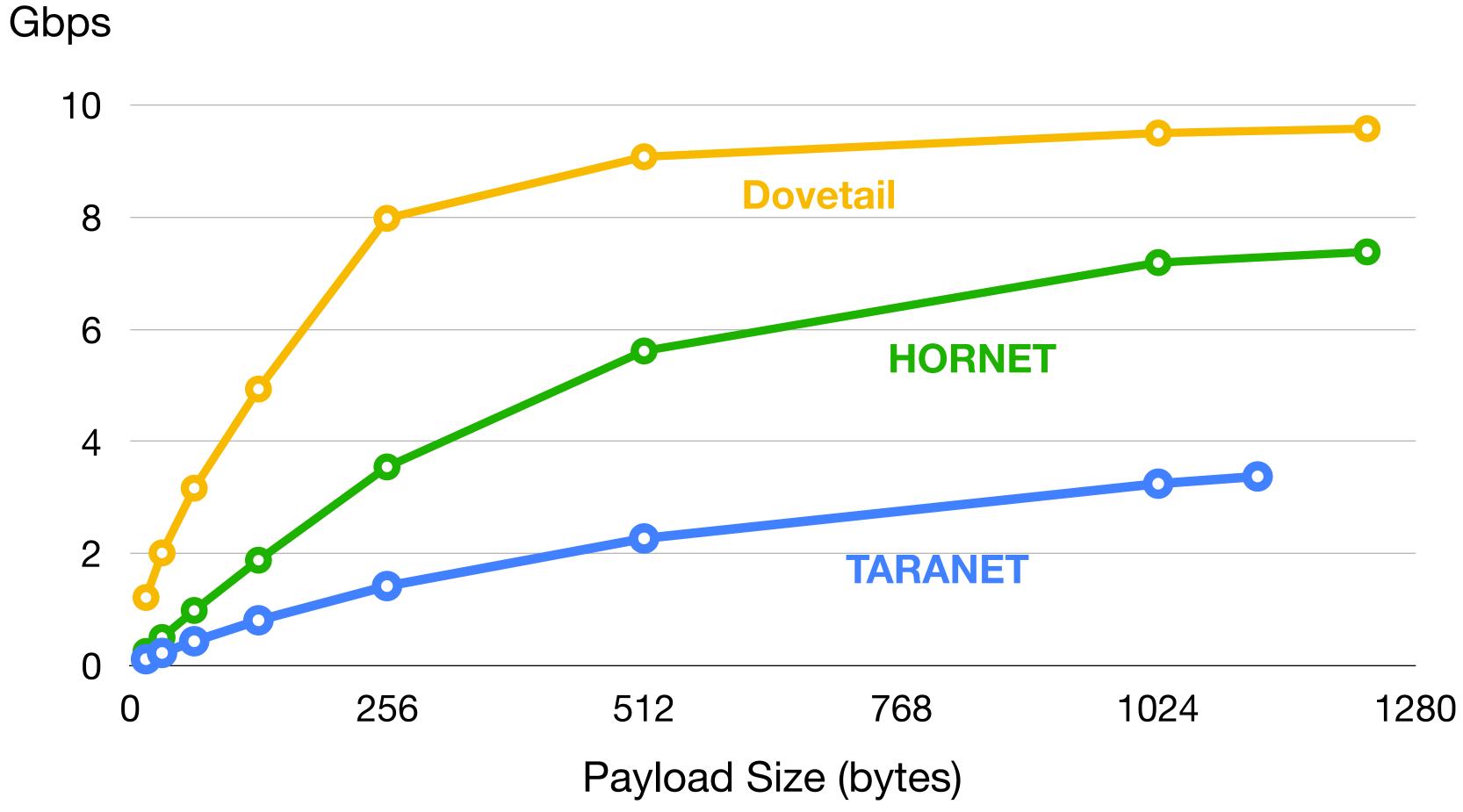








#### Path length: 7 nodes

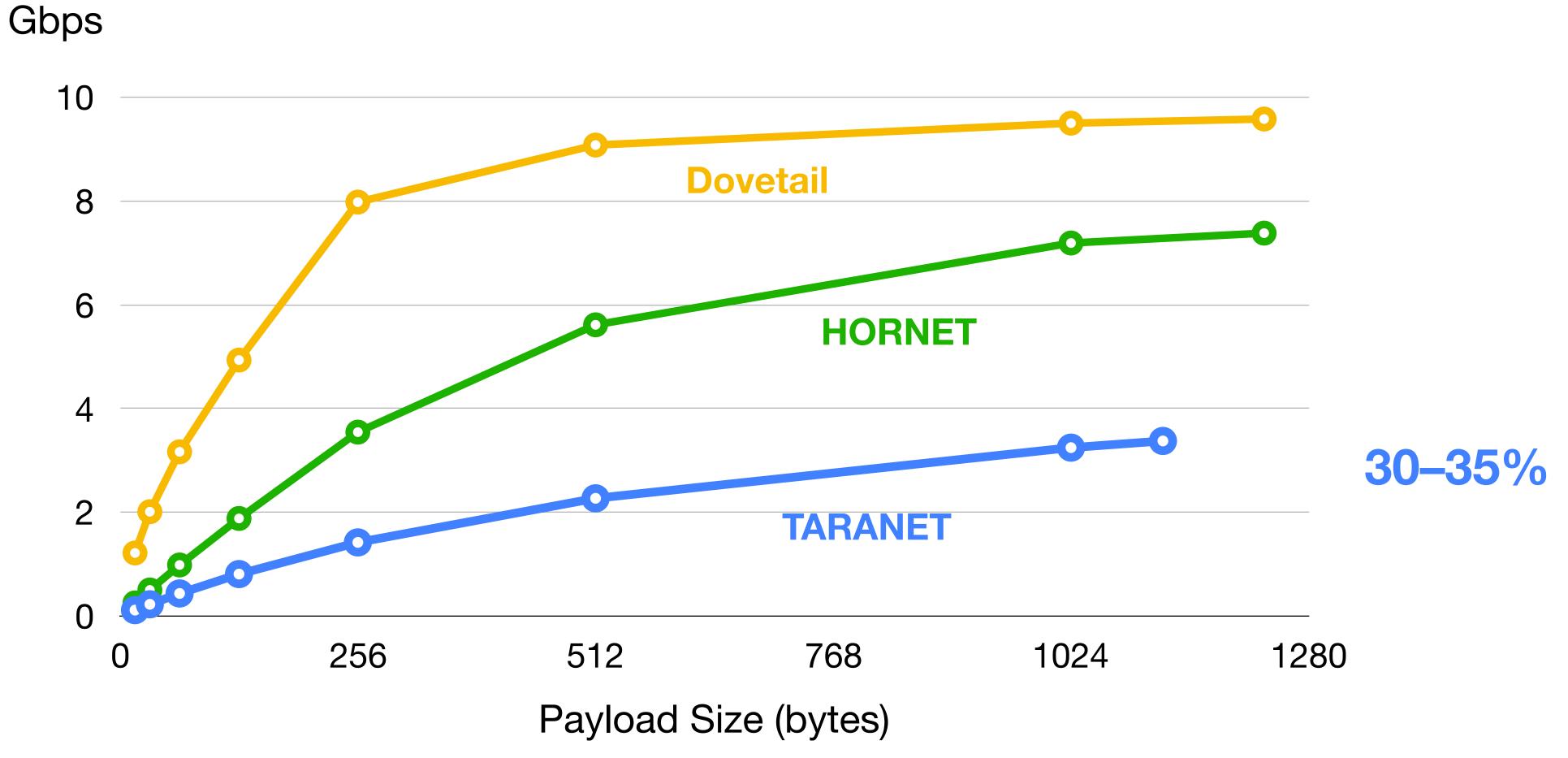






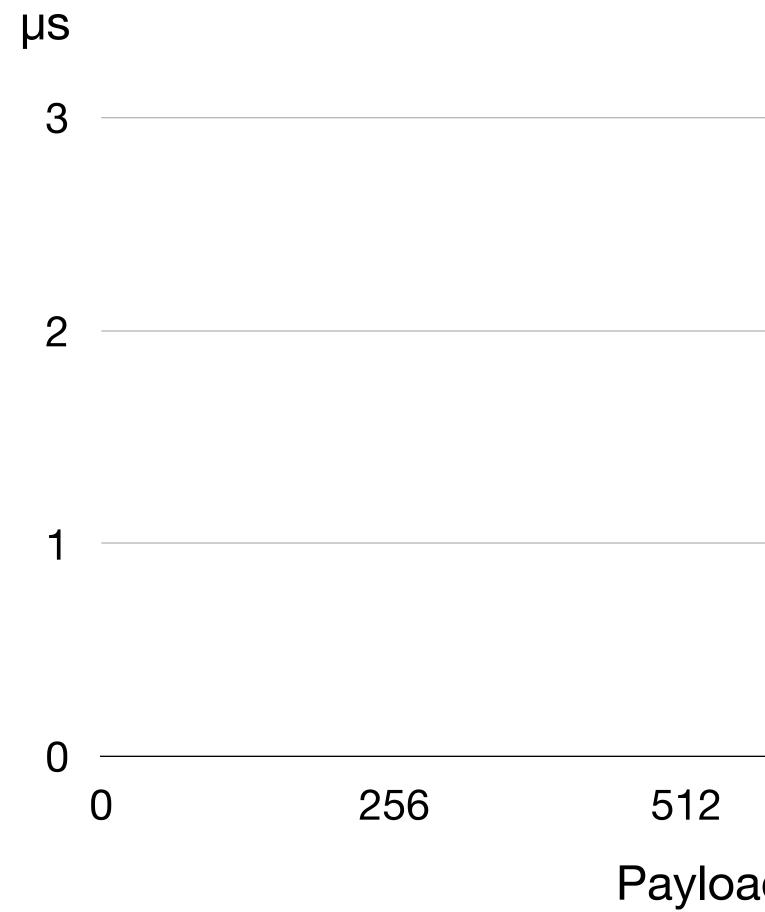


#### Path length: 7 nodes







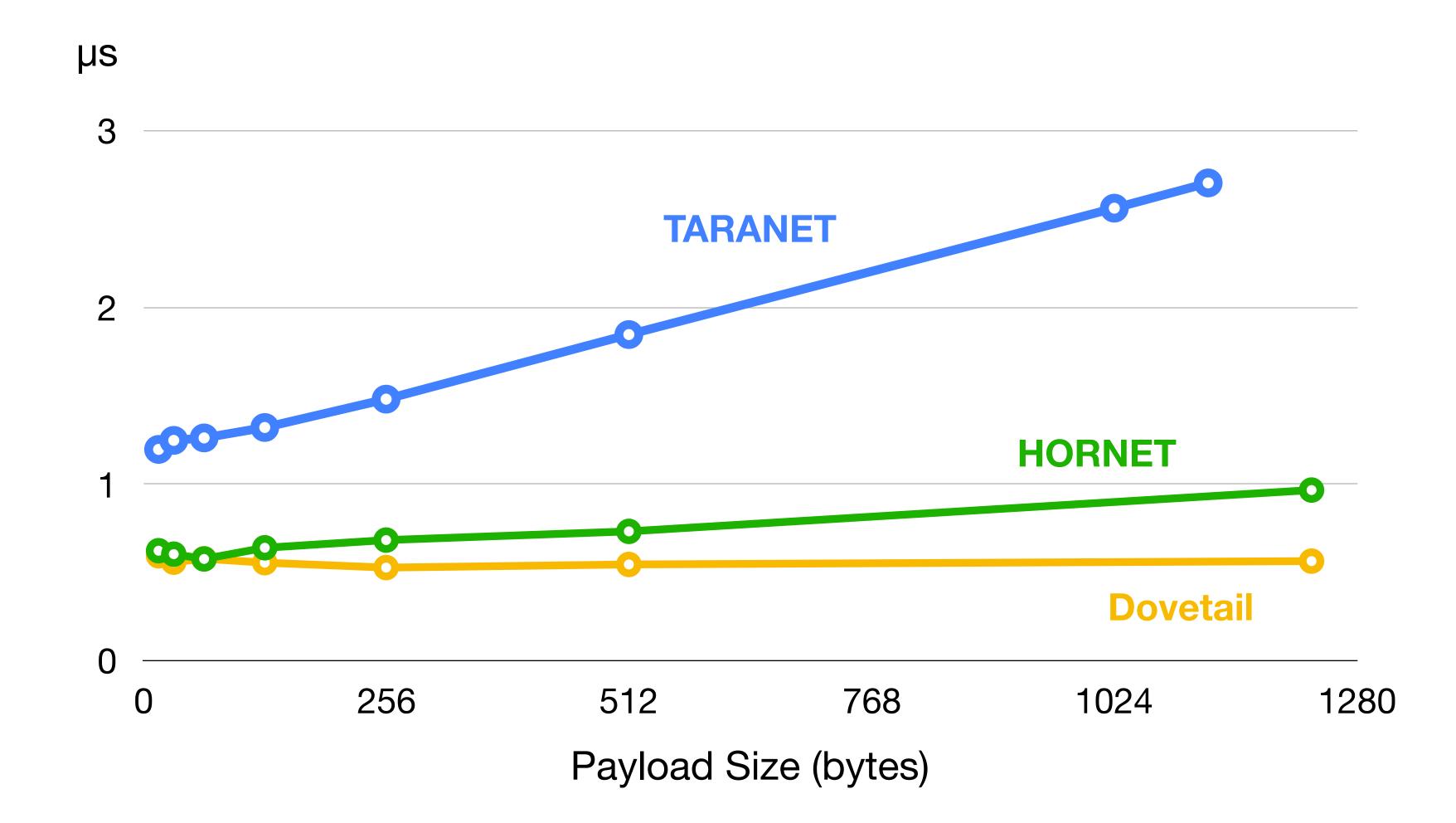




## Latency

768	1024	1280
ad Size (bytes)		



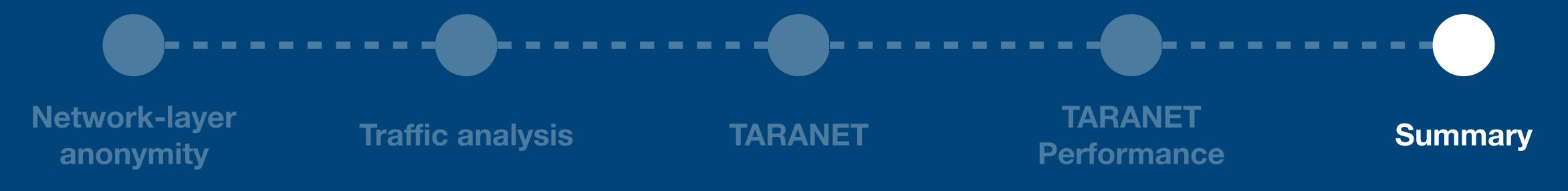




## Latency



### Summary Highlights, Limitations



- TARANET highlights:
  - Protection against passive traffic analysis with *flowlets*
  - Protection against active traffic analysis with packet splitting





- TARANET highlights:
  - Protection against passive traffic analysis with *flowlets*
  - Protection against active traffic analysis with packet splitting
  - Good performance 3 Gbps on single core, acceptable latency





- TARANET highlights:
  - Protection against passive traffic analysis with *flowlets*
  - Protection against active traffic analysis with packet splitting
  - Good performance <u>3 Gbps</u> on single core, acceptable latency
- Limitations: lacksquare





- TARANET highlights:
  - Protection against passive traffic analysis with *flowlets*
  - Protection against active traffic analysis with packet splitting
  - Good performance 3 Gbps on single core, acceptable latency
- Limitations:
  - Chaff traffic creates a non-negligible bandwidth overhead





- TARANET highlights:
  - Protection against passive traffic analysis with *flowlets*
  - Protection against active traffic analysis with packet splitting
  - Good performance 3 Gbps on single core, acceptable latency
- Limitations:
  - Chaff traffic creates a non-negligible bandwidth overhead
  - Third-party anonymity





# In the paper

#### **TARANET: Traffic-Analysis Resistant Anonymity** at the Network Layer

Chen Chen chenche1@andrew.cmu.edu Carnegie Mellon University

David Barrera david.barrera@polymtl.ca Polytechnique Montreal

Daniele E. Asoni daniele.asoni@inf.ethz.ch ETH Zürich

George Danezis g.danezis@ucl.ac.uk University College London

Adrian Perrig adrian.perrig@inf.ethz.ch ETH Zürich

Carmela Troncoso carmela.troncoso@epfl.ch EPFL

Abstract—Modern low-latency anonymity systems, no matter whether constructed as an overlay or implemented at the network layer, offer limited security guarantees against traffic analysis. On the other hand, high-latency anonymity systems offer strong security guarantees at the cost of computational overhead and long delays, which are excessive for interactive applications. We propose TARANET, an anonymity system that implements protection against traffic analysis at the network layer, and limits the incurred latency and overhead. In TARANET's setup phase, traffic analysis is thwarted by mixing. In the data transmission phase, end hosts and ASes coordinate to shape traffic into constant-rate transmission using packet splitting. Our prototype implementation shows that TARANET can forward anonymous traffic at over 50 Gbps using commodity hardware.

#### **1. Introduction**

**ETH** zürich

Users are increasingly aware of their lack of privacy and

in forwarding anonymous traffic. Intermediate anonymity supporting network nodes (or nodes for short) first cooperate with senders to establish anonymous sessions or circuits, and then process and forward traffic from those senders to receivers. While these systems achieve high throughput and low latency, the security guarantees of these systems are no stronger than Tor's. Moreover, LAP and Dovetail leak the position of intermediate nodes on the path and the total path length, which reduces the anonymity set size, facilitating deanonymization [21].

The problem space appears to have an unavoidable tradeoff: strong anonymity appears achievable only through *drastically higher overhead* [27]. In this paper, we aim to push the boundaries of this anonymity/performance tradeoff by combining the speed of network-layer anonymity systems with strong defenses.

To improve the anonymity guarantees, traffic analysis attacks need to be prevented, or made significantly harder for the adversary to perform. The common method to achieve

- Flowlet setup (asymm. crypto)
- Link padding (security in depth)
- Anonymity set size analysis
- Security analysis

- Chaff/setup packet trade-off
- **Deployment incentives**





**Contacts:** 

Chen Chen: chenche1@andrew.cmu.edu Daniele E. Asoni: daniele.asoni@inf.ethz.ch

# Thank you!